Week 11: Lecture A Secure Authentication

Tuesday, November 5, 2024



Stefan Nagy

Project 3: WebSec released

Deadline: this Thursday, November 7th by 11:59PM

Project 3: Web Security

Deadline: Thursday, November 7 by 11:59PM.

Before you start, review the course syllabus for the Lateness, Collaboration, and Ethical Use policies.

You may optionally work alone, or in teams of **at most two** and submit **one project per team**. If you have difficulties forming a team, post on **Piazza's Search for Teammates** forum. Note that the final exam will cover project material, so you and your partner should collaborate on each part.

The code and other answers your group submits must be entirely your own work, and you are bound by the University's Student Code. You may consult with other students about the conceptualization of the project and the meaning of the questions, but you may not look at any part of someone else's solution or collaborate with anyone outside your group. You may consult published references, provided that you appropriately cite them (e.g., in your code comments). **Don't risk your grade and degree by cheating!**

Complete your work in the **CS 4440 VM**—we will use this same environment for grading. You may not use any **external dependencies**. Use only default Python 3 libraries and/or modules we provide you.

Project 3 progress

Working on Part 1	0%
Finished Part 1, working on Part 2	006
Finished Part 2 working on Part 3	070
	0%
Finished with everything!	0%
Haven't started vet :(
	0%



Start the presentation to see live content. For screen share software, share the entire screen. Get help at **pollev.com/app**

Project 4: NetSec released

Deadline: Thursday, December 5th by 11:59PM

Project 4: Network Security

Deadline: Thursday, December 5 by 11:59PM.

Before you start, review the course syllabus for the Lateness, Collaboration, and Ethical Use policies.

You may optionally work alone, or in teams of **at most two** and submit **one project per team**. If you have difficulties forming a team, post on **Piazza's Search for Teammates** forum. Note that the final exam will cover project material, so you and your partner should collaborate on each part.

The code and other answers your group submits must be entirely your own work, and you are bound by the University's Student Code. You may consult with other students about the conceptualization of the project and the meaning of the questions, but you may not look at any part of someone else's solution or collaborate with anyone outside your group. You may consult published references, provided that you appropriately cite them (e.g., in your code comments). **Don't risk your grade and degree by cheating!**

Complete your work in the **CS 4440 VM**—we will use this same environment for grading. You may not use any **external dependencies**. Use only default Python 3 libraries and/or modules we provide you.

New **Wiki pages** to help you on Project 4:

CS 4440 Wiki: Wireshark Tutorial

Table of Contents:

Below is an abridged cheat sheet of Wireshark fundamentals that you'll use in this course.

This page is by no means comprehensive – we encourage you to bookmark and familiarize yourself with one of the many in-depth Wireshark tutorials on the web. Some great examples are:

• Wireshark User's Guide

Wireshark Basics

Wireshark is a graphical packet analyzer. We recommend using Wireshark for manual packet inspection to aid in developing your solutions for Project 4.

Running Wireshark Inside the VM

To run Wireshark inside your VM, launch the interactive menu (accessible via the bottom-left-most icon). Navigate to the Internet tab, and select Wireshark.



- Wireshark Basics
 Running Wireshark
- Opening PCAP Files
 Navigating Wireshark
- Main Window
- Packet List
 Packet Details
- Link Layer
- Network Layer
- Transport Layer
- Application Layer
- Packet Bytes

