

# Lecture 12: Hardware for Arithmetic

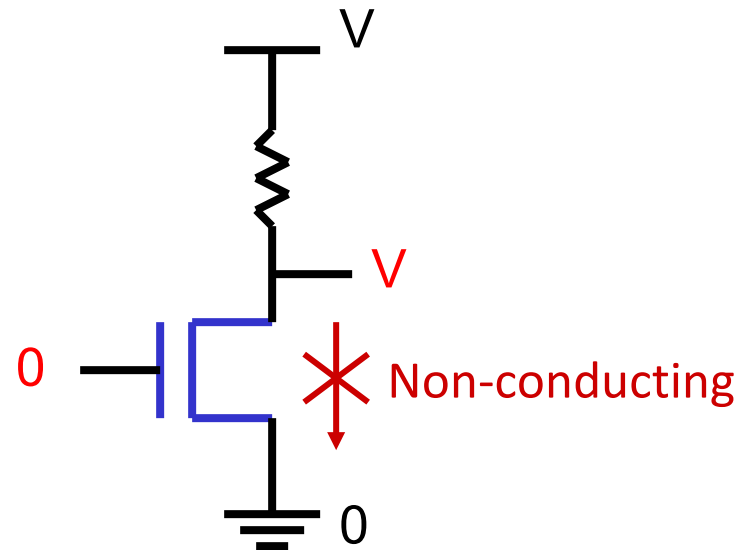
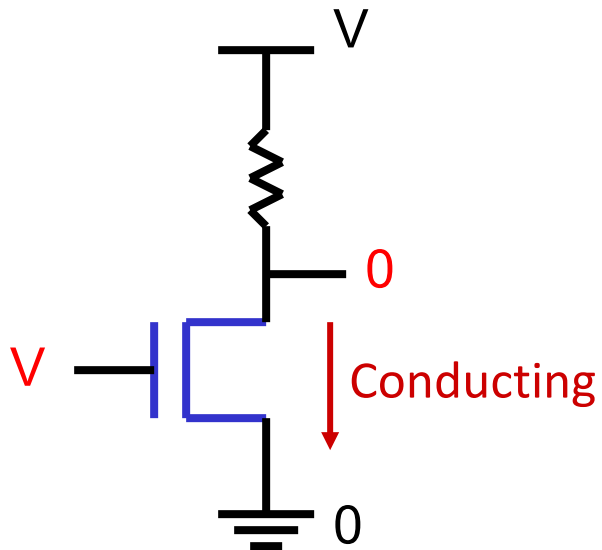
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- Today's topics:
  - Digital logic intro
  - Logic for common operations
  - Designing an ALU

# Digital Design Basics

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- Two voltage levels – high and low (1 and 0, true and false)  
Hence, the use of binary arithmetic/logic in all computers
- A transistor is a 3-terminal device that acts as a switch



# Logic Blocks

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- A logic block has a number of binary inputs and produces a number of binary outputs – the simplest logic block is composed of a few transistors
- A logic block is termed *combinational* if the output is only a function of the inputs
- A logic block is termed *sequential* if the block has some internal memory (state) that also influences the output
- A basic logic block is termed a *gate* (AND, OR, NOT, etc.)

We will only deal with combinational circuits today

# Truth Table

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- A truth table defines the outputs of a logic block for each set of inputs
- Consider a block with 3 inputs A, B, C and an output E that is true only if *exactly* 2 inputs are true

A	B	C	E

# Truth Table

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- A truth table defines the outputs of a logic block for each set of inputs
- Consider a block with 3 inputs A, B, C and an output E that is true only if *exactly* 2 inputs are true

A	B	C	E
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

Can be compressed by only representing cases that have an output of 1

# Boolean Algebra

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- Equations involving two values and three primary operators:
  - OR : symbol  $+$  ,  $X = A + B \rightarrow$  X is true if at least one of A or B is true
  - AND : symbol  $\cdot$  ,  $X = A \cdot B \rightarrow$  X is true if both A and B are true
  - NOT : symbol  $\bar{\quad}$  ,  $X = \bar{A} \rightarrow$  X is the inverted value of A

