Week 10: Lecture A Application-layer Network Attacks

Tuesday, October 29, 2024



Announcements

Project 3: WebSec released

Deadline: Thursday, November 7th by 11:59PM (next week)

Project 3: Web Security

Deadline: Thursday, November 7 by 11:59PM.

Before you start, review the course syllabus for the Lateness, Collaboration, and Ethical Use policies.

You may optionally work alone, or in teams of **at most two** and submit **one project per team**. If you have difficulties forming a team, post on **Piazza's Search for Teammates** forum. Note that the final exam will cover project material, so you and your partner should collaborate on each part.

The code and other answers your group submits must be entirely your own work, and you are bound by the University's Student Code. You may consult with other students about the conceptualization of the project and the meaning of the questions, but you may not look at any part of someone else's solution or collaborate with anyone outside your group. You may consult published references, provided that you appropriately cite them (e.g., in your code comments). **Don't risk your grade and degree by cheating!**

Complete your work in the **CS 4440 VM**—we will use this same environment for grading. You may not use any **external dependencies**. Use only default Python 3 libraries and/or modules we provide you.

Project 3 progress

Working on Part 1

Finished Part 1, working on Part 2

Finished Part 2, working on Part 3

Haven't started :(

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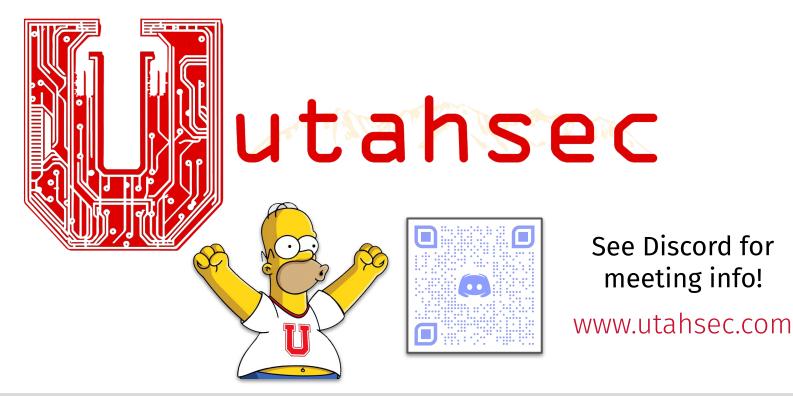
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Announcements



Questions?





Last time on CS 4440...

Introduction to Networking The Physical, Link, Network, Transport, and Application Layers



What is the Internet?

What is it?

- How you trash-talk players in COD game lobbies
- How Wall Street trades shares faster than you
- How the CS 4440 website is distributed to you

MUTEP	LAYERS	
VOICE CHAT	ON CON	
GAME VOICE CHANNEL	ALL LOBBY	
VOICE CHAT DEVICE	ALL LOBBY	
VOICE CHAT RECORDING MODE	PARTY ONLY	
VOICE CHAT (PUSH TO TALK)	FRIENDS ONLY	
■ OPEN MIC RECORDING THRESHOLD	70	



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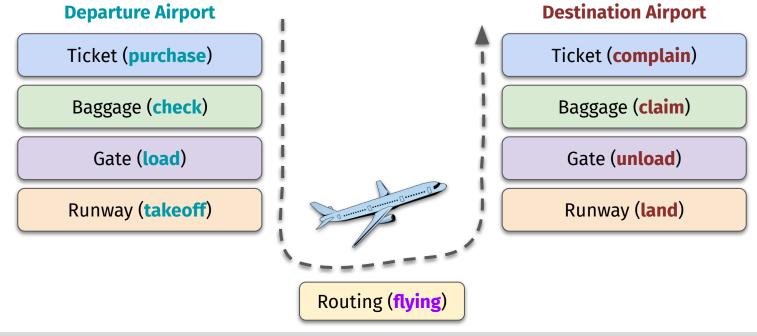
CS 4440: Introduction to Computer Security

This course teaches the security mindset and introduces the principles and practices of computer security as applied to software, host systems, and networks. It covers the foundations of building, using, and managing secure systems. Topics include standard cryptographic functions and protocols, threats and defenses for realworld systems, incident response, and computer forensics.

This class is open to undergraduates. It is recommended that you have a solid grasp over topics like software engineering, computer organization, basic networking, SQL, scripting languages, and C/C++.

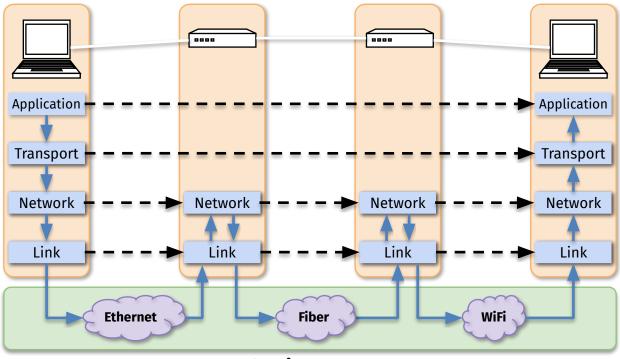
What really is the internet?

• **A group of layers**—each implementing a **service**





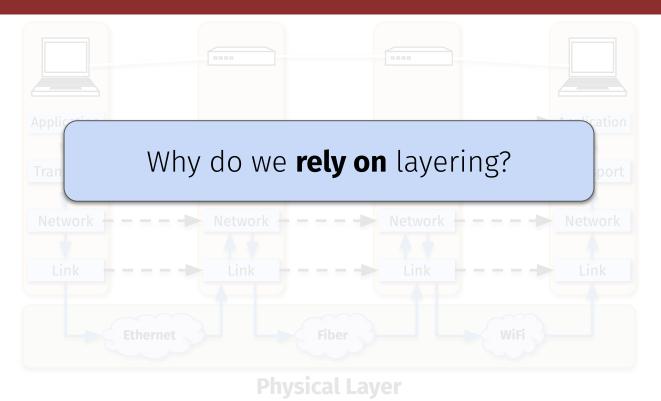
The 5-layer Internet



Physical Layer

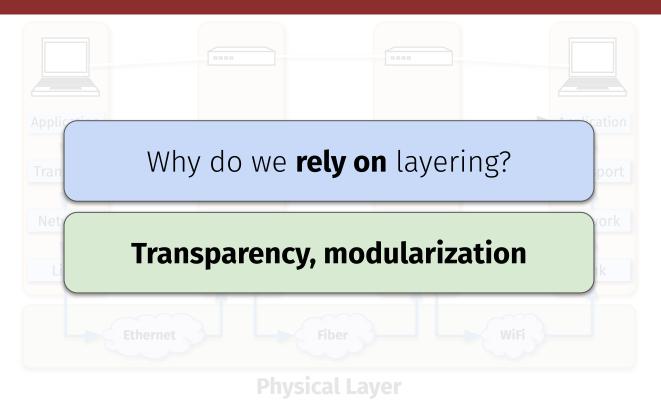


The 5-layer Internet





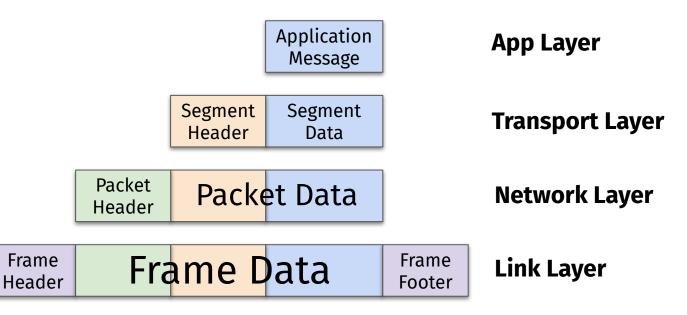
The 5-layer Internet





Internet Packet Encapsulation

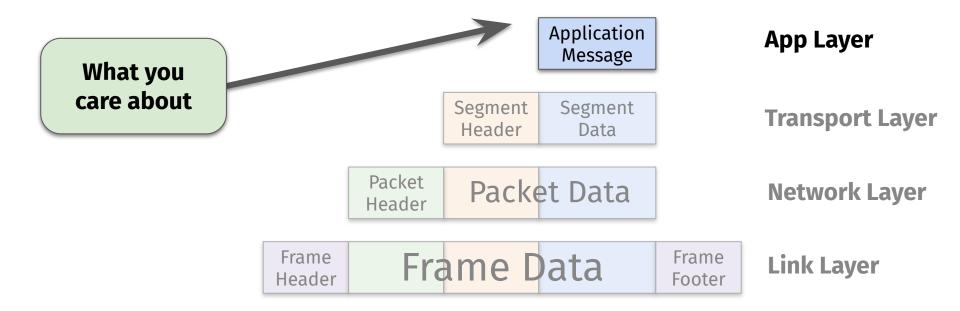
How packets are generated and sent





Internet Packet Encapsulation

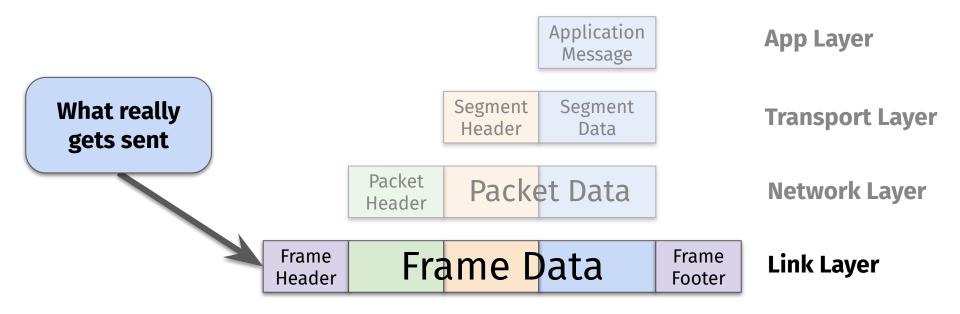
How packets are generated and sent





Internet Packet Encapsulation

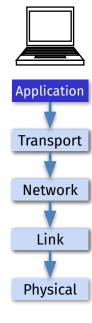
How packets are generated and sent



The Application Layer

What is it?

