

Basic Concepts in Information Theory

Spring 2024

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Coding theory

- Let us start with discrete random variables

Coding theory

- How to represent the information contained in the random variables?

$$h(\mathbf{x}) \geq 0$$

$$h(\mathbf{x}, \mathbf{y}) = h(\mathbf{x}) + h(\mathbf{y}) \quad \mathbf{x}, \mathbf{y} \text{ are independent}$$

$$p(\mathbf{x}, \mathbf{y}) = p(\mathbf{x})p(\mathbf{y})$$



$$h(\mathbf{x}) = -\log(p(\mathbf{x}))$$

Entropy

- The average amount of information needed to transmit

$$H(\mathbf{x}) = - \sum_{\mathbf{x}} p(\mathbf{x}) \log (p(\mathbf{x}))$$

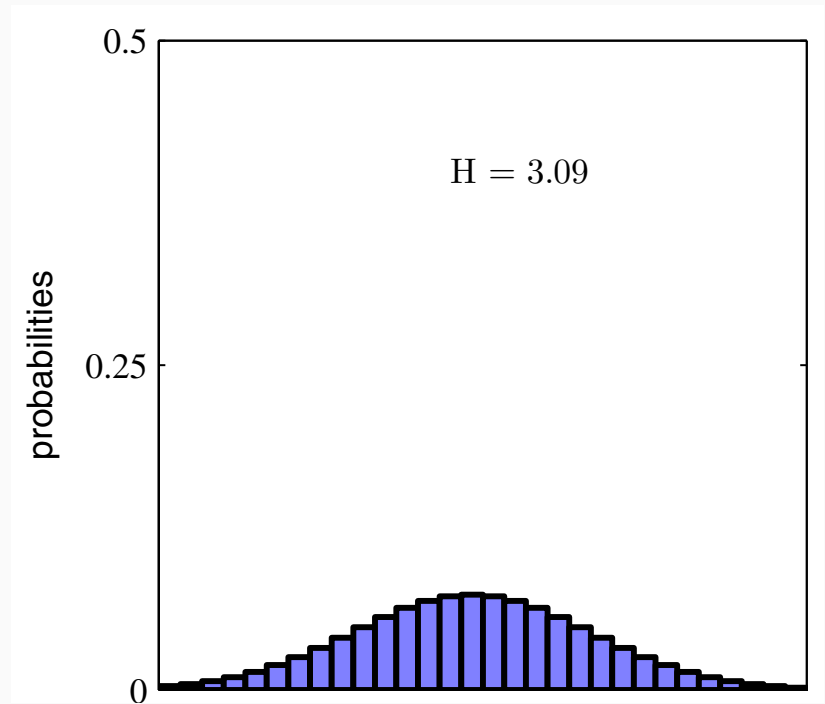
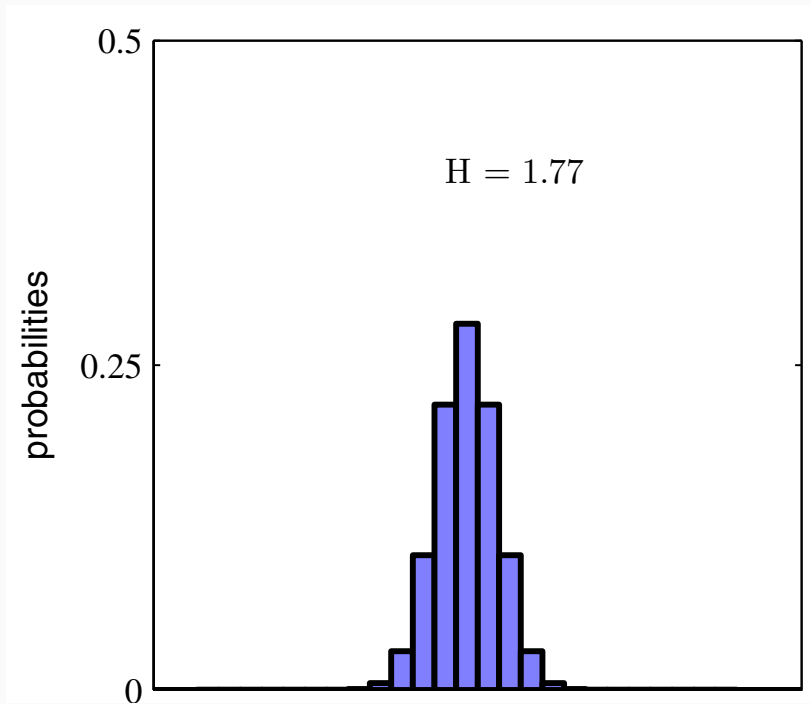
Entropy

