

From FP to OOP

Start with data:

```
; A posn is  
; (make-posn num num)  
(define-struct posn (x y))
```



```
class Posn {  
    int x;  
    int y;  
    Posn(int x, int y) {  
        this.x = x;  
        this.y = y;  
    }  
}
```

From FP to OOP

Start with data:

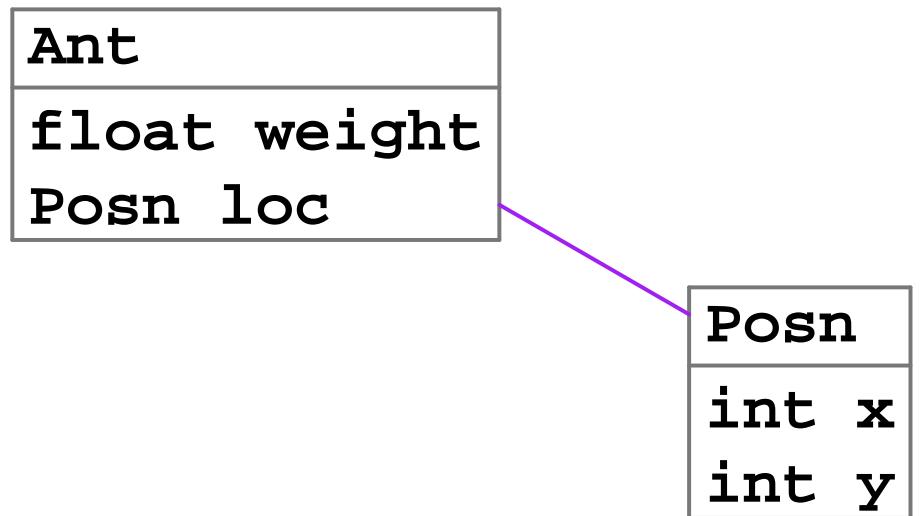
```
; A posn is  
; (make-posn num num)  
(define-struct posn (x y))
```



Posn
int x
int y

From FP to OOP

```
; An ant is  
; (make-ant num posn)  
  
; A posn is  
; (make-posn num num)
```