???



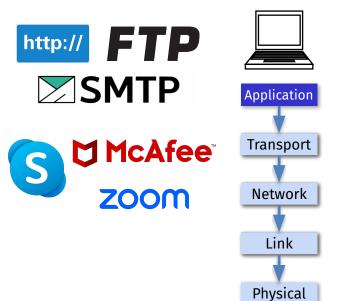


The Application Layer

What is it?

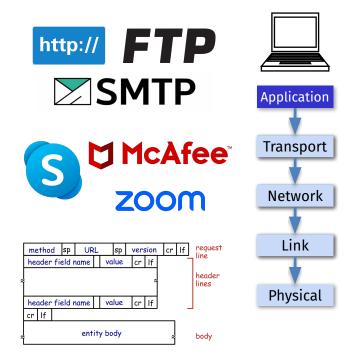
- The top-most layer in the 5-layer network model
- Where applications send and receive messages
- What does it **define**? **Application protocols**

???



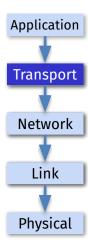
The Application Layer

- The top-most layer in the 5-layer network model
- Where applications send and receive messages
- What does it **define**? Application protocols
 - Message types
 - What is the purpose of this message?
 - Message syntax
 - How should this message be structured?
 - Message semantics
 - What does each message field really mean?
 - Rules for sending/receiving messages
 - When/how should this application respond?



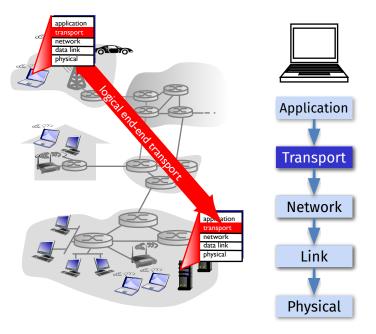
- What is it?
 - ???







- The **second layer** in the 5-layer network model
- Communication between apps on different hosts
- What are its two main **protocols**?
 - ???

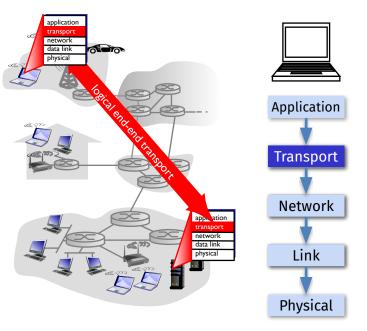




What is it?

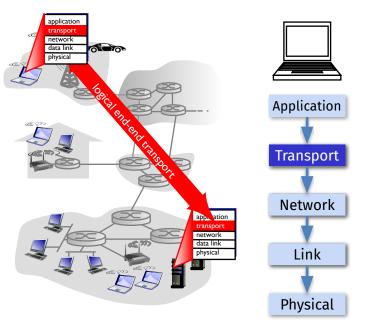
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- The **second layer** in the 5-layer network model
- Communication between apps on different hosts
- What are its two main **protocols**? **TCP**, **UDP**
 - TCP—Transmission Control Protocol
 - Characteristics: slow/complex but reliable
 - UDP—User Datagram Protocol
 - Characteristics: fast/simple but unreliable
- What are ideal use cases for TCP and UDP?



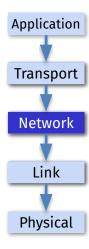


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 - **TCP**—Transmission Control Protocol
 - Characteristics: slow/complex but reliable
 - UDP—User Datagram Protocol
 - Characteristics: fast/simple but unreliable
- What are ideal use cases for TCP and UDP?
 - TCP: reliability matters (file transfer, SSH, e-mail)
 - UDP: speed matters (video calls, gaming, livestream)



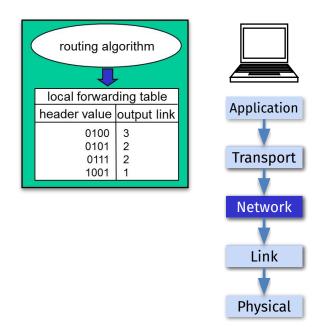
- What is it?
 - ???







- The third layer in the 5-layer network model
- Sends data from host on one network to another
- What are its two functions?
 - ???





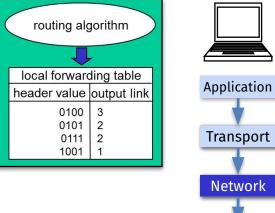
What is it?

- The **third layer** in the 5-layer network model
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What are its two functions? Routing, forwarding

- Routing: find shortest possible path to send a packet
- Forwarding: sending packets on to the next hop
- What is its addressing based on?

???



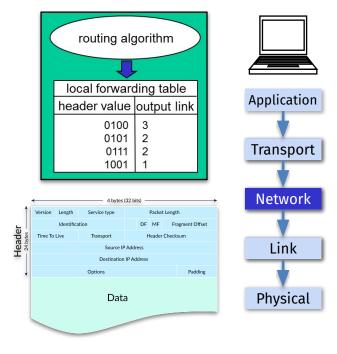


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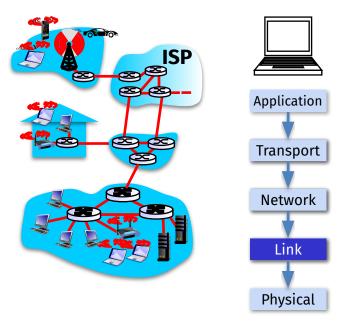
What are its two functions? Routing, forwarding

- Routing: find shortest possible path to send a packet
- Forwarding: sending packets on to the next hop
- What is its addressing based on?
 - IP addresses—a logical address
 - Network-internal IP assigned by your router
 - Public IP assigned by Internet Service Provider



What is it?

• ???

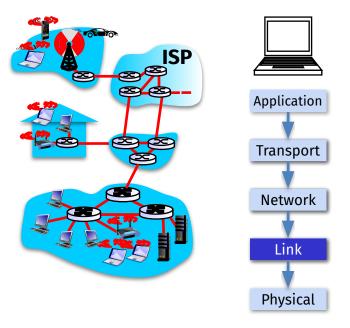




What is it?

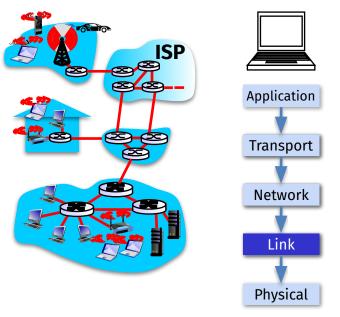
- The **fourth layer** in the 5-layer network model
- Responsible for the node-to-node delivery of data
- What are "nodes"?

???



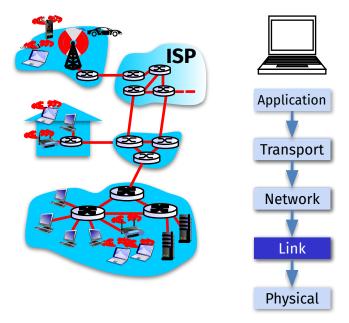


- The **fourth layer** in the 5-layer network model
- Responsible for the node-to-node delivery of data
- What are "nodes"? Hosts, switches
 - Hosts: the physical devices within a network
 - Switches: interface to all hosts on the network
- What is its addressing based on?





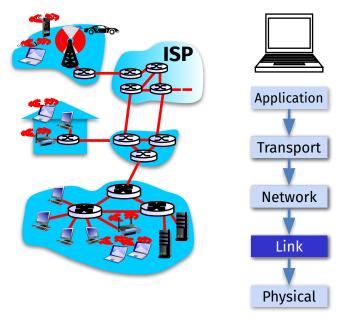
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 - MAC addresses—a physical identifier for hardware
- Do MAC addresses guarantee **authenticity**?
 - ???



What is it?

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 - MAC addresses—a physical identifier for hardware
- Do MAC addresses guarantee **authenticity**?
 - Reconfigurable via network interface
 - Attacker-spoofable

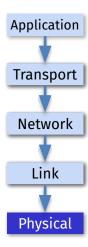


Stefan Nagy

What is it?

• ???

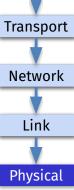






- The **last layer** in the 5-layer network model
- The physical means of sending/receiving data
- **Examples** of physical layers?
 - ???