Helpful Tips & Tricks

- Filter Toolbar
 Filter Operators
- Filter Examples
- · Following Streams

CS 4440 Wiki: Scapy Cheat Sheet **Table of Contents:** Scapy Basics Below is an abridged cheat sheet of Scapy fundamentals that you'll use in this course. Installing Scapy This page is by no means comprehensive-we encourage you to bookmark and familiarize yourself with one • PCAP of the many in-depth Scapy tutorials on the web. Some great examples are: Show Layers Scapy Library Has Laver Scapy Usage Pavload Link Layer Source Scapy Basics Destination Network Layer In case Scapy is not installed on your VM, just run \$ pip3 install scapy in your VM's terminal. Once installed, Version here's how you can import Scapy in Python: Source from scapy,all import * Destination Transport Layer rdpcap('(pcap)') : Retrieve packets from a packet capture file. Source Port Destination Port >>> packets = rdpcap('crack0.pcap') >>> packet = packets[0] Flags Application Layer packet.show() : Show available protocols and fields for a packet. Data >>> packet.show() ###[Ethernet]### dst = 08:00:27:6e:cf:4a = 0a:00:27:00:00:00 src ###[IP]### version = 4 ihl = 5 ###[TCP]### sport = 54017

SCHOOL OF COMPUTING UNIVERSITY OF UTAH

dport = ftp



Interested in fuzzing?

Spring 2025: CS 5963/6963: Applied Software Security Testing

- Everything you'd ever want to know about fuzzing for finding security bugs!
- Course project: team up to fuzz **a real program** (of your choice), and find and report its bugs!
- https://cs.utah.edu/~snagy/courses/cs5963/

CS 5963/6963: Applied Software Security Testing

This special topics course will dive into today's state-of-the-art techniques for uncovering hidden security vulnerabilities in software. Projects will provide hands-on experience with real-world security tools like AFL++ and AddressSanitizer, culminating in a final project where **you'll team up to hunt down, analyze, and report security bugs in a real application or system of your choice**.

This class is open to graduate students and upper-level undergraduates. It is recommended you have a solid grasp over topics like software security, systems programming, and C/C++.

Professor





Questions?





Last time on CS 4440...

Attacks on Security Properties Denial of Service Attacks



Basic Security Properties

- Confidentiality: ???
- Authenticity: ???
- Integrity: ???
- Access Control: ???
- Availability: ???

Basic Security Properties

- **Confidentiality:** Concealment of information or resources
 - Attacks: intercept credentials, info
- Authenticity: Identification and assurance of info origin
 - Attacks: SMTP header spoofing
- Integrity: Preventing improper and unauthorized changes
 - Attacks: tampering HTML over HTTP
- Access Control: Enforce who is allowed access to what
 - Attacks: web app code injection
- **Availability:** Ability to use desired information or resource
 - Attacks: denial of service

DoS: Denial of Service

• Goal: ???



DoS: Denial of Service

- **Goal:** make a service unusable, usually by overloading the server or network
- How?



DoS: Denial of Service

- **Goal:** make a service unusable, usually by overloading the server or network
- How?

ERSITY OF UTAF

- Trigger the host to **crash**
 - Application-based DoS
 - Memory corruption
- Consume host's resources
 - TCP SYN floods
 - ICMP ECHO (ping) floods
- Consume host's **bandwidth**
 - UDP floods
 - ICMP floods











Stefan Nagy





Stefan Nagy

















- Reflection:
 - ???





- Reflection:
 - IP spoofing to redirect response to a victim
- Amplification:
 - ???





Reflection:

IP spoofing to redirect response to a victim

Amplification:

- Technique that increases the amount of traffic or packet size that the victim sees versus what the attacker originally sent
- How do these make detection harder?
 ???



Reflection:

IP spoofing to redirect response to a victim

Amplification:

 Technique that increases the amount of traffic or packet size that the victim sees versus what the attacker originally sent

How do these make detection harder?

- Source remains obfuscated
- Source constantly changes





DDoS or legitimate traffic?



The TCP Three-way Handshake

Recall: TCP is a connection-oriented protocol

- Initiate with three-way "handshake": SYN, SYN-ACK, ACK
- Server waits until client responds with ACK





SYN Flooding Attack











SYN Flooding Attack

- Attack: spam SYN packets to server, with spoofed origin address
 - Server's resources completely reserved—now can't serve legitimate clients



ICMP: Internet Control Message Protocol

- ICMP: pings to determine whether a system is connected to the Internet
 - Analogous to "Hello, are you still there?"