# Boolean Algebra Rules

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- Identity law :  $A + 0 = A$  ;  $A \cdot 1 = A$
- Zero and One laws :  $A + 1 = 1$  ;  $A \cdot 0 = 0$
- Inverse laws :  $A \cdot \overline{A} = 0$  ;  $A + \overline{A} = 1$
- Commutative laws :  $A + B = B + A$  ;  $A \cdot B = B \cdot A$
- Associative laws :  $A + (B + C) = (A + B) + C$   
 $A \cdot (B \cdot C) = (A \cdot B) \cdot C$
- Distributive laws :  $A \cdot (B + C) = (A \cdot B) + (A \cdot C)$   
 $A + (B \cdot C) = (A + B) \cdot (A + C)$

# DeMorgan's Laws

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- $\overline{A + B} = \overline{A} \cdot \overline{B}$

- $\overline{A \cdot B} = \overline{A} + \overline{B}$

- Confirm that these are indeed true



# Pictorial Representations

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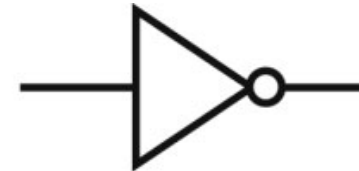
AND



OR



NOT



Source: H&P textbook

What logic function is this?



Source: H&P textbook

# Boolean Equation

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- Consider the logic block that has an output E that is true only if exactly two of the three inputs A, B, C are true

Multiple correct equations:

Two must be true, but all three cannot be true:

$$E = ((A \cdot B) + (B \cdot C) + (A \cdot C)) \cdot \overline{(A \cdot B \cdot C)}$$

Identify the three cases where it is true:

$$E = (A \cdot B \cdot \overline{C}) + (A \cdot C \cdot \overline{B}) + (C \cdot B \cdot \overline{A})$$

# Sum of Products

- Can represent any logic block with the AND, OR, NOT operators
  - Draw the truth table
  - For each true output, represent the corresponding inputs as a product
  - The final equation is a sum of these products

A	B	C	E
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

$$(A \cdot B \cdot \overline{C}) + (A \cdot C \cdot \overline{B}) + (C \cdot B \cdot \overline{A})$$

- Can also use “product of sums”
- Any equation can be implemented with an array of ANDs, followed by an array of ORs

# NAND and NOR

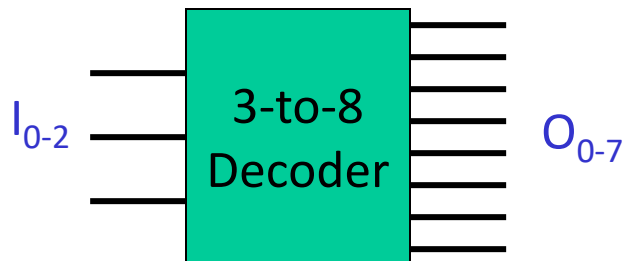
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- NAND : NOT of AND :  $A \text{ nand } B = \overline{A \cdot B}$
- NOR : NOT of OR :  $A \text{ nor } B = \overline{A + B}$
- NAND and NOR are *universal gates*, i.e., they can be used to construct any complex logical function