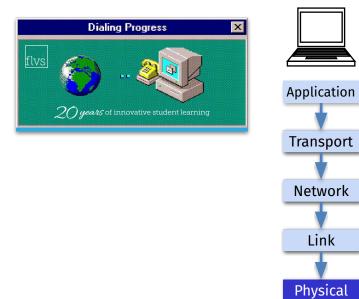


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- The **last layer** in the 5-layer network model
- The physical means of sending/receiving data
- **Examples** of physical layers?
 - Radio waves

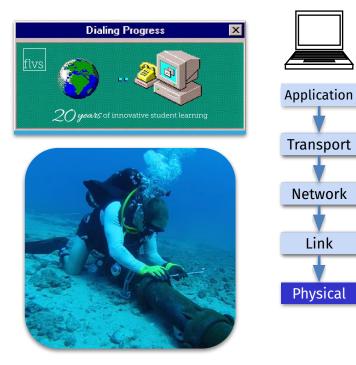
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- Telephone lines
- Fiber optic cables
- Undersea submarine cables
- Does physical layer guarantee **availability**?



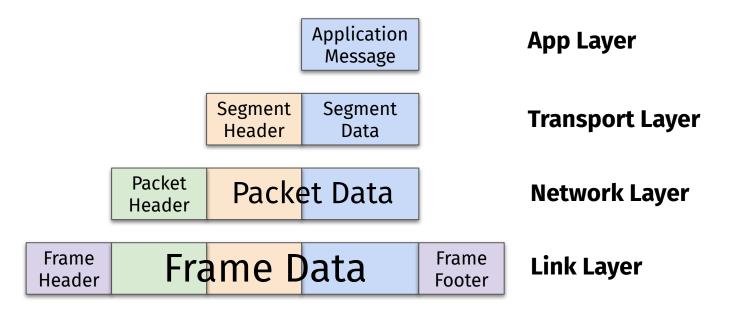


- The **last layer** in the 5-layer network model
- The physical means of sending/receiving data
- **Examples** of physical layers?
 - Radio waves
 - Telephone lines
 - Fiber optic cables
 - Undersea submarine cables
- Does physical layer guarantee availability?
 - No—tamperable by third parties!



Food for Thought

• Are any of the five network layers susceptible to **attacks**? If so, **which ones**?





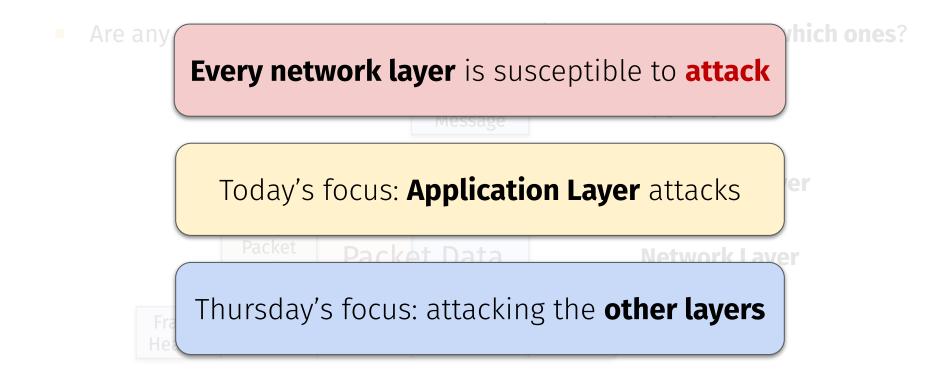
Which network layers are susceptible to attack?

Physical	
	0%
Link	
	0%
Network	
	0%
Transport	
	0%
Application	
	0%
None of the above	
	0%
All of the above	
	0%



Start the presentation to see live content. For screen share software, share the entire screen. Get help at **pollev.com/app**

Food for Thought





Questions?





This time on CS 4440...

Application Layer Attacks HTTP Content Injection SMTP Header Spoofing DNS Hijacking Network Packet Analysis



Recap: The 5-Layer Internet

Application Layer:

- Sends/receives app messages
- Transport Layer:
 - Communication between apps

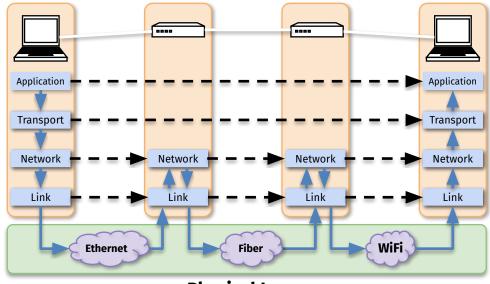
Network Layer:

- Communication between hosts
- Data Link Layer:
 - Node-to-node delivery of data

Physical Layer:

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Send/receive the physical signals



Physical Layer

Application Layer: ???

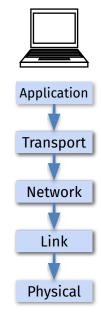
Transport Layer: ???

Network Layer: ???

Data Link Layer: ???



- **1. Host Device** (e.g., your laptop)
- **2. Other Devices** (e.g., switch, router)
- 3. Both!



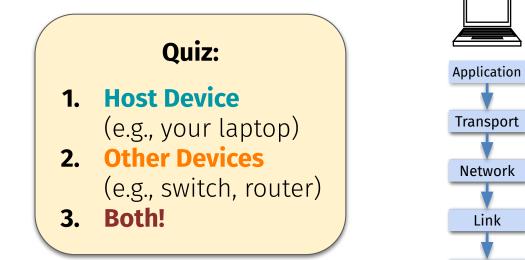


Stefan Nagy

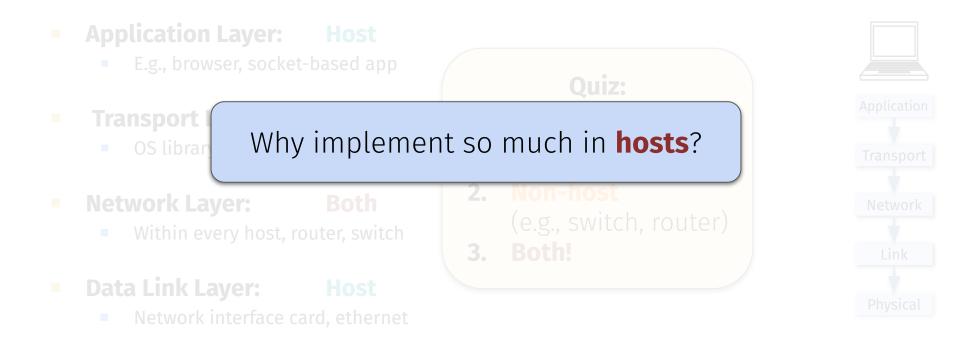
- Application Layer: Host
 - E.g., browser, socket-based app
- Transport Layer: Host
 - OS library and necessary drivers
- Network Layer: Both
 - Within every host, router, switch
- Data Link Layer: Host

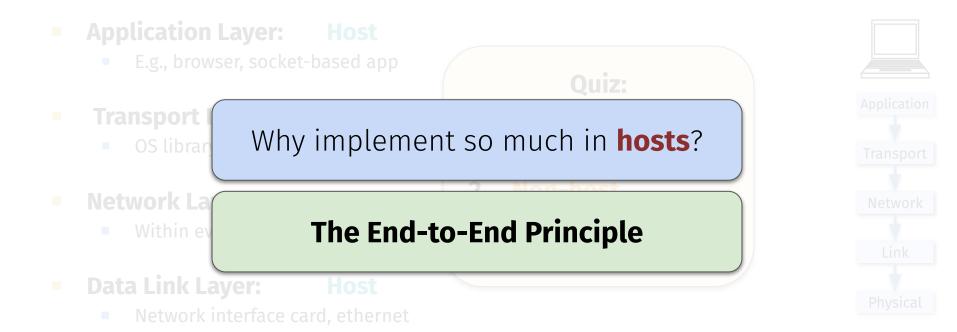
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Network interface card, ethernet



Physical





 Based on Paul Baran's 1960's work "reliability from unreliable parts"

Key idea:

- Application-specific functions ought to reside in the end hosts of a network
- Rather than in intermediary nodes
- All this assumes that the end host can implement it "completely and correctly"

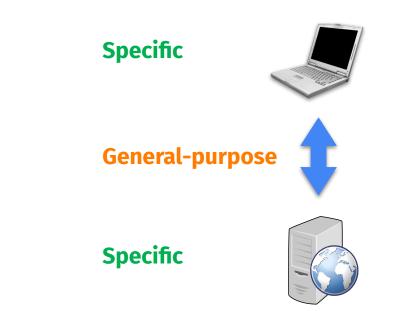




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Based on Paul Baran's 1960's work

"reliability from unreliable parts"

Key idea:

Application-

Why not move more functionality to the **intermediate nodes**?

reside in the end hosts of a network

- Rather than in intermediary nodes
- All this assumes that the end host can implement it "completely and correctly"

Specific



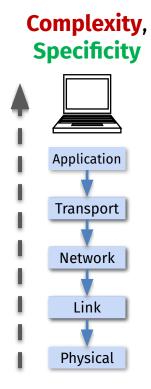


"reliability from unreliable parts" Why not move more functionality to the **intermediate nodes**? reside in the end hosts of a network More **latency**, less **flexibility**, higher-**complexity** midpoints

Implications of The E2E Principle

Nothing guaranteed by default

- No default integrity
- No default confidentiality
- No default availability
- No default **authentication**





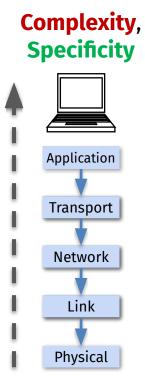
Implications of The E2E Principle

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- No default integrity
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- No default availability
- No default authentication

Solution: bolt-on more security!