ICMP Smurf Attacks











ICMP Smurf Attacks

- Attack: takes advantage of broadcast-enabled hosts to amplify attack
- Attacker spams spoofed-source ICMP requests, reflected to victim's IP



ICMP Ping of Death Attack

Internet Protocol: IPV4 packets should be less than 65,536 bytes

- Packets can be sent in **fragments** and **reassembled** by receiver
- Attack: ???

IP	ICMP	ICMP
Header	Header	Data
20 bytes	8 bytes	65,508 bytes



ICMP Ping of Death Attack

- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in **fragments** and **reassembled** by receiver
- Attack: send packet in fragments that reassemble to 64K+ bytes
 - Many historical computer systems could not handle larger packets
- Result: crash by buffer overflow
 - Can't serve clients until restart!




ARP: Address Resolution Protocol

• **ARP:** query to **resolve the MAC address** given a desired host IP

• How we know which **physical** address to transmit data to from its logical address





ARP Flooding Attack











ARP Flooding Attack

• Attack: same idea as ICMP Smurfing; spoof source to victim and spam away!

Victim gets overwhelmed by ARP replies and bandwith crashes



Physical Layer DoS

Russian Spy Submarines Are Tampering with Undersea Cables That Make the Internet Work. Should We Be Worried?

A massive cable attack is probably an over-hyped scenario, at least for a country with as many redundant cables as the United States pitted against a limited number of Russian special-operations submarines.



CNN Exclusive: FBI investigation determined Chinese-made Huawei equipment could disrupt US nuclear arsenal communications

Thwarting DoS/DDoS Attacks

How?





Thwarting DoS/DDoS Attacks

Limit connection rate

Reduce to N total requests

Detect anomalous activity

- IP geo-filtering
- Packet similarity detection

Avoid holding connection state

Don't wait on "half-open" connections

Don't be part of the problem!

- Disable potential amplifiers
- Prevent botnet infection



Questions?





This time on CS 4440...

Authentication Multiple Authentication Factors One-time PINs Secure Password Storage



What is it?



What is it?

- That password you re-use for every website
- An ever-changing set of rules to frustrate you
- The most annoying thing about attending UofU



• Goal: ???





- Goal: establish trust in the identity of another communicating party
- Problem: ???





- Goal: establish trust in the identity of another communicating party
- Problem: cannot directly interact with them to verify their identity
 - Must be performed remotely
- Challenge: how can someone prove they are who they say they are?





The Three Factors of Authentication

Something you ???

Something you ???

Something you ???



The Three Factors of Authentication

Something you have

- Smartphone
- Laptop
- Email account

Something you are

- Your fingerprint
- Your DNA

SCHOOL OF COMPUTING

UNIVERSITY OF UTAH

Your iris, retina

Something you know

- Account password, banking PIN number
- Nuclear strike challenge-response code







Stefan Nagy

Single- vs. Multi-factor Authentication

- **N-factor authentication:** how many factors are used to authenticate
 - Password-only login is a single-factor authentication
- What are the **trade-offs**?
 - ???



Single- vs. Multi-factor Authentication

- **N-factor authentication:** how many factors are used to authenticate
 - Password-only login is a single-factor authentication
- What are the trade-offs?
 - Fewer factors = worse security
 - Compromise of one factor is total authentication violation
 - More factors = increased security
 - To fully violate authentication, attacker must compromise all
 - Trade-off: more annoying for user
 - Who cares? Security >> UX



SEND VERIFICATION EMAIL

Questions?





One-time PINs



Proof of Possession

How can you prove—remotely—that you possess something?

Proof of possession: make the user perform some **object-specific action** that requires their **physical interaction**



One-time PINs

• One-time PINs / Passwords:

 Password valid for only one login session or transaction

Delivering One-time PINs:

???



One-time PINs

• One-time PINs / Passwords:

 Password valid for only one login session or transaction

Delivering One-time PINs:

- SMS
 - Phone call
 - Text message
- Hardware
 - Yubico YubiKey
 - RSA SecureID
- Application
 - DUO Mobile
 - Google authenticator





9**10:00**

• Idea: call an API (e.g., math.random), send random to user, user re-enters it



random — Generate pseudorandom numbers

Source code: Lib/random.py

This module implements pseudo-random number generators for various distributions.