x	a	b	c	d	e	f	g	h
$p(x)$	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{8}$	$\frac{1}{16}$	$\frac{1}{64}$	$\frac{1}{64}$	$\frac{1}{64}$	$\frac{1}{64}$

$$\begin{aligned} H[x] &= -\frac{1}{2} \log_2 \frac{1}{2} - \frac{1}{4} \log_2 \frac{1}{4} - \frac{1}{8} \log_2 \frac{1}{8} - \frac{1}{16} \log_2 \frac{1}{16} - \frac{4}{64} \log_2 \frac{1}{64} \\ &= 2 \text{ bits} \end{aligned}$$


Entropy is also the average code length


Entropy reflects uncertainty



Maximum entropy

- Consider a discrete R.V. with M possible status. We want to find the distribution has the the maximum entropy $H[p] = - \sum_i p(x_i) \ln p(x_i)$.


$$\tilde{H} = - \sum_i p(x_i) \ln p(x_i) + \lambda \left(\sum_i p(x_i) - 1 \right)$$

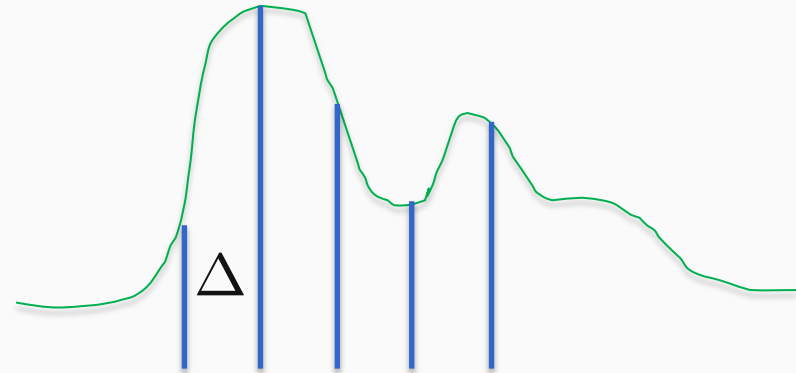

$$p(x_i) = 1/M \quad \text{uniform distribution}$$

Differential entropy

- Entropy is naturally defined on discrete random variables.
- But how about continuous variables?

Differential entropy

- Let us divide x into bins of Δ



Mean-value theorem


$$\int_{i\Delta}^{(i+1)\Delta} p(x) dx = p(x_i)\Delta$$

Entropy on discretized probability



$$H_{\Delta} = - \sum_i p(x_i)\Delta \ln (p(x_i)\Delta) = - \sum_i p(x_i)\Delta \ln p(x_i) - \ln \Delta$$

$$\sum_i p(x_i)\Delta = 1$$

Differential entropy

$$H_{\Delta} = - \sum_i p(x_i) \Delta \ln (p(x_i) \Delta) = - \sum_i p(x_i) \Delta \ln p(x_i) - \ln \Delta$$


Goes to infinity
Throw out it

$$\lim_{\Delta \rightarrow 0} \left\{ \sum_i p(x_i) \Delta \ln p(x_i) \right\} = \int p(x) \ln p(x) dx$$


$$H[\mathbf{x}] = - \int p(\mathbf{x}) \ln p(\mathbf{x}) d\mathbf{x}$$

Differential entropy

- The term that is thrown out reflects that to specify a continuous variable very precisely requires many many bits
- Note: differential entropy can be negative!