- IPsec (Network layer)
 - Integrity, Confidentiality, Authentication
- DNSSec (App layer)
 - Integrity, Authenticity
- TLS (Session layer)
 - Integrity, Confidentiality, Authentication



Application Layer Network Attacks





Application Layer Attacks

- Application Layer: where network-facing apps send/receive message
 - Application-specific protocols (message semantics, structure, processing rules, etc.)
- Attacking the application layer:
 - ???



Application Layer Attacks

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 - Application-specific protocols (message semantics, structure, processing rules, etc.)

Attacking the application layer:

- Command Injection
 - SQL injection, CSRF, XSS
- Denial of Service
 - Crash a remote application
 - Prevent others from using it
- Message Tampering / Sniffing
 - Injecting data into messages
 - Capturing unencrypted data

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 - Injecting data into messages
 - Capturing unencrypted data
- Other protocol-specific attacks



Application Layer Attacks HTTP Content Injection



Recap: HyperText Markup Language (HTML)

What is HTML?

<form action="home.html">
 First Name:

 <input type="text" name="first_name">
 </br>
 Last Name:

 cinput type="text" name="last_name">
 </br>
 Email:

 cinput type="text" name="last_name">
 </br>
 cinput type="text" name="last_name">
 </br
 cinput type="text" name="last_name">
 </br
 cinput type="text" name="last_name">
 </br
 cinput type="text" name="last_name">
 c/br>
 cinput type="text" name="email">
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 cinput type="text" name="submit">
 cinput type="text" name="email">
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 cinput type="text" name="submit">
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 cinpu

First Name:
Last Name:
Email:
Submit Query



Recap: HyperText Markup Language (HTML)

What is HTML?

- Describes content and formatting of web pages
- Rendered within browser window

HTML features

- Static document description language
- Links to external pages, images by reference
- User input sent to server via forms

HTML extensions

- Additional media (e.g., PDF, videos) via **plugins**
- Embedding programs in other languages (e.g., Java) provides dynamic content that can:
 - Interacts with the user
 - Modify the browser user interface
 - Access the client computer environment

<form action<="" th=""><th>on="home.html"></th></form>	on="home.html">
First I	Name:
<input< td=""><td><pre>type="text" name="first_name"</pre></td></input<>	<pre>type="text" name="first_name"</pre>
Last Na	ame:
<input< td=""><td><pre>type="text" name="last_name"></pre></td></input<>	<pre>type="text" name="last_name"></pre>
Email:	
<input< td=""><td>type="text" name="email"></td></input<>	type="text" name="email">
<input< td=""><td>type="submit" name="Submit"></td></input<>	type="submit" name="Submit">
	(First Name:
	L (NI
	Last Name:
-	
	Email:
	I E MAIL

Submit Query

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Recap: HyperText Transfer Protocol (HTTP)

What is HTTP?

Helio Accounts Transfer & Pay	You have the dreams and the funds. You easing forme easily in all oriest can fund home Personners programs large particulars and more.	ORTINET	
8 Send Money with Zelle*	DEPOSITS & INVESTMENTS		
	5/3 Essential Checking	Available	
5/3	5/3 Preferred Checking	Available	
	5/3 Essential Checking	Available	
r ID	Maxsaver	Available	
soword	Roth Ira	Available	
Remember Me Forgot Login	Ira	Available	
Log In	CREDIT CARDS & LOANS		
	Equity Line Princip	pat Balance	
NDP 50 FM Open Support Erach Accessit profile 2022 PRN The Back, National Association All Stight Interved			



What is HTTP?

How we transmit hyper-media objects over the web!

An unencrypted, stateless protocol for transmitting information from server to client (and back).

0%

0%

Halloween Toblerone transfer protocol 🎃 🍫

None of the above

0%

0%



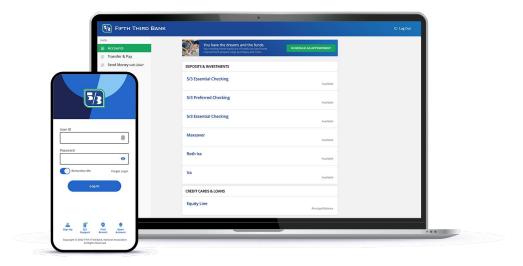
Recap: HyperText Transfer Protocol (HTTP)

What is HTTP?

Protocol for transmitting hypermedia documents (e.g., web pages)

HTTP's Characteristics:

- Widely used
- Simple
- Unencrypted





Recap: HyperText Transfer Protocol (HTTP)

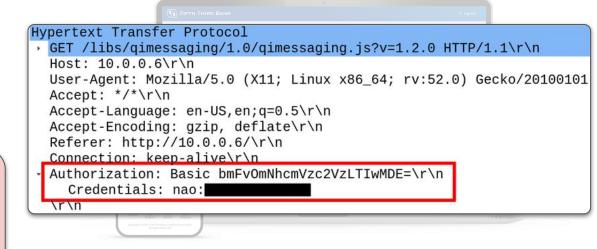
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Tampering with HTTP-transmitted HTML

- Capitalizes on HTTP's insecurity
 - Nothing is encrypted!
- Attacker intercepts requested webpage and modifies it
 - User receives modified webpage
- Attacker capabilities?
 - ???



Tampering with HTTP-transmitted HTML

- Capitalizes on HTTP's insecurity
 - Nothing is encrypted!
- Attacker intercepts requested webpage and modifies it
 - User receives modified webpage
- Attacker capabilities?
 - Inject malicious content
 - Inject malicious code





- Do you trust your **ISP**?
 - Could they tamper with data?



- Do you trust your **ISP**?
 - Could they tamper with data?
- Attack: ISP intercepts HTML pages transmitted via HTTP
 - Injects advertisements
 - They make commission



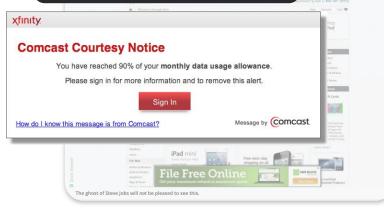
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- Attack: ISP intercepts HTML pages transmitted via HTTP
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 - Injects pop-up messages
 - E.g., monthly data usage

Comcast caught hijacking web traffic

On November 20th, 20th 2 Comcast hijacked my HTTP traffic and ra-routed it through their own servers, injecting a 'notice' on the page before completing the request. What this means is instead of my web request being routed to the website I wanted to visit, Comcast took it upon themselves to hijack my web traffic, forcing it to go through their servers instead. This poses a massive security risk for users since there's no telling what type of logging Comcast uses on their end. Why did they do all this? To force a "courtesy notice" on every webpage I visit until I logged into my Comcast account because I was within 90% of my new 300GB limit?

appeared on

In my testing I discovered that this only affects HTTP traffic and not HTTPS traffic. What this means is while your online banking may be safe, any other website you visit over HTTP may cause your privacy to be at risk. This is a prime example of why SSL encryption on websites is so important. However, it may only be a matter of time before Comcast starts executing man in the middle attacks on SSL traffic.





- Do you trust your **ISP**?
 - Could they tamper with data?
- Attack: ISP intercepts HTML pages transmitted via HTTP
 - Injects advertisements
 - They make commission
 - Injects pop-up messages
 - E.g., monthly data usage
 - Redirect search engine results
 - More commission!

Comcast caught hijacking web

traffic

wember 20th, 2012 Comcast hijacked my HTTP traffic and re-routed it through own servers, injecting a "notice" on the page before completing the request. What eans is instead of my web request being routed to the website I wanted to visit,

ISPs Accused Of Hijacking Search Terms, Redirecting Browser Results To Marketer's Websites

from the yikes dept

Fri, Aug 5th 2011 02:36pm - Mike Masnick

It's really quite stunning that ISPs and marketers haven't yet realized that hijacking users' browser functions and redirecting them for marketing purposes could get them into serious trouble. They just keep doing it. The latest involves "more than 10 ISPs" in the US who have been secretly hijacking search terms and redirecting users directly to marketers'

(Mis)Uses of W Technology

of websites. That is, if you typed "apple" into a browser search box, the service could take you directly to Apple's website, rather than to search results. In this case, the search query

never even reaches your search engine of choice, being intercepted by the ISP, via a partner called Paxfire. Christian Kreibich and Nicholas Weaver, at Berkeley, discovered this and have been tracking it for a few months. Apparently, they found 165 search terms being used in this manner, including: "apple" and "dell" and "safeway" and "bloomingdales."





- Do you trust your government?
 - Could they tamper with data?



- Do you trust your government?
 - Could they tamper with data?
- Attack: government forces ISPs to inject code into HTTP content
 - Steal HTTP-transmitted passwords
 - E.g., Facebook, GMail, Twitter

How The Tunisian Government Tried To Steal The Entire Country's Facebook Passwords

Pascal-Emmanuel Gobry Jan 24, 2011, 10:01 AM

Tunisia is in the midst of what increasingly looks like a happy, democratic revolution. People are wondering about the role social media played in that revolution. It turns out Facebook played a great role --for good and for bad.



AP

- Do you trust your government?
 - Could they tamper with data?
- Attack: government forces ISPs to inject code into HTTP content
 - Steal HTTP-transmitted **passwords**
 - E.g., Facebook, GMail, Twitter
 - **Result: persistent XSS** attack
 - Passes Same-origin Policy!