For integers, there is uniform selection from a range. For sequences, there is uniform selection of a random element, a function to generate a random permutation of a list inplace, and a function for random sampling without replacement.



- Idea: call an API (e.g., math.random), send random to user, user re-enters it
- Authentication **offline**? **No!**
 - User needs internet to receive the OTP code
 - Without a connection, they can't authenticate
- Demonstrably secure? No!
 - Most "random" APIs have small/predictable seeds
 - Also vulnerable to man-in-the-middle attacks

random — Generate pseudorandom numbers

Source code: Lib/random.py

Warning: The pseudo-random generators of this module should not be used for security purposes. For security or cryptographic uses, see the secrets module.

For integers, there is uniform selection from a range. For sequences, there is uniform selection of a random element, a function to generate a random permutation of a list inplace, and a function for random sampling without replacement.



Attack: SIM Swap

SIM: Subscriber Identity Module

- A small card inserted into your phone
- Connects you to your carrier's network





Attack: SIM Swap

SIM: Subscriber Identity Module

- A small card inserted into your phone
- Connects you to your carrier's network

Social engineering attack:

- Learn key info about victim. E.g.:
 - Mothers' maiden name
 - Childhood street address
- Trick carrier to issue new SIM card
 - "I'm Jeff Bezos, my phone broke!"
 - Attacker "appears to be" victim





Attack: SIM Swap

SIM: Subscriber Identity Module

- A small card inserted into your phone
- Connects you to your carrier's network

Social engineering attack:

- Learn key info about victim. E.g.:
 - Mothers' maiden name
 - Childhood street address
- Trick carrier to issue new SIM card
 - "I'm Jeff Bezos, my phone broke!"
 - Attacker "appears to be" victim
- Result: attacker is man-in-the-middle
 - Receives any OTPs transmitted by SMS!



Hackers steal thousands of dollars through victims' cell phones using SIM swap fraud

Hackers Hit Twitter C.E.O. Jack Dorsey in a 'SIM Swap.' You're at Risk, Too.

- Better idea: independently generate OTP codes based on a moving factor
 - E.g., intervals of time, unique session count, etc.





- Better idea: independently generate OTP codes based on a moving factor
 - E.g., intervals of time, unique session count, etc.





- Better idea: independently generate OTP codes based on a moving factor
 - E.g., intervals of **time**, unique session **count**, etc.

Common OTP protocols:

- HMAC-based OTP (HOTP)
 - Use session count as factor
- Time-based OTP (TOTP)
 - Use time interval as factor



- Better idea: independently generate OTP codes based on a moving factor
 - E.g., intervals of **time**, unique session **count**, etc.

Common OTP protocols:

- HMAC-based OTP (HOTP)
 - Use session count as factor
- Time-based OTP (TOTP)
 - Use time interval as factor
- Problem: desynchronization
 - E.g., user hits "login" one too many times

- Better idea: independently generate OTP codes based on a moving factor
 - E.g., intervals of **time**, unique session **count**, etc.

Common OTP protocols:

- HMAC-based OTP (HOTP)
 - Use session count as factor
- Time-based OTP (TOTP)
 - Use time interval as factor
- Problem: desynchronization
 - E.g., user hits "login" one too many times
 - Solution: make a few OTPs; user matches once



Pre-generated OTPs

Questions?





Biometrics



Biometrics

Provides proof of ???







Biometrics

Provides proof of physical identity






Biometrics

- Provides proof of physical identity
- Something unique to you (hopefully)
 - Fingerprint, iris, retina, DNA
- Security = unlikely match probability
 - Fingerprint match chance: ???
 - Iris pattern match chance: ???