# Common Logic Blocks – Decoder

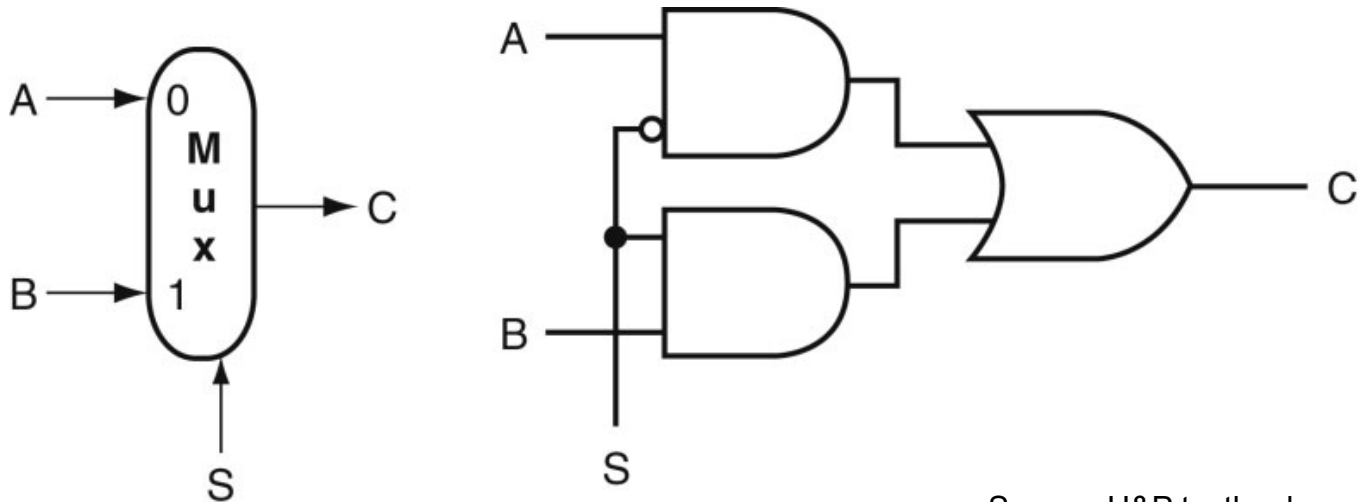
Takes in N inputs and activates one of  $2^N$  outputs

$I_0$	$I_1$	$I_2$	$O_0$	$O_1$	$O_2$	$O_3$	$O_4$	$O_5$	$O_6$	$O_7$
0	0	0	1	0	0	0	0	0	0	0
0	0	1	0	1	0	0	0	0	0	0
0	1	0	0	0	1	0	0	0	0	0
0	1	1	0	0	0	1	0	0	0	0
1	0	0	0	0	0	0	1	0	0	0
1	0	1	0	0	0	0	0	1	0	0
1	1	0	0	0	0	0	0	0	1	0
1	1	1	0	0	0	0	0	0	0	1



# Common Logic Blocks – Multiplexor

- Multiplexor or selector: one of N inputs is reflected on the output depending on the value of the  $\log_2 N$  selector bits



2-input mux

Source: H&P textbook

# Adder Algorithm

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	1	0	0	1
	0	1	0	1
<hr/>				
Sum	1	1	1	0
Carry	0	0	0	1

Truth Table for the above operations:

A	B	Cin	Sum	Cout
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

# Adder Algorithm

	1	0	0	1
	0	1	0	1
Sum	1	1	1	0
Carry	0	0	0	1

Truth Table for the above operations:

A	B	Cin	Sum	Cout
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

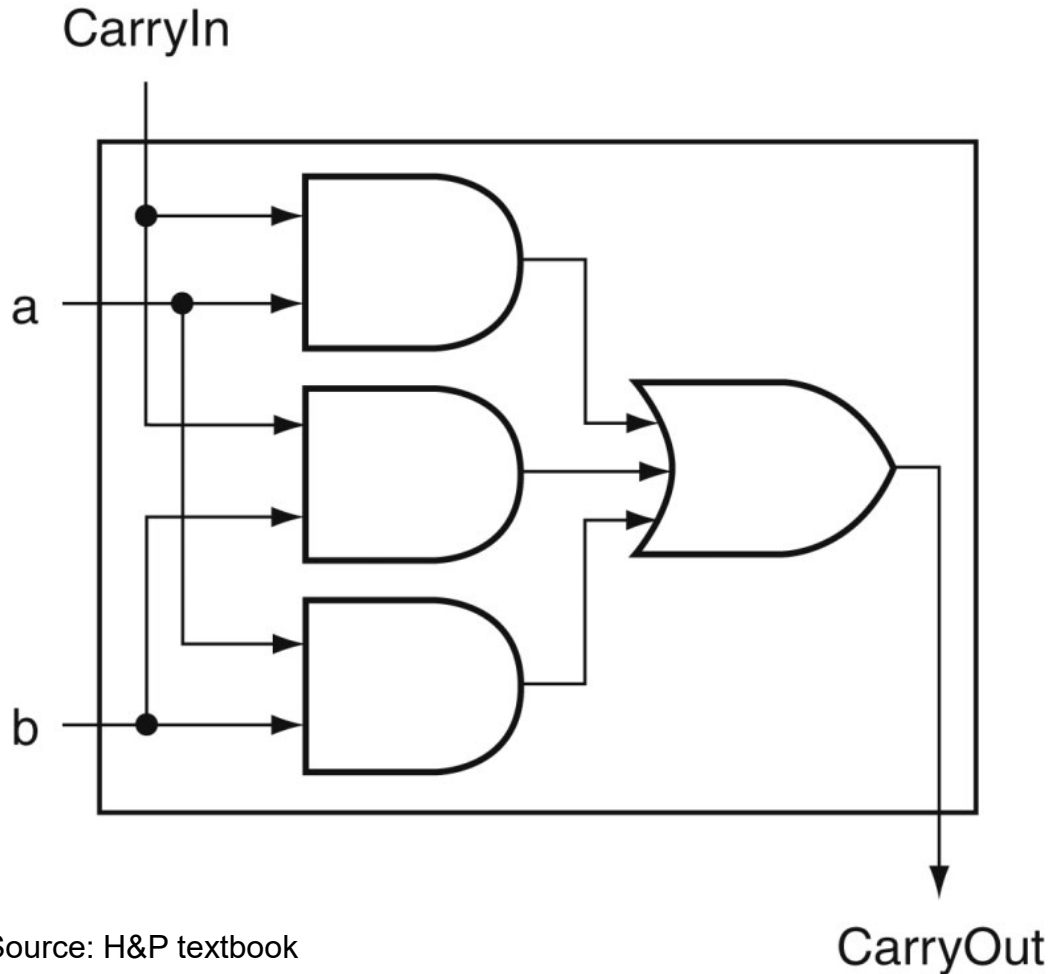
Equations:

$$\begin{aligned} \text{Sum} = & \text{Cin} \cdot \overline{A} \cdot \overline{B} + \\ & B \cdot \overline{\text{Cin}} \cdot \overline{A} + \\ & A \cdot \overline{\text{Cin}} \cdot \overline{B} + \\ & A \cdot B \cdot \text{Cin} \end{aligned}$$

$$\begin{aligned} \text{Cout} = & A \cdot B \cdot \text{Cin} + \\ & A \cdot B \cdot \overline{\text{Cin}} + \\ & A \cdot \text{Cin} \cdot \overline{B} + \\ & B \cdot \text{Cin} \cdot \overline{A} \\ = & A \cdot B + \\ & A \cdot \text{Cin} + \\ & B \cdot \text{Cin} \end{aligned}$$



# Carry Out Logic



Equations:

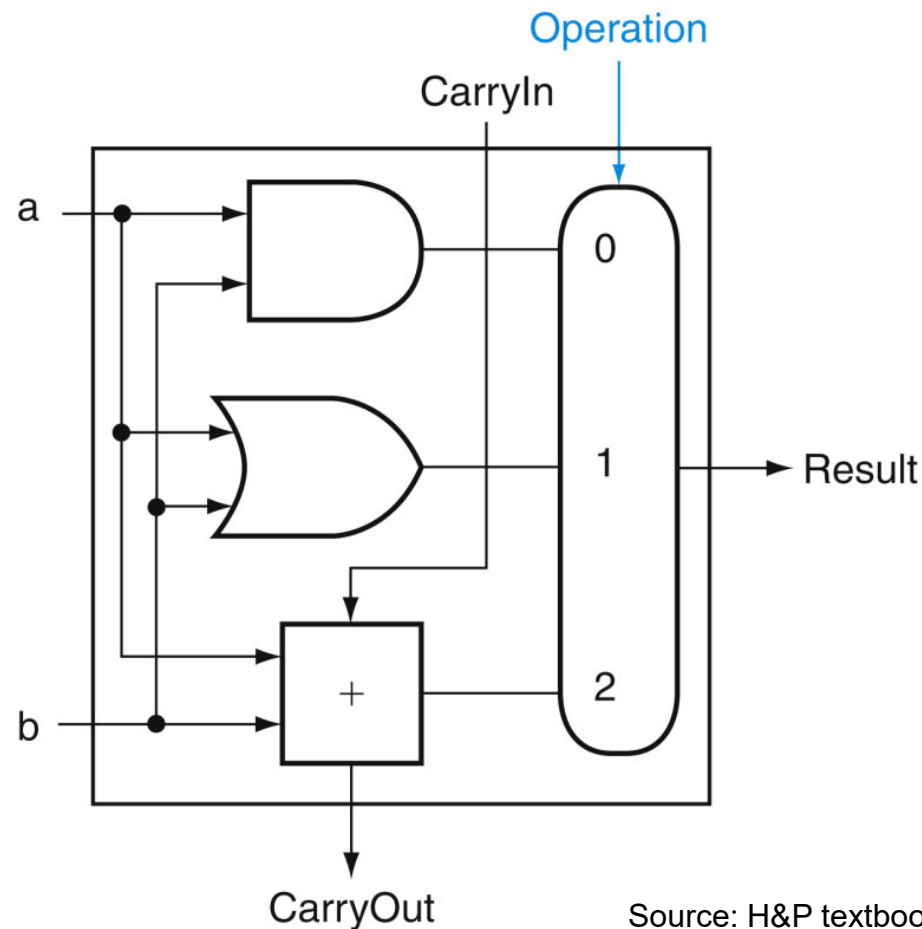
$$\begin{aligned} \text{Sum} = & \text{Cin} \cdot \bar{A} \cdot \bar{B} + \\ & B \cdot \bar{\text{Cin}} \cdot \bar{A} + \\ & A \cdot \bar{\text{Cin}} \cdot B + \\ & A \cdot B \cdot \text{Cin} \end{aligned}$$

$$\begin{aligned} \text{Cout} = & A \cdot B \cdot \text{Cin} + \\ & A \cdot B \cdot \bar{\text{Cin}} + \\ & A \cdot \text{Cin} \cdot \bar{B} + \\ & B \cdot \text{Cin} \cdot \bar{A} \\ = & A \cdot B + \\ & A \cdot \text{Cin} + \\ & B \cdot \text{Cin} \end{aligned}$$