How The Tunisian Government Tried To Steal The Entire Country's Facebook Passwords

The rogue JavaScript, which was individually customized to steal passwords for each site, worked when users tried to login without availing themselves of the secure sockets layer protection designed to prevent man-in-the-middle attacks. It was found injected into Tunisian versions of Facebook, Gmail, and Yahoo! in late December, around the same time that protestors began demanding the ouster of Zine el-Abidine Ben Ali, the president who ruled the country from 1987 until his ouster 10 days ago. <

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Thwarting HTTP Injection

How do we prevent code and content injection in HTTP-transmitted data?



Thwarting HTTP Injection

How do we prevent code and content injection in HTTP-transmitted data?

Answer: completely ditch HTTP!

- As web and app developers, enforce strict **HTTPS** compliance
 - Necessary to prevent HTTPS → HTTP downgrade attacks



Questions?





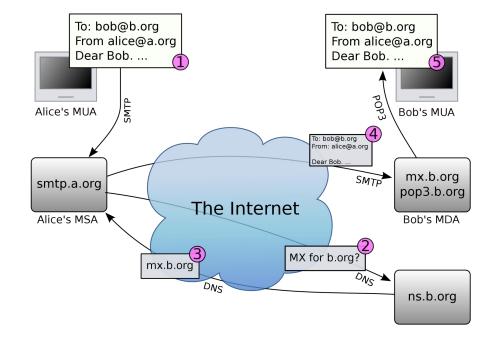
Application Layer Attacks SMTP Header Spoofing



How does sending E-mail work?

 Nigerian Prince writes me a great investment opportunity

Sends it!



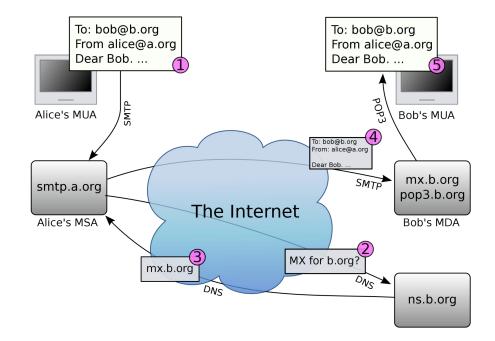


How does sending E-mail work?

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Sends it!

- Protocol SMTP (aka the Simple Mail Transfer Protocol)
- Sender's SMTP server breaks up the message into body/receiver
- Sender's SMTP server queries DNS to find receiver's server IP
- Receiver's SMTP server gets msg, then queries its POP3 server to find the correct user mailbox

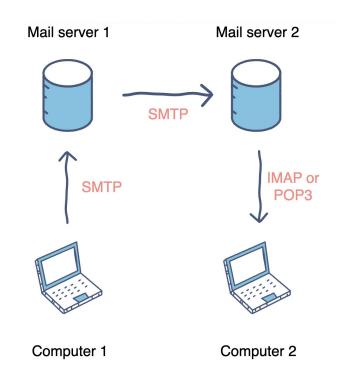


SMTP Protocol

- SMTP: Simple Mail Transfer Protocol
 - Implemented in the application layer

Characteristics:

- Text-based
- Connection-oriented
- Uses TCP ports 25/587



SMTP Protocol

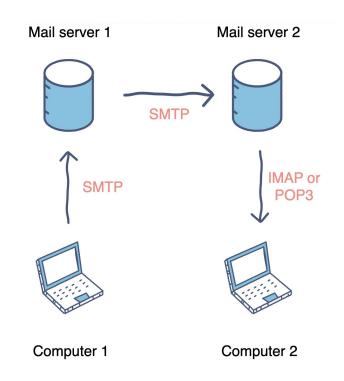
- SMTP: Simple Mail Transfer Protocol
 - Implemented in the application layer

Characteristics:

- Text-based
- Connection-oriented
- Uses TCP ports 25/587

Security guarantees:

- Message integrity—no!
- Confidentiality—no!
- Authentication—no!



Example SMTP Connection

Plain SMTP:

No encryption whatsoever

Key Protocol fields:

- HELO: setup sender's server
- MAIL FROM: sender address
- RCPT TO: recipient address
- DATA: subject, body, files

S: 220 smtp.example.com ESMTP Postfix C: HELO relay.example.org S: 250 Hello relay.example.org, I am glad to meet you C: MAIL FROM: < bob@example.org> S: 250 Ok C: RCPT TO:<alice@example.com> S: 250 Ok C: RCPT TO:<theboss@example.com> S: 250 Ok C: DATA S: 354 End data with <CR><LF>.<CR><LF> C: From: "Bob Example" <bob@example.org> C: To: "Alice Example" <alice@example.com> C: Cc: theboss@example.com C: Date: Tue, 15 January 2008 16:02:43 -0500 C: Subject: Test message C: C: Hello Alice. C: This is a test message with 5 header fields and 4 lines in the message body. C: Your friend, C: Bob C: . S: 250 Ok: queued as 12345 C: OUIT S: 221 Bye {The server closes the connection}

Example SMTP Connection

Plain SMTP:

No encryption whatsoever

Key Protocol fields:

- HELO: setup sender's server
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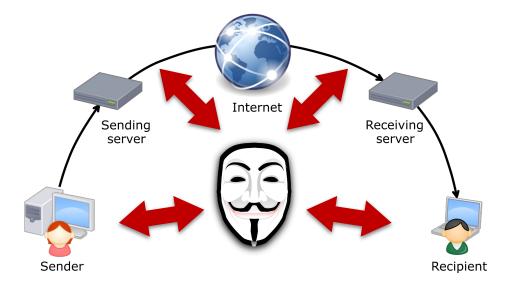
What could an attacker do?

S: 220 smtp.example.com ESMTP Postfix C: HELO relay.example.org S: 250 Hello relay.example.org, I am glad to meet you C: MAIL FROM:<bob@example.org> S: 250 Ok C: RCPT TO:<alice@example.com> S: 250 Ok C: RCPT TO:<theboss@example.com> S: 250 Ok C: DATA S: 354 End data with <CR><LF>.<CR><LF> C: From: "Bob Example" <bob@example.org> C: To: "Alice Example" <alice@example.com> C: Cc: theboss@example.com C: Date: Tue, 15 January 2008 16:02:43 -0500 C: Subject: Test message C: C: Hello Alice. C: This is a test message with 5 header fields and 4 lines in the message body. C: Your friend, C: Bob C: . S: 250 Ok: queued as 12345 C: OUIT S: 221 Bye {The server closes the connection}

SMTP Attacks

No message integrity

- Tamper with messages
- Block messages



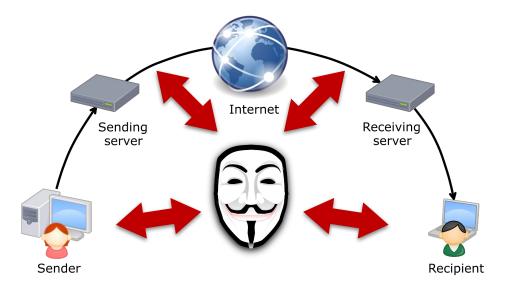
SMTP Attacks

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No confidentiality

- Find sender/recipient
- Read message contents



SMTP Attacks

No message integrity

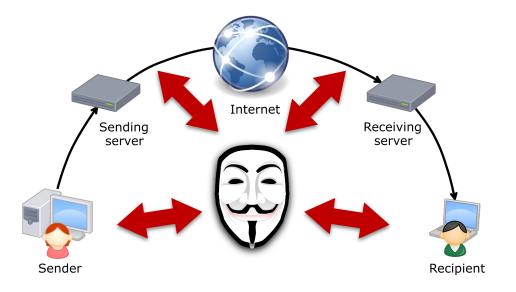
- Tamper with messages
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No confidentiality

- Find sender/recipient
- Read message contents

No authentication

Spoof sender identity



Case Study: Email Header Spoofing

• Attack: spoof email header to **mislead recipient** about sender of the email

S: 220 attacker.com SMTP Exim C: HELO attacker.com S: 250 Hello attacker.com C: MAIL FROM: <ceo@company.com> S: 250 0k C: RCPT TO: <bob@company.com> S: 250 Accepted C: DATA S: 354 Enter a message, ending with "." on a line by itself C: Subject: Download this urgently C: From: ceo@company.com C: To: bob@company.com C: C: Hi Bob, C: Please download this urgently: https://some-malicious-link.com C: Regards C: . S: 250 OK C: OUIT S: 221 attacker.com closing connection

To: robertbateman@email.com Subject: Hi There From: "Mickey Mouse" <m.mouse@disney.com> X-Priority: 3 (Normal) Importance: Normal Errors-To: m.mouse@disney.com Reply-To: m.mouse@disney.com Content-Type: text/plain