Biometrics

- Provides proof of physical identity
- Something unique to you (hopefully)
 - Fingerprint, iris, retina, DNA
- Security = unlikely match probability
 - Fingerprint match chance: 1 in 64 * 10¹³
 - Iris pattern match chance: 1 in 10⁷⁸
- Trade-offs?
 - ???





Biometrics

- Provides proof of physical identity
- Something unique to you (hopefully)
 - Fingerprint, iris, retina, DNA
- Security = unlikely match probability
 - Fingerprint match chance: 1 in 64 * 10¹³
 - Iris pattern match chance: 1 in 10⁷⁸
- Trade-offs?
 - Engineering effort, storage size, privacy concerns







Biometric Challenges





Biometric Challenges

Replay attacks

Spoofs an enrolled user

Poisoning attacks

- Alter enrollment template
- Alter one user's enrollment

Noisy sensors

 Gives attackers "leeway" in crafting adversarial inputs

Change / loss of biometric

- Change: cataracts surgery
- Loss: losing your finger



After an initial analysis, the Indian and American scientists used three iris sensors and two commercial iris biometric matchers to check if the new irises passed biometric authentication. They found that the iris sensors' success rate dropped to 75% after surgery. The biometric matchers did better authenticating 93% of the irises.



Crane horror *Reg* reader uses his severed finger to unlock Samsung Galaxy phone

On the other hand he was fine



Questions?







Something that you ???

	uNID: (e.g	. u8675309)	_
			Forgot your uNID?
	Password	d:	
			Forgot your password?
	LOGIN		
auti our l	on: Before ent	ering your uNID or pass ting you to a University	sword, verify that the address in the URL bar o of Utah web site.
mpo eque ne se Some	rtant security sted and to oth ervices you are browsers, incl	information: This login er protected University using and exit your bro uding Google Chrome,	uses cookies to provide access to the site you of Utah websites. For your security, log out of wser when you have finished your session. retain cookie information by default even after
ou c	lose your brow	ser. Review your brows	er's support documentation to set your browse

- Something that you know
 - Something that you forget?
- A secret string of data that confirms a user's identity

L	ogin	
	uNID: (e.g. u8675309)	Forgot your uNID?
	Password:	Forgot your password?
	LOGIN	
auti	ion: Before entering your uNID or pa	assword, verify that the address in the URL bar of ity of Utah web site.
mpo eque ne se Some ou c	rtant security information: This lo asted and to other protected Univers ervices you are using and exit your l e browsers, including Google Chrom close your browser. Review your bro	in uses cookies to provide access to the site you ity of Utah websites. For your security, log out of prowser when you have finished your session. e, retain cookie information by default even after wser's support documentation to set your browser



- Something that you know
 - Something that you forget?
- A secret string of data that confirms a user's identity
 - Letters (ABCDEFGH)
 - Digits (0123456789)
 - Other symbols (\$#%-_!)

Login	
uNID: (e.g. u8675309)	Forgot your uNID?
Password:	Forgot your password?
LOGIN	
Caution: Before entering your uNID or pass rour browser is directing you to a University	word, verify that the address in the URL bar of of Utah web site.
mportant security information: This login equested and to other protected University of he services you are using and exit your brow Some browsers, including Google Chrome, r ou close your browser. Review your browse o clear cookies automatically upon exit. Inst	uses cookies to provide access to the site you of Utah websites. For your security, log out of vser when you have finished your session. etain cookie information by default even after n's support documentation to set your browser ructions for Google Chrome.

- Something that you know
 - Something that you forget?
- A secret string of data that confirms a user's identity
 - Letters (ABCDEFGH)
 - Digits (0123456789)
 - Other symbols (\$#%-_!)
- Cryptographically secure?

Lo	ogin
	uNID: (e.g. u8675309) Forgot your uNID?
	Password: <u>Forgot your password?</u>
	LOGIN
Cautio	on: Before entering your uNID or password, verify that the address in the URL bar rowser is directing you to a University of Utah web site.
mpor eque he se Some rou cl o clea	tant security information: This login uses cookies to provide access to the site you sted and to other protected University of Utah websites. For your security, log out o rvices you are using and exit your browser when you have finished your session. browsers, including Google Chrome, retain cookie information by default even afte ose your browser. Review your browser's support documentation to set your brows ar cookies automatically upon exit. Instructions for Google Chrome.