From FP to OOP

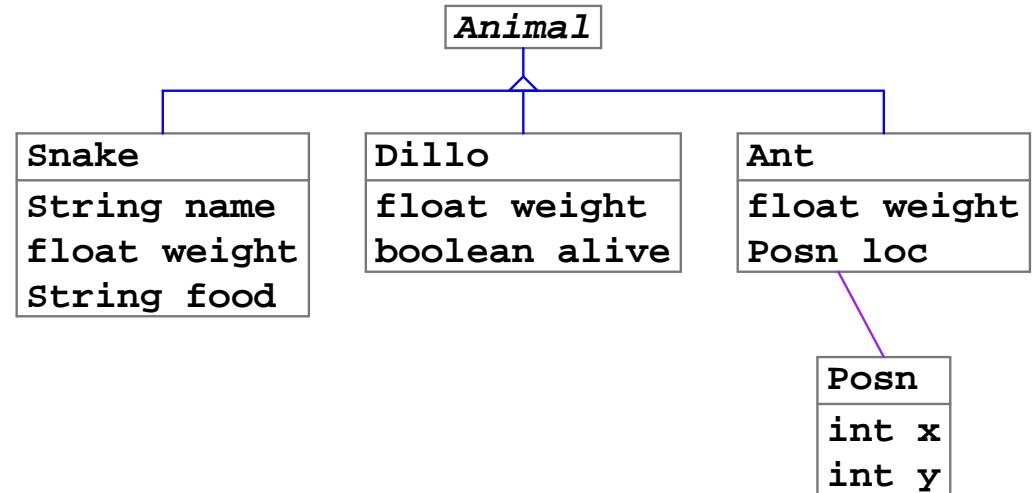
```
; An animal is either
; - snake
; - dillo
; - ant

; A snake is
; (make-snake sym num sym)

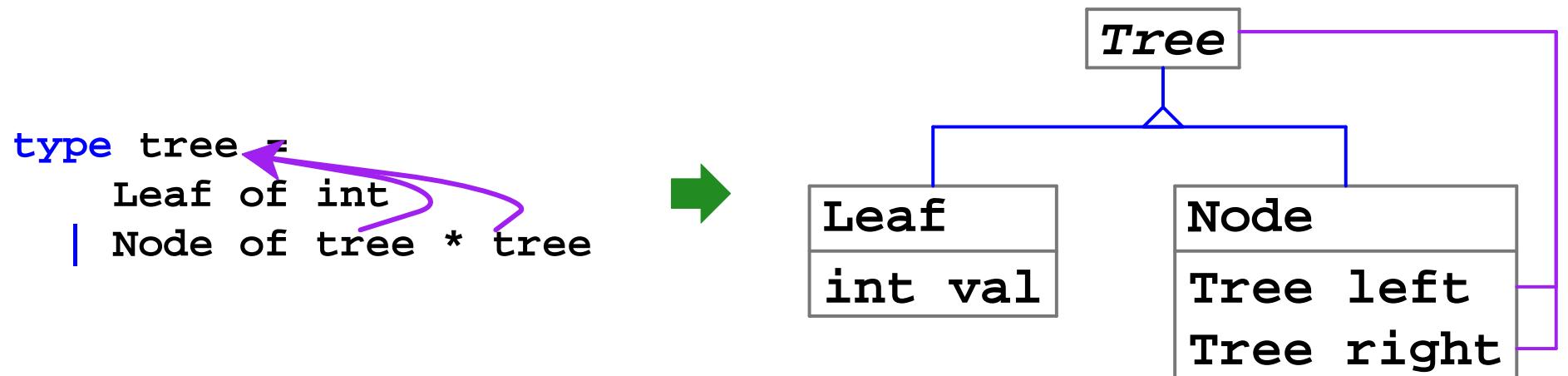
; A dillo is
; (make-dillo num bool)

; An ant is
; (make-ant num posn)

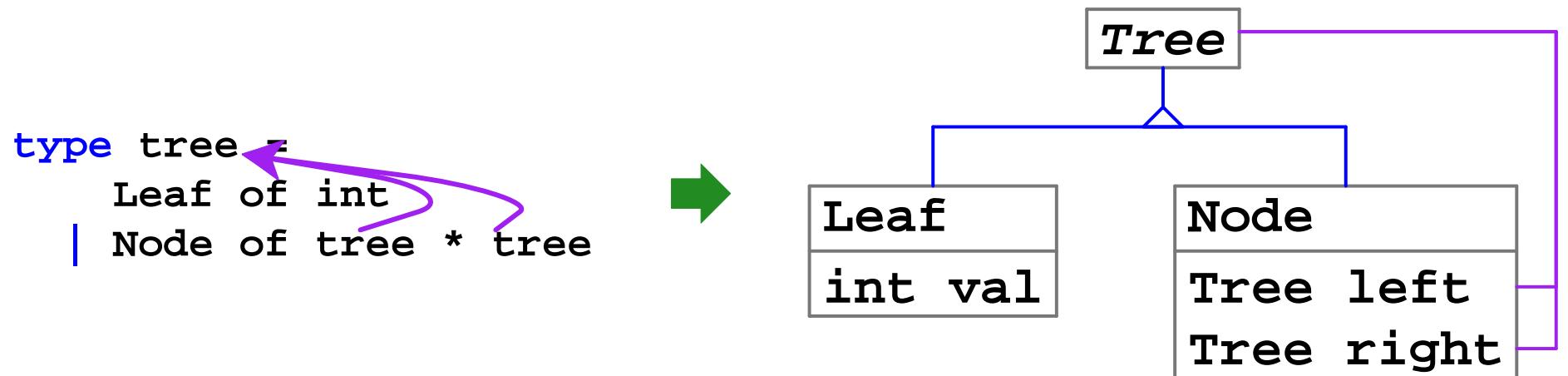
; A posn is
; (make-posn num num)
```



From FP to OOP



From FP to OOP



And so on (for mutually referential data definitions)...

From Functions to Methods

```
; An animal is either
; - snake
; - dillo
; - ant
; ...

; animal-is-lighter? : animal num -> bool
(define (animal-is-lighter? a n)
  (cond
    [(snake? a) (snake-is-lighter? s n)]
    [(dillo? a) (dillo-is-lighter? s n)]
    [(ant? a) (ant-is-lighter? s n)]))

; snake-is-lighter? : snake num -> bool
(define (snake-is-lighter? s n) ...)

; dillo-is-lighter? : dillo num -> bool
(define (dillo-is-lighter? d n) ...)

; ant-is-lighter? : ant num -> bool
(define (ant-is-lighter? a n) ...)
```

```
interface Animal {
  boolean isLighter(double n);
}

class Snake extends Animal {
  ...
  boolean isLighter(double n) { ... }
}

class Dillo extends Animal {
  ...
  boolean isLighter(double n) { ... }
}

class Ant extends Animal {
  ...
  boolean isLighter(double n) { ... }
}
```