Case Study: Email Header Spoofing

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```
To: robertbateman@email.com
Subject: Hi There
From: "Mickey Mouse" <m.mouse@disney.com>
X-Priority: 3 (Normal)
Importance: Normal
Errors-To: m.mouse@disney.com
Reply-To: m.mouse@disney.com
Content-Type: text/plain
```



Thwarting Email Spoofing

Checking email bodies

- Included links
- Attached files
- Text analysis (e.g., known spam campaigns)

FROM to 👻	I MR.FRED URGENT BUSINESS ASSISTANCE consultfred61@gaza.org.in yia 💷 🖙 Mon, May 6, 9:19 AM (4 days ago)	☆	•
•	This message seems dangerous		
	Similar messages were used to steal people's personal information. Avoid clicking links, downloading attachments, or replying with personal information.		
	Looks safe		
Please			
Please	Looks safe		ł
Please			
	prove this to inbox and open the pdf file. Downloading this attachment is disabled. This email has been identified as phishing. If you want to		
	e move this to inbox and open the pdf file. Downloading this attachment is disabled. This email has been identified as phishing. If you want to download it and you trust this message, click 'Not spam' in the barner above.		
A state and state Reg of View of Reg of View of View of Reg of View of View of Reg of View of View of View of Reg of View of View of View of View of Reg of View of View of View of View of View of View of Reg of View	e move this to inbox and open the pdf file. Downloading this attachment is disabled. This email has been identified as phishing. If you want to download it and you trust this message, click 'Not spam' in the barner above.		



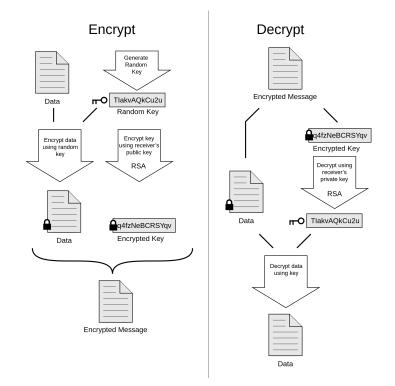
Thwarting Email Spoofing

Checking email bodies

- Included links
- Attached files
- Text analysis (e.g., known spam campaigns)

Checking email headers

- Egress server domain registration
 - Check that sender is who it says it is
- Pretty Good Privacy (PGP)
 - Sender and Receiver authentication
 - Confidentiality
 - Integrity



Questions?





Application Layer Attacks DNS Hijacking





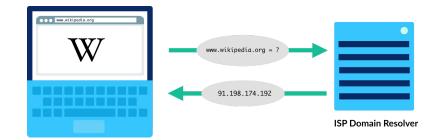
Identification on the Web

- How do we identify **people**?
 - ???



Identification on the Web

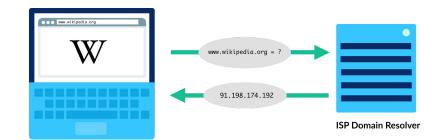
- How do we identify people?
 - Social security numbers
 - Passports, drivers licenses
 - Their unique fingerprints
- How can we identify internet hosts?
 - Network layer: location via IP addresses
 - A logical addressing system
 - 32-bit (IPV4) addressing datagrams





Identification on the Web

- How do we identify **people**?
 - Social security numbers
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 - Network layer: location via IP addresses
 - A logical addressing system
 - 32-bit (IPV4) addressing datagrams
 - What you care about: the domain name
 - E.g., www.wikipedia.org

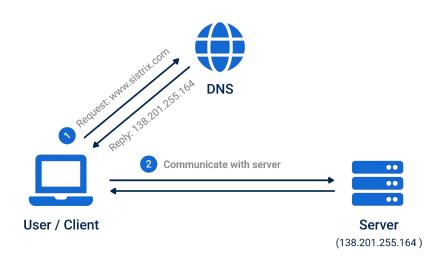






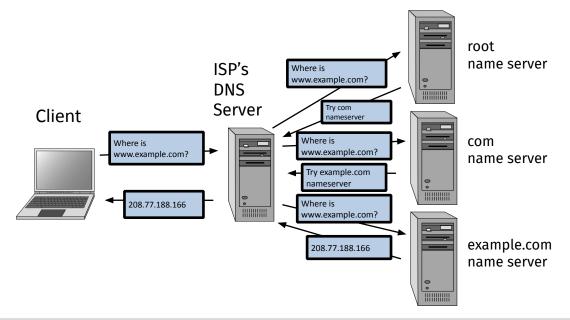
The Domain Name System

- Distributed database implemented in hierarchy of many name servers
- Application-layer protocol:
 - Hosts and domain name servers communicate to resolve domain names
 - Address-name translation
- Result: user requests domain name
 - But their host really gets its IP address
 - Convenient!



DNS Name Resolution

- **Zone:** collection of connected nodes with the same authoritative DNS server
- Resolution method when answer not in cache:

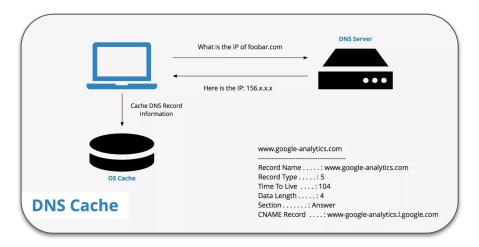


- How can we optimize DNS resolution?
 - ???



How can we optimize DNS resolution?

- Cache look-ups to amortize initial look-up, reduce system load
- Temporal locality of requests:
 - www.espn.com/page1
 - www.espn.com/page2
- Popular destinations:
 - google.com
 - Facebook.com





stefan@cs4440:~\$	time	nslookup	facebook.com
Server:	127.0	.0.53	
Address:	127.0	.0.53#53	

Non-authoritative answer: Name: facebook.com Address: 31.13.70.36 Name: facebook.com Address: 2a03:2880:f10d:83:face:b00c:0:25de

real 0m0.474s user 0m0.000s sys 0m0.015s

First Lookup (non-cached)

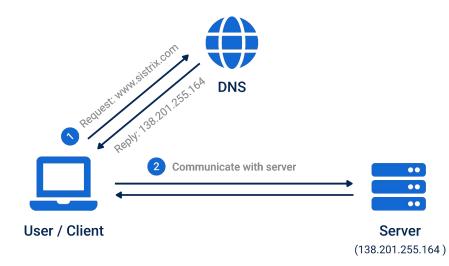
stefan@cs44	40:~\$ ti	me nslookup	facebook.com	stefan	@cs4440:~\$	5 time nslooku	p facebook.com
Server:	127	7.0.0.53		Server	:	127.0.0.53	
Address:	127	7.0.0.53#53		Addres	s:	127.0.0.53#53	
Non-authori	tative a	inswer:		Non-au	thoritativ	/e answer:	
Name: fac	ebook.co	om		Name:	facebook	<.com	
Address: 31	.13.70.3	6		Addres	s: 31.13.7	70.36	
Name: fac	ebook.co	om		Name:	facebook	<.com	
Address: 2a	03:2880:	f10d:83:fac	e:b00c:0:25de	Addres	s: 2a03:28	380:f10d:83:fa	ce:b00c:0:25de
real OmO	.474s			real	0m0.023s	5	
user 0m0	.000s			user	0m0.000s	5	
sys 0m0	.015s			sys	0m0.011s	5	

Second Lookup (cached)

First Lookup (non-cached)

Attacking DNS

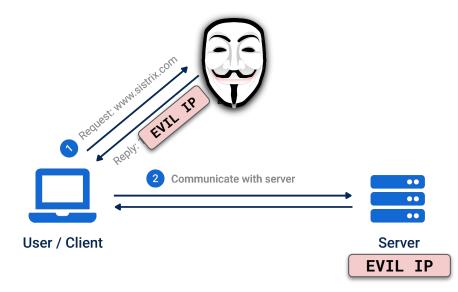
What can an attacker do if they control a DNS server?





Attacking DNS

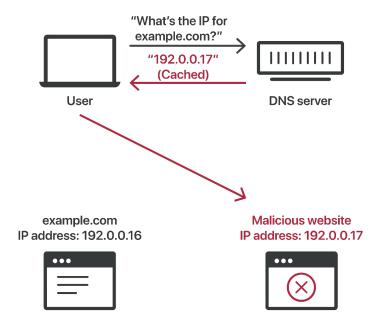
- What can an attacker do if they **control a DNS server**?
 - Control how users of that DNS server view the internet!
 - Assuming they use domain names





Case Study: DNS Cache Poisoning

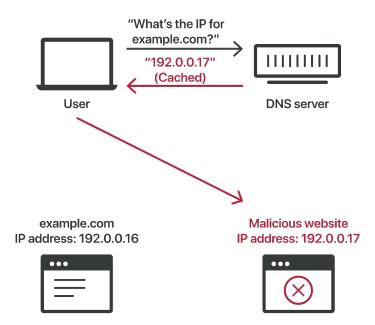
- Attack: pre-empt DNS lookup by injecting malicious cache contents
 - Exploits DNS lookup optimization!





Case Study: DNS Cache Poisoning

- Attack: pre-empt DNS lookup by injecting malicious cache contents
 - Exploits DNS lookup optimization!
- Victim performs cache lookup, instead gets malicious domain IP
 - Attacker can redirect the victim's browser to the malicious website



Case Study: DNS Cache Poisoning

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- Victim performs cache lookup, instead gets malicious domain IP
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• A massive vulnerability in 2008!

The Great DNS Vulnerability of 2008 by Dan Kaminsky

The Internet was never designed to be secure. The Internet was designed to move pictures of cats.

In 2008, Security Researcher Dan Kaminsky presented on the massively widespread and critical <u>Domain Name System (DNS)</u> <u>vulnerability</u> that allowed attackers to send users to malicious sites and hijack email at Black Hat, the information security conference. The exploit would allow attackers to impersonate any legitimate website and steal data.

This fundamental design flaw allowed for arbitrary DNS cache poisoning - affecting nearly every DNS server on the planet, including vendors and products that worked with DNS. To explain what that is here's some background on DNS:

Thwarting DNS Hijacking

Attack points:

- Local host
- Router
- ISP



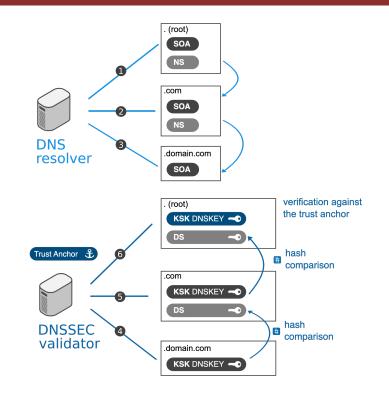
Thwarting DNS Hijacking

Attack points:

- Local host
- Router
- ISP

DNS-level authentication

- DNSSec
- Public-key crypto to "sign" DNS records
- Endpoint authentication
 - Certify that what I am seeing really is bank.com
 - Transport Layer Security (TLS)



Questions?





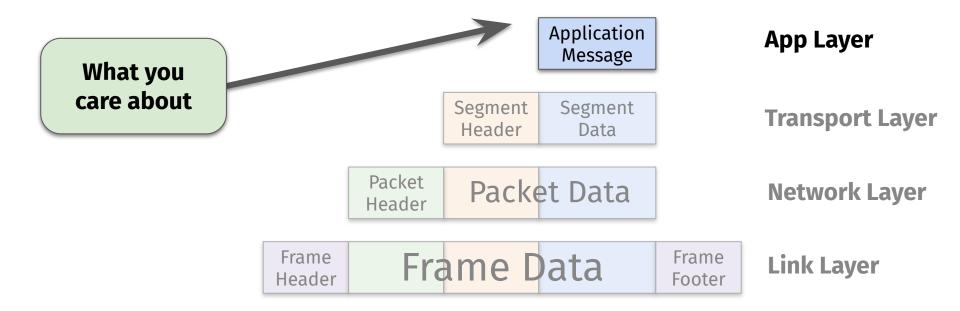
Analyzing Network Packets





Recap: Internet Packet Encapsulation

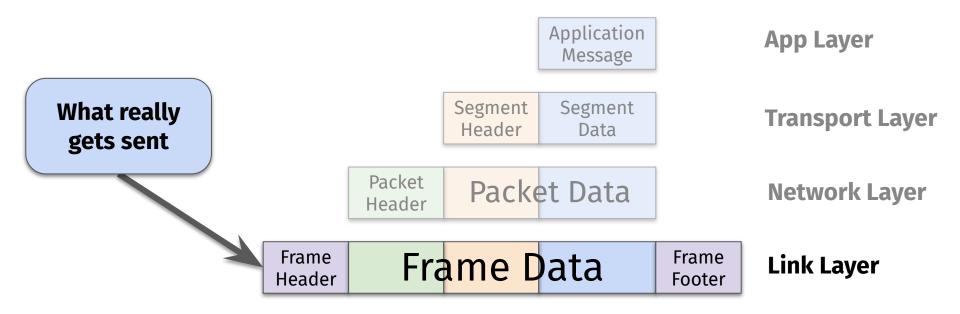
How packets are generated and sent





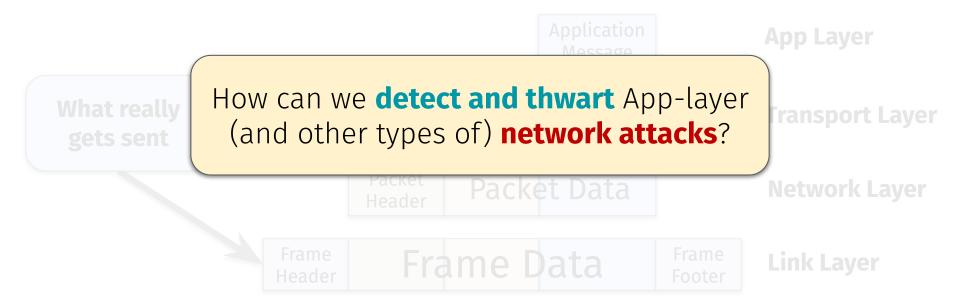
Recap: Internet Packet Encapsulation

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Recap: Internet Packet Encapsulation

How packets are generated and sent





Tools of the Trade

Packet Analyzers:

- Tools for dissecting network packets
- Packet Analyzers allow you to:
 - Identify unusual packets
 - Characterize network activity
 - Pinpoint malicious traffic
- The basis of modern-day network security (e.g., firewalls, antivirus)



Familiarity with packet analysis tools?

I eat NetSec CTF challenges like a kid eats candy on Halloween. 🧙

Some (e.g., Wireshark, DPKT, Scapy, or something else)

None (but that's totally okay!)





0%

0%

0%



Tools of the Trade: Wireshark

- A "graphical interface" for manual packet analysis
 - Completely open-source and free
- General workflow:
 - Load up a PCAP (packet capture)
 - Wireshark will display each packet
 - Inspect particular fields of interest

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lo.		Time			So	urce				Destir	ation				Proto	col	Lengt	h Ir	
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	442	4.402	758		19	2.16	8.1.1	.06		192.	168.1	1.213	3		syne		e	8	
	453	5.042	758		19	2.16	8.1.1	.06		192.	168.1	1.213	3		syne	rgy	6	8	
	464	5.290	740		19	2.16	8.1.1	06		192.	168.1	L.213	3		syne	rgy	e	8	
	486	5.826	760		19	2.16	8.1.1	.06		192.	168.1	L.213	3		syne	rgy	e	8	
	494	5.978	736		19	2.16	8.1.1	.06		192.	168.1	1.213	3		syne	rgy	e	8	
	512	6.186	737		19	2.16	8.1.1	.06		192.	168.1	1.213	3		syne	rgy	e	8	
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Fraid Etl Traid Syrift Syrift	532 ame 4 herne ansm: nergy Pack Pack Key Key Key Key	6.420 428: et II et Pr issio y Pro ket L ket T Rele key M key B	68 b , Sr otoco n Co toco engt ype: ased d: 1 odif utto	c: / ol \ ntro l h: 1 Key 16 ier n: 2	s on ASUS /ers: bl Pi 10 / Rel Mask	wir TekC ion roto	e (54 _14:f 4, Sr col, ed (D	4 bit 6:e8 c: 19 Src F KUP)	(38:0 92.167 Port:	58 by 15:47 3.1.10 2480	tes c :14:f 06, D 0, Ds	aptu 6:e8 0st: t Po	red), D 192. ort:	(544 st:) 168.: 4972	bit: Apple 1.21 7, Se	s) o e_9d 3 eq:	n in :dc: 1091	> terf 83 (ace a0:7 k: 7
<pre>c Fra Fra Fra Tra Syn V 00000</pre>	ame 4 herne ansm: Pack Pack Key k k k k a0	428: et II et Pro issio y Pro ket L ket T Rele (ey I (ey M	68 b , Sr otoco engt ype: ased d: 1: odif utto	c: / ol \ ntro l h: 1 Key 16 ier n: 2 dc	s on ASUS /ers: bl Pi 10 (Re] Mask 20	wir TekC roto Lease	e (54 _14:f 4, Sr col, ed (D) .92	4 bit 6:e8 c: 19 Src F KUP)	(38:0 92.16	08 00	tes c :14:f 06, D 0, Ds	aptu 6:e8 Pst: t Po	ired i), D 192. rrt:	(544 st: / 168.: 4972	bit: Appla 1.211 7, 50	s) o e_9d 3 eq:	n in :dc: 1091 E.	> terf 83 (a0:7
<pre>c Fra Fra Fra Tra Syn V 00000</pre>	ame A herne terne ansm: Pack Key K k k k k k k a0 00	6.420 428: et II et Pro issio y Pro ket L ket T Rele key I key B key B	68 b , Sr otoco engt ype: ased d: 1. odif utto 7 9d 3 da	c: / ol \ ntro l h: 1 Key 16 ier n: 2 dc 40	5 on ASUS /ers: 51 Pi 10 7 Re] 00 83 83 00 8	wir TekC roto Lease : 81	e (54 _14:f 4, Sr col, ed (D) .92	4 bit 6:e8 5rc F KUP) 14 f 00 c	(38:0 92.168 Port:	08 00 01 6a	tes c :14:f 06, D 0, Ds 0, Ds 0, 45 1 c0	000 a8	ired 3), D 192. ort: .x.	(544 st: / 168.: 4972	bit: Appld 1.211 7, Se	s) o =_9d 3 eq:	n in :dc: 1091 E.	> terf 83 (ace a0:7 k: 7

Tools of the Trade: Wireshark

	Time	Source	Destination	Protocol	Length Info	
	1 0.000000	10.0.0.2	10.128.0.2	TCP	54 3341 → 80 [SYN] Seq=0 Win=512 Len=0	
	2 0.003987	10.128.0.2	10.0.0.2	ТСР	58 80 → 3222 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	3 0.005514	10.128.0.2	10.0.0.2	TCP	58 80 → 3341 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	4 0.008429	10.0.0.2	10.128.0.2	TCP	54 3342 → 80 [SYN] Seq=0 Win=512 Len=0	
	5 0.010233	10.128.0.2	10.0.0.2	TCP	58 80 → 3220 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	6 0.014072	10.128.0.2	10.0.0.2	TCP	58 80 → 3342 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	7 0.016830	10.0.0.2	10.128.0.2	TCP	54 3343 → 80 [SYN] Seq=0 Win=512 Len=0	
	8 0.022220	10.128.0.2	10.0.0.2	TCP	58 80 → 3343 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	9 0.023496	10.128.0.2	10.0.0.2	TCP	58 80 → 3219 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	10 0.025243	10.0.0.2	10.128.0.2	TCP	54 3344 → 80 [SYN] Seq=0 Win=512 Len=0	
	11 0.026672	10.128.0.2	10.0.0.2	TCP	58 80 → 3218 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	12 0.028038	10.128.0.2	10.0.0.2	TCP	58 80 → 3221 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
	13 0.030523	10.128.0.2	10.0.0.2	TCP	58 80 → 3344 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460	
Tra S D [nsmission Contr Gource Port: 80 Destination Por Stream index:	ol Protocol, Src P t: 3222 1]	.128.0.2, Dst: 10.0.0 ort: 80, Dst Port: 32		Ack: 1, Len: θ	
Tra S [[S	nsmission Contr Source Port: 80 Destination Por Stream index: TCP Segment Le Sequence number Next sequence	ol Protocol, Src P 1: 3222 1: 0] : 0 (relative se number: 0 (relat	ort: 80, Dst Port: 32 quence number) ive sequence number)]	222, Seq: 0, A	Ack: 1, Len: θ	
Tra S [[S A	nsmission Contr Gource Port: 80 Destination Por Stream index: TCP Segment Len Sequence number Next sequence Acknowledgment	ol Protocol, Src P 1: 3222 1: 0] : 0 (relative se number: 0 (relat umber: 1 (relat	ort: 80, Dst Port: 32 quence number) ive sequence number)] ive ack number)	222, Seq: 0, A	Ack: 1, Len: 0	
Tra S D [S A O	nsmission Contr Source Port: 80 Destination Por Stream index: TCP Segment Lei Sequence number Next sequence n Next sequence I cknowledgment I D110 = Head	ol Protocol, Src P 1] 1] 10] 10] 10] 10] 10] 11] 11] 11] 1	ort: 80, Dst Port: 32 quence number) ive sequence number)] ive ack number)	222, Seq: 0, A	Ack: 1, Len: Θ	