- Something that you know
 - Something that you forget?
- A secret string of data that confirms a user's identity
 - Letters (ABCDEFGH)
 - Digits (0123456789)
 - Other symbols (\$#%-_!)
- Cryptographically secure?
 - Not at all!

Lc	ogin	
	uNID: (e.g. u8675309)	argot your uNID?
	Password:	orgot your password?
	LOGIN	
Caution by the second s	on: Before entering your uNID or password, ver rowser is directing you to a University of Utah w	ify that the address in the URL bar of reb site.
Impor reque the se Some you cl to clea	rtant security information: This login uses coo sted and to other protected University of Utah w rvices you are using and exit your browser whe browsers, including Google Chrome, retain coo ose your browser. Review your browser's support ar cookies automatically upon exit. Instructions f	kies to provide access to the site you ebsites. For your security, log out of n you have finished your session. kie information by default even after rt documentation to set your browse or Google Chrome.

Cryptographically Secure = ???



- Cryptographically Secure = unbiased output, cannot be predicted
 - E.g., a cryptographically-secure pseudo-random number generator





Are most passwords biased or predictable?

Analysis of Sony and Gawker breached passwords:





Are most passwords biased or predictable?

Patterns across **all** passwords

Analysis of Sony and Gawker breached passwords:





Direct

Numbers

Symbols

Reverse

55%

Are most passwords biased or predictable?

Patterns across **all** passwords

Analysis of Sony and Gawker breached passwords:





Direct

Numbers

Symbols

Reverse

64%

Are most passwords biased or predictable?

SCHOOL OF COMPUTING

UNIVERSITY OF UTAH

Analysis of Sony and Gawker breached passwords:



Are most passwords biased or predictable?

Analysis of Sony and Gawker breached passwords:



Passwords derived from **numbers**

Stefan Nagy

Are most passwords biased or predictable?

Analysis of Sony and Gawker breached passwords:



SCHOOL OF COMPUTING

UNIVERSITY OF UTAH

Patterns across **all** passwords

Passwords derived from keyboard patterns



Are most passwords biased or predictable?

Analysis of Sony and Gawker breached passwords:



Patterns across **all** passwords

Passwords derived from **pop culture references**



Attack: Guessing Passwords

- Known **default** passwords:
 - Device manufacturers don't care
 - E.g., password, 12345, etc.
 - How Mirai Botnet spread itself

Username	Password
666666	666666
888888	888888
admin	(none)
admin	1111
admin	111111
admin	1234
admin	12345
admin	123456
admin	54321
admin	7ujMko0admin
admin	admin

Attack: Guessing Passwords

- Known **default** passwords:
 - Device manufacturers don't care
 - E.g., password, 12345, etc.
 - How Mirai Botnet spread itself

Social engineering attacks:

- Trick victim to revealing key info
 - E.g., date of birth, nickname pet's name, favorite team
- Try to guess their password
 - E.g., GoChiefs94, Chiefs1994

Username	Password
666666	666666
888888	888888
admin	(none)
admin	1111
admin	1111111
admin	1234
admin	12345
admin	123456
admin	54321
admin	7ujMko0admin
admin	admin

1 in 3 U.S. Pet Parents Have Used Their Pet's Name as Their Password





Server-side Password Storage



Why is storing passwords in **plaintext** problematic?



Server-side Password Storage



If database **breached**, attacker has **all passwords**!



Server-side Password Storage



- Assumption: attacker has full access to our database of hashed passwords
 - E.g., SQL injection, other web app attacks



- Assumption: attacker has full access to our database of hashed passwords
 - E.g., SQL injection, other web app attacks
- What if a weak hash function is used?
 ???