Source: H&P textbook

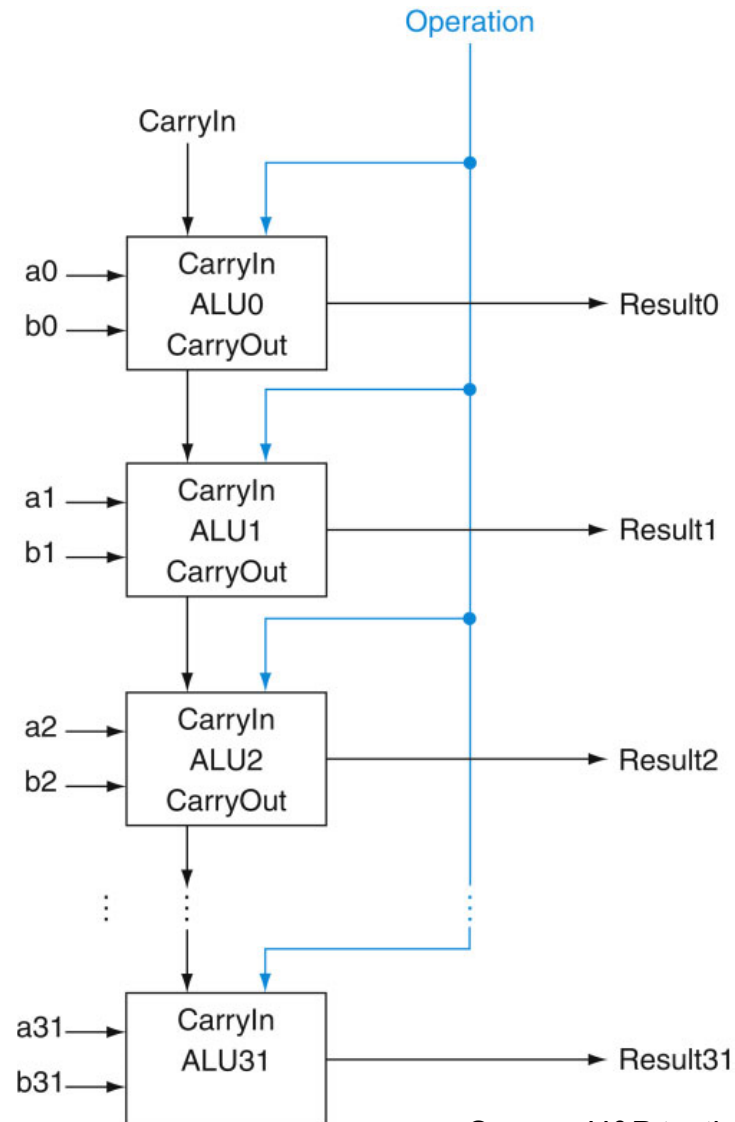
# 1-Bit ALU with Add, Or, And

- Multiplexor selects between Add, Or, And operations



# 32-bit Ripple Carry Adder

1-bit ALUs are connected  
“in series” with the  
carry-out of 1 box  
going into the carry-in  
of the next box