Tra S D S S S S F	nsmission Contr Gource Port: 80 Destination Por Stream index: TCP Segment Len Sequence number Next sequence Acknowledgment	ol Protocol, Src P t: 3222 1] 1: 0] 1: 0 (relative se 1: 0 (relat	ort: 80, Dst Port: 32 quence number) ive sequence number)] ive ack number)	222, Seq: 0, A	Ack: 1, Len: θ	
Tra S D L S S S S S S S S S S S S S S S S S	nsmission Conti Source Port: 800 lestination Por Stream index: : TCP Segment Le Gequence number Next sequence Laknowledgment : 110 = Heaa Lags: 0x012 (S indow size val	ol Protocol, Src P t: 3222 1] 1: 0] 1: 0 (relative se 1: 0 (relat	ort: 80, Dst Port: 32 quence number) ive sequence number)] ive ack number)	222, Seq: 0, A	Ack: 1, Len: θ	
Tra S C C S C C C C C S C C C C C C C C C	nsmission Conti Source Port: 800 lestination Por Stream index: : TCP Segment Le Gequence number Next sequence Laknowledgment : 110 = Heaa Lags: 0x012 (S indow size val	ol Protocol, Src P t: 3222 1] 1: 0] : 0 (relative se humber: 0 (relat humber: 1 (relat der Length: 24 byte (N, ACK) Je: 29200 dow size: 29200]	ort: 80, Dst Port: 32 quence number) ive sequence number)] ive ack number)	222, Seq: 0, A	Ack: 1, Len: θ	
Tra S C L S S C L S S C C C C C C C C C C C	nsmission Conti Source Port: 80 Destination Por Stream index: : TCP Segment Le equence number Next sequence : Next sequence : Next sequence : Next sequence : lags: 0x012 (S' findow size val Calculated win checksum: 0x426 Checksum Statu	ol Protocol, Src P t: 3222 1] 1: 0] 2: 0 (relative se 1) 1: 0 (relative se 1) 1: 0 (relative se 1: 0 (relative se 1: 0 (relative se 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1	ort: 80, Dst Port: 32 quence number) ive sequence number)] ive ack number)	222, Seq: 0, A	Ack: 1, Len: 0	
Tra S C L S S C C C C C C C C C C C C C C C	nsmission Conti Source Port: 800 Lestination Por Stream index: : TCP Segment Les Gequence number Next sequence : Next sequence : Lindow size val Calculated win checksum: 0x426 Checksum Statu rgent pointer:	ol Protocol, Src P t: 3222 1] 1: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte (N, ACK) 16: 29200 dow size: 29200 3 [unverified] 0	ort: 80, Dst Port: 32 equence number) ive sequence number)] ive ack number) s (6)	222, Seq: 0, A	Ack: 1, Len: θ	
Tra D C C C C C C C C C C C C C C C C C C	nsmission Contr Source Port: 80 Destination Por Stream index: : TCP Segment Lei equence number Next sequence i cknowledgment i 110 = Hear flags: 0x012 (S findow size val Calculated win checksum: 0x426 Checksum Statu urgent pointer: ptions: (4 byt	ol Protocol, Src P t: 3222 1] 1: 0] 2: 0 (relative se 1) 1: 0 (relative se 1) 1: 0 (relative se 1: 0 (relative se 1: 0 (relative se 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1: 1	ort: 80, Dst Port: 32 equence number) ive sequence number)] ive ack number) s (6)	222, Seq: 0, A	Ack: 1, Len: θ	
Tra D C C C C C C C C C C C C C C C C C C	nsmission Conti Source Port: 800 Lestination Por Stream index: : TCP Segment Les Gequence number Next sequence : Next sequence : Lindow size val Calculated win checksum: 0x426 Checksum Statu rgent pointer:	ol Protocol, Src P t: 3222 1] 1: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte (N, ACK) 16: 29200 dow size: 29200 3 [unverified] 0	ort: 80, Dst Port: 32 equence number) ive sequence number)] ive ack number) s (6)	222, Seq: 0, A	Ack: 1, Len: θ	
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Tra S C C S C C C C C C C C C C C C C C C	nsmission Contr Source Port: 80 Destination Por Stream index: : TCP Segment Lei equence number Next sequence i cknowledgment i 110 = Hear flags: 0x012 (S findow size val Calculated win checksum: 0x426 Checksum Statu urgent pointer: ptions: (4 byt	ol Protocol, Src P t: 3222 1] 1: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte (N, ACK) 16: 29200 dow size: 29200 3 [unverified] 0	ort: 80, Dst Port: 32 equence number) ive sequence number)] ive ack number) s (6)	222, Seq: 0, A	Ack: 1, Len: θ	

Tools of the Trade: Scapy

- Python API for programmatic packet capture and analysis
 - Think of it as "Wireshark in API form"
 - Project 4: you will use Scapy to write your own packet analysis scripts

Scapy API reference		
scapy.asn1 package	希 / Scapy API reference	O Edit on GitHut
scapy.contrib package		
scapy.layers package	Scony ADI reference	
scapy.ansmachine	Scapy API reference	
scapy.as_resolvers	Scapy: create, send, sniff, dissect and manipulate network	nackets
scapy.asn1fields	ocapy, ereate, seria, sinn, assect and manipulate network	puercesi
scapy.asn1packet	Usable either from an interactive console or as a Python li	brary. https://scapy.net
scapy.automaton		
scapy.autorun	Subpackages	
scapy.base_classes	 scapy.asn1 package 	
scapy.config	scapy.contrib package	
scapy.consts	scapy.layers package	
scapy.dadict		
scapy.data	Submodules	
scapy.error		
scapy.fields	 scapy.ansmachine scapy.as resolvers 	
scapy.interfaces	 scapy.as_resolvers scapy.asn1fields 	
scapy.main	 scapy.asn1packet 	
scapy.packet	scapy.automaton	
scapy.pipetool	scapy.autorun	
	 scapy.base_classes 	

Tools of the Trade: Scapy

- Python API for programmatic packet capture and analysis
 - Think of it as "Wireshark in API form"
 - Project 4: you will use Scapy to write your own packet analysis scripts
- We'll provide the PCAP traces...
 - You'll write code to analyze them!
 - Examples:
 - Detecting attacks on a network
 - Finding user credentials
 - Sniffing a user's browsing history

Scapy API reference		
scapy.asn1 package	I Scapy API reference	O Edit on GitHub
scapy.contrib package		
scapy.layers package	Scapy API reference	
scapy.ansmachine	Scapy API reference	
scapy.as_resolvers	Scapy: create, send, sniff, dissect and manipulate network packets.	
scapy.asn1fields		
scapy.asn1packet	Usable either from an interactive console or as a Python library. https://scap	y.net
scapy.automaton		
scapy.autorun	Subpackages	
scapy.base_classes	 scapy.asn1 package 	
scapy.config	scapy.asi1 package scapy.contrib package	
scapy.consts	scapy.layers package	
scapy.dadict		
scapy.data	Submodules	
scapy.error		
scapy.fields	 scapy.ansmachine scapy.as resolvers 	
scapy.interfaces	 scapy.as_resolvers scapy.asn1fields 	
scapy.main	• scapy.asn1packet	
scapy.packet	scapy.automaton	
scapy.pipetool	scapy.autorun	
scapy.pipetooi	 scapy.base_classes 	

Questions?





Next time on CS 4440...

Transport, Network, Link, and Physical Layer Attacks