Differential entropy

- Given a continuous variable x with mean μ and variance σ^2 , which distribution has the largest entropy?

$$\int_{-\infty}^{\infty} p(x) dx = 1$$

$$\int_{-\infty}^{\infty} xp(x) dx = \mu$$

$$\int_{-\infty}^{\infty} (x - \mu)^2 p(x) dx = \sigma^2.$$

Differential entropy

$$\begin{aligned} \max \quad & - \int_{-\infty}^{\infty} p(x) \ln p(x) dx + \lambda_1 \left(\int_{-\infty}^{\infty} p(x) dx - 1 \right) \\ & + \lambda_2 \left(\int_{-\infty}^{\infty} xp(x) dx - \mu \right) + \lambda_3 \left(\int_{-\infty}^{\infty} (x - \mu)^2 p(x) dx - \sigma^2 \right) \end{aligned}$$



$$p(x) = \frac{1}{(2\pi\sigma^2)^{1/2}} \exp \left\{ -\frac{(x - \mu)^2}{2\sigma^2} \right\}$$

Gaussian distribution!

Conditional entropy

- Given \mathbf{x} , how much information is left for \mathbf{y}

$$H[\mathbf{y}|\mathbf{x}] = - \iint p(\mathbf{y}, \mathbf{x}) \ln p(\mathbf{y}|\mathbf{x}) d\mathbf{y} d\mathbf{x}$$

$$H[\mathbf{x}, \mathbf{y}] = H[\mathbf{y}|\mathbf{x}] + H[\mathbf{x}] \quad \text{Prove it by yourself}$$

Kullback-Leibler (KL) divergence

- Also called relative entropy

$$\begin{aligned} \text{KL}(p||q) &= - \int p(\mathbf{x}) \ln q(\mathbf{x}) \, d\mathbf{x} - \left(- \int p(\mathbf{x}) \ln p(\mathbf{x}) \, d\mathbf{x} \right) \\ &= - \int p(\mathbf{x}) \ln \left\{ \frac{q(\mathbf{x})}{p(\mathbf{x})} \right\} \, d\mathbf{x}. \end{aligned}$$

If we use q to transmit information for p , how much extra information do we need

Kullback-Leibler (KL) divergence

- KL divergence is widely used to measure the difference between two distributions

$$\text{KL}(p||q) \geq 0 \quad =0 \text{ iff } p = q$$

Prove it with convexity
And Jensen's inequality

- However, it is not symmetric!

$$\text{KL}(p||q) \neq \text{KL}(q||p)$$

KL Divergence

- KL divergence plays the key role in approximate inference
- All the deterministic approximate methods aim to minimize the KL divergence between the true and approximate posteriors (or in the reversed direction)
- In general, we have alpha divergence
- We will discuss these in detail later

Mutual information

How many information do the two random variables share?

$$\begin{aligned} I[\mathbf{x}, \mathbf{y}] &\equiv \text{KL}(p(\mathbf{x}, \mathbf{y}) \| p(\mathbf{x})p(\mathbf{y})) \\ &= - \iint p(\mathbf{x}, \mathbf{y}) \ln \left(\frac{p(\mathbf{x})p(\mathbf{y})}{p(\mathbf{x}, \mathbf{y})} \right) d\mathbf{x} d\mathbf{y} \end{aligned}$$



$$I[\mathbf{x}, \mathbf{y}] = H[\mathbf{x}] - H[\mathbf{x}|\mathbf{y}] = H[\mathbf{y}] - H[\mathbf{y}|\mathbf{x}]$$

Prove it by
yourself

What you need to know

- Definition of entropy
- How is differential entropy is derived
- Entropy is an indicator for uncertainty
- KL divergence and properties (especially asymmetric)
- Mutual information