From Functions to Methods

```
; An animal is either  
; - snake  
; - dillo  
; - ant  
; ...
```

```
; animal-is-lighter? : animal num -> bool  
(define (animal-is-lighter? a n)  
  (cond  
    [(snake? a) (snake-is-lighter? s n)]  
    [(dillo? a) (dillo-is-lighter? s n)]  
    [(ant? a) (ant-is-lighter? s n)]))  
  
; snake-is-lighter? : snake num -> bool  
(define (snake-is-lighter? s n) ...)  
  
; dillo-is-lighter? : dillo num -> bool  
(define (dillo-is-lighter? d n) ...)  
  
; ant-is-lighter? : ant num -> bool  
(define (ant-is-lighter? a n) ...)
```

Data definition turns into class declarations

```
interface Animal {  
  boolean isLighter(double n);  
}  
  
class Snake extends Animal {  
  ...  
  boolean isLighter(double n) { ... }  
}  
  
class Dillo extends Animal {  
  ...  
  boolean isLighter(double n) { ... }  
}  
  
class Ant extends Animal {  
  ...  
  boolean isLighter(double n) { ... }  
}
```

From Functions to Methods

```
; An animal is either
; - snake
; - dillo
; - ant
; ...

; animal-is-lighter? : animal num -> bool
(define (animal-is-lighter? a n)
  (cond
    [(snake? a) (snake-is-lighter? s n)]
    [(dillo? a) (dillo-is-lighter? s n)]
    [(ant? a) (ant-is-lighter? s n)]))

; snake-is-lighter? : snake num -> bool
(define (snake-is-lighter? s n) ...)

; dillo-is-lighter? : dillo num -> bool
(define (dillo-is-lighter? d n) ...)

; ant-is-lighter? : ant num -> bool
(define (ant-is-lighter? a n) ...)
```

Variant functions turn into variant methods – all with the same contract after the implicit argument

```
interface Animal {
  boolean isLighter(double n);
}

class Snake extends Animal {
  ...
  boolean isLighter(double n) { ... }
}

class Dillo extends Animal {
  ...
  boolean isLighter(double n) { ... }
}

class Ant extends Animal {
  ...
  boolean isLighter(double n) { ... }
}
```

From Functions to Methods

```
; An animal is either
; - snake
; - dillo
; - ant
; ...
```

```
; animal-is-lighter? : animal num -> bool
(define (animal-is-lighter? a n)
  (cond
    [(snake? a) (snake-is-lighter? s n)]
    [(dillo? a) (dillo-is-lighter? s n)])
    [(ant? a) (ant-is-lighter? s n)]))
```

```
; snake-is-lighter? : snake num -> bool
(define (snake-is-lighter? s n) ...)
```

```
; dillo-is-lighter? : dillo num -> bool
(define (dillo-is-lighter? d n) ...)
```

```
; ant-is-lighter? : ant num -> bool
(define (ant-is-lighter? a n) ...)
```

Function with variant-based **cond**
turns into just an **abstract**
method declaration

```
interface Animal {
  boolean isLighter(double n);
}

class Snake extends Animal {
  ...
  boolean isLighter(double n) { ... }
}

class Dillo extends Animal {
  ...
  boolean isLighter(double n) { ... }
}

class Ant extends Animal {
  ...
  boolean isLighter(double n) { ... }
}
```

Extensibility Problem

If we need new **animal** operations...

- FP: add a function (no change to old code)
- OOP: change interfaces & classes to add method

If we need new **animal** variants...

- FP: change all functions to add case
- OOP: add a subclass (no change to old code)

Design patterns in each style provide some benefits of the other

The Object Pattern for FP

```
(define-struct animal (lighter?-proc))

(define (animal-is-lighter? a n)
  ((animal-lighter?-proc a) n))

(define (make-snake name weight food)
  (make-animal (lambda (n) (< weight n)))))

...
```

The Visitor Pattern for OOP

```
interface Animal {
    <T> accept(Visitor<T> v);
}
class Snake implements Animal {
    ...
    <T> accept(Visitor<T> v) { return v.visit(this); }
}
...
interface Visitor<T> {
    <T> visit(Snake s);
    ...
}
class IsLighter implements Visitor<boolean> {
    int n;
    ...
    boolean visit(Snake s) { return s.weight < n; }
    ...
}
```

Language Cores

FP core:

- Closures
- Datatype case dispatch
- Parametric polymorphism

OOP core:

- Object creation
- Dynamic method dispatch
- Static method dispatch (e.g., `super`)