- Assumption: attacker has full access to our database of hashed passwords
 - E.g., SQL injection, other web app attacks
- What if a **weak** hash function is used?
 - **Pre-image attacks:** find the original string
 - **Collision attacks:** find a different string that produces same hash as password



function



image

- Assumption: attacker has full access to our database of hashed passwords
 - E.g., SQL injection, other web app attacks
- What if a weak hash function is used?
 - Pre-image attacks: find the original string
 - Collision attacks: find a different string that produces same hash as password

What if a fast hash function is used?







- Assumption: attacker has full access to our database of hashed passwords
 - E.g., SQL injection, other web app attacks
- What if a weak hash function is used?
 - Pre-image attacks: find the original string
 - Collision attacks: find a different string that produces same hash as password
- What if a fast hash function is used?
 - Attacker can quickly pre-generate hashes for all possible password possibilities



103

Attack: Rainbow Tables

- Similar to a lookup table—attacker can trade-off disk space vs. CPU time
 - Attacker wants something that uses less time, less storage than a brute-force attack



Attack: Rainbow Tables

- Similar to a lookup table—attacker can trade-off disk space vs. CPU time
 - Attacker wants something that uses less time, less storage than a brute-force attack
- Idea: iteratively hash and reduce to form a connected "chain" of hashes
 - **Simple reduction function:** truncate to just the first **10** characters of every hash



Attack: Rainbow Tables

• To find a **password** from its hash, **perform reductions** and check for a match

• For efficiency, only the starting and ending links are stored per each chain



Better Password Generation

• Why is reusing the same password bad practice?







Better Password Generation

Why is reusing the same password bad practice?

If a breached server stores it in **plaintext**, your credentials are now stolen!






- Slower hash functions
 - ???



Slower hash functions

- Makes rainbow table generation more computationally expensive for attackers!
- E.g., **Bcrypt, Scrypt**—perform multiple rounds of hashing (**much slower**)



Slower hash functions

- Makes rainbow table generation more computationally expensive for attackers!
- E.g., **Bcrypt, Scrypt**—perform multiple rounds of hashing (**much slower**)

Salted passwords:

- Add **extra data** when generating hash
- **Goal:** same input = different output





Slower hash functions

- Makes rainbow table generation more computationally expensive for attackers!
- E.g., Bcrypt, Scrypt—perform multiple rounds of hashing (much slower)

Salted passwords:

- Add **extra data** when generating hash
- **Goal:** same input = different output
- Salting considerations:
 - Salt should not be short
 - Should be **unique** per user





Slower hash functions

- Makes rainbow table generation more computationally expensive for attackers!
- E.g., **Bcrypt, Scrypt**—perform multiple rounds of hashing (**much slower**)

Salted passwords:

- Add **extra data** when generating hash
- **Goal:** same input = different output
- Salting considerations:
 - Salt should not be short
 - Should be **unique** per user

Better: salting + slow hashing!



Attack: Password Cracking

- Assume attacker knows hash function and wants to **find a single password**
 - Rapidly becoming more doable with advances in hardware!





Attack: Client-side Password Theft

How?



Attack: Client-side Password Theft

How?

Keyloggers, unencrypted transit, phishing, angry ex-partner







Forgetting and Recovering Passwords

- Security questions:
 - What's your childhood pet?
- Password recovery email
 - Click here to reset your password!
- Send in plaintext to email
 - Your password is "in\$3cur3"





Forgetting and Recovering Passwords

- Security questions:
 - What's your childhood pet?
- Password recovery email
 - Click here to reset your password!
- Send in plaintext to email
 - Your password is "in\$3cur3"

Bad security! Attacker might have control of the victim's **email**!

Forgetting and Recovering Passwords

- Security questions:
 - What's your childhood pet?
- Password recovery email
 - Click here to reset your password!
- Send in plaintext to email
 - Your password is "in\$3cur3"
- Other approaches:
 - Phone call
 - Session-specific PIN

Bad security! Attacker might have control of the victim's **email**!



Questions?





Next time on CS 4440...

Tor: The Onion Router Project 4 Tips



Stefan Nagy