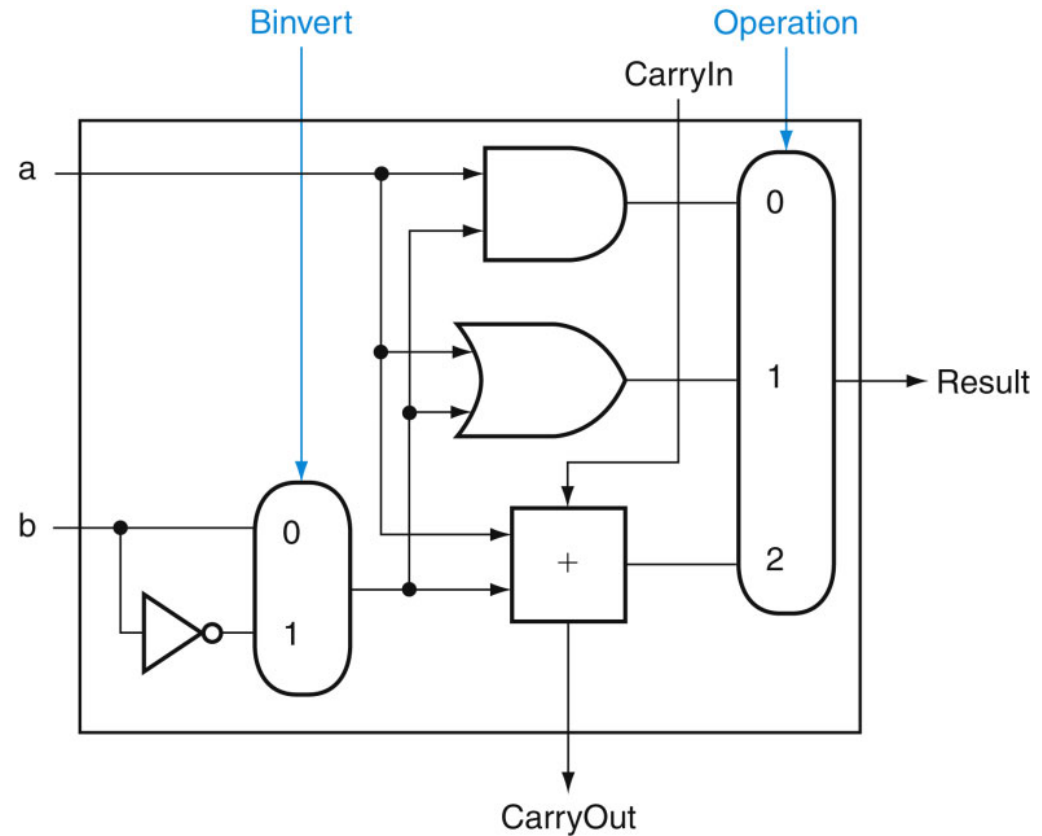


Source: H&P textbook

# Incorporating Subtraction

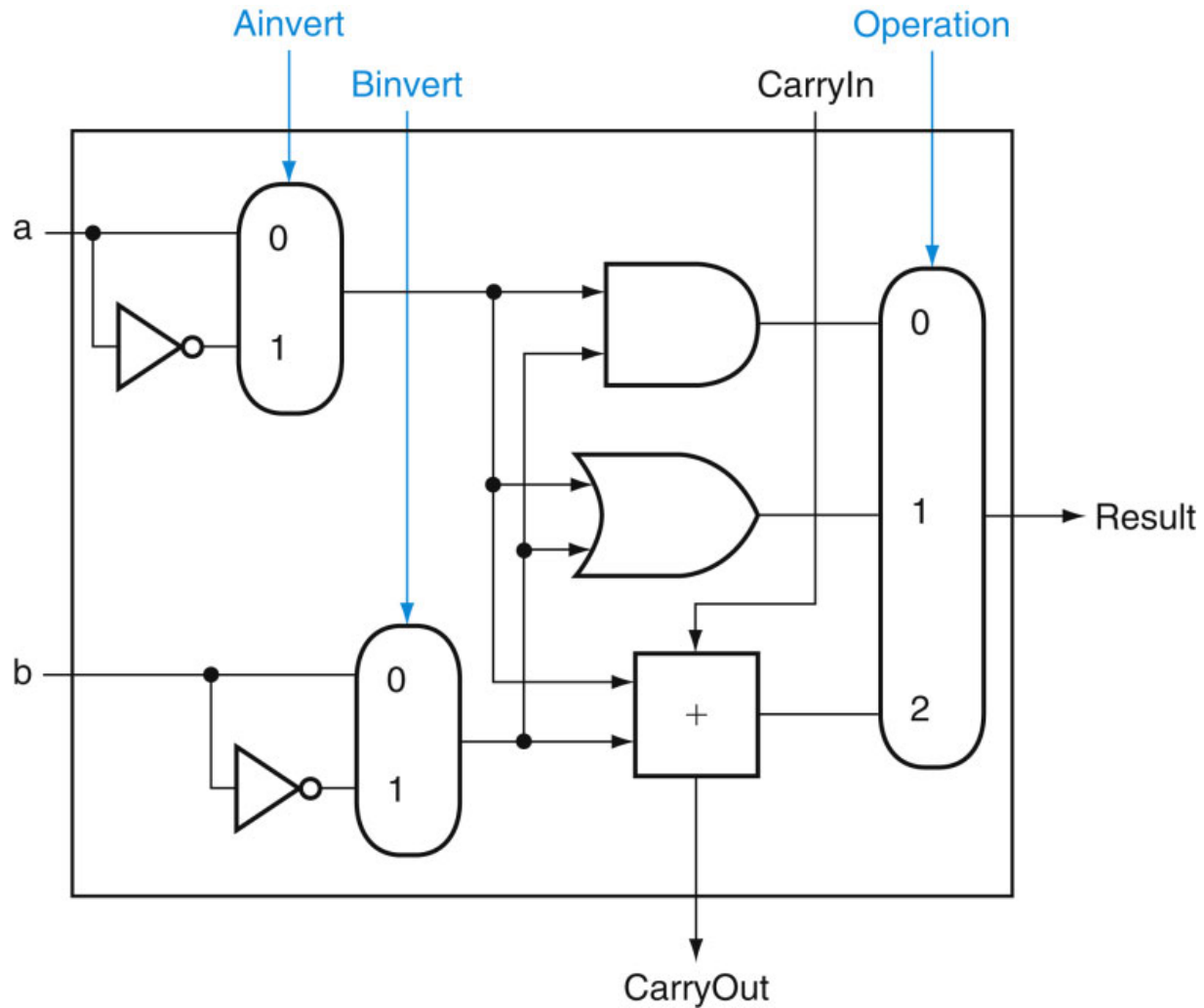
Must invert bits of B and add a 1

- Include an inverter
- CarryIn for the first bit is 1
- The CarryIn signal (for the first bit) can be the same as the Binvert signal



Source: H&P textbook

# Incorporating NOR and NAND



# Control Lines

What are the values of the control lines and what operations do they correspond to?

	Ai	Bn	Op
AND	0	0	00
OR	0	0	01
Add	0	0	10
Sub	0	1	10
NAND	1	1	01
NOR	1	1	00