Static and Dynamic Dispatch

```
class Snake implements Animal {  
    ...  
    boolean endangers(Animal a) {  
        return (a.slowerThan(100)  
                && a.isLighter(this.weight/2));  
    }  
}  
  
class Rattlesnake extends Snake {  
    ...  
    boolean endangers(Animal a) {  
        return (!a.hasThickSkin()  
                || super.endangers(a))  
    }  
}  
  
Animal a = new Rattlesnake(...);  
Animal b = new Dillo(...);  
  
a.endangers(b);
```

dynamic
static

CAE Grammar

```
<prog> ::= <decl>* <CAE>
<decl> ::= {class <cid> <fid>* <meth>*}
<meth> ::= {<mid> <CAE>}
<CAE> ::= <num>
         | {+ <CAE> <CAE>}
         | {- <CAE> <CAE>}
         | {if0 <CAE> <CAE> <CAE>}
         | arg
         | this
         | {new <cid> <CAE>*}
         | {get <CAE> <fid>}
         | {dsend <CAE> <mid> <CAE>}
         | {ssend <CAE> <cid> <mid> <CAE>}
```

```
{class posn
  x y
  {mdist {+ {get this x} {get this y}}}
  {addDist {+ {dsend arg mdist 0}
            {dsend this mdist 0}}}}
{dsend {new posn 1 2} addDist {new posn 3 4}}
```

Analogous Java code:

```
class Posn {
    int x, y; ...
    int mdist() {
        return this.x + this.y;
    }
    int addDist (Posn p) {
        return p.mdist() + this.mdist();
    }
}
new Posn(1,2).addDist(new Posn(3,4))
```

CAE Grammar

```
<prog> ::= <decl>* <CAE>
<decl> ::= {class <cid> <fid>* <meth>*}
<meth> ::= {<mid> <CAE>}
<CAE> ::= <num>
          | {+ <CAE> <CAE>}
          | {- <CAE> <CAE>}
          | {if0 <CAE> <CAE> <CAE>}
          | arg
          | this
          | {new <cid> <CAE>*}
          | {get <CAE> <fid>}
          | {dsend <CAE> <mid> <CAE>}
          | {ssend <CAE> <cid> <mid> <CAE>}

{class posn ...
 {addDist {+ {dsend arg mdist 0}
           {dsend this mdist 0}}}}
{class posn3D
 x y z
 {mdist {+ {get this z}
         {ssend this posn mdist arg}}}
 {addDist {ssend this posn addDist arg}}}
{send {new posn3D 1 2 3} addDist {new posn 3 4}}
```

```
Analogous Java code:
class Posn {
  ... as before ...
}
class Posn3D extends Posn {
  int z; ...
  int mdist() {
    return this.z + super.mdist();
  }
}
new Posn3D(1,2,3).addDist(new Posn(3,4))
```

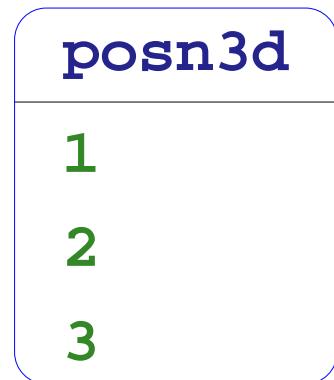
Object Values

How does

```
{send {new posn3D 1 2 3} mdist ...}
```

dispatch to the right `mdist`?

The result of `{new posn3D 1 2 3}` must contain a class tag and field values:



CAE Datatypes

```
type cae =
  Num of int
  | Add of cae * cae
  | Sub of cae * cae
  | IfZ of cae * cae * cae
  | Arg
  | This
  | New of string * cae list
  | Get of cae * string
  | DSend of cae * string * cae
  | SSend of cae * string * string * cae

and cdecl = Class of string * field list * meth list
and field = Field of string
and meth = Method of string * cae

and caeValue =
  NumV of int
  | ObjV of cdecl * caeValue list
```

CAE Interpreter

```
let rec interp : (cae * cdecl list * caeValue * caeValue
                  -> caeValue )
  = function (expr, cdecls, this, arg) ->
    let recur = fun e -> interp(e, cdecls, this, arg)
    in match expr with
      Num(n) -> NumV(n)
      | Add(l, r) -> numPlus(recur l, recur r)
      | Sub(l, r) -> numMinus(recur l, recur r)
      | IfZ(tst, thn, els) ->
          if (recur tst = NumV(0))
          then recur thn
          else recur els
      | This -> this
      | Arg -> arg
    ...
    ...
```

CAE Interpreter

```
let rec interp : (cae * cdecl list * caeValue * caeValue
                  -> caeValue )
  = function (expr, cdecls, this, arg) ->
    let recur = fun e -> interp(e, cdecls, this, arg)
    in match expr with
      ...
    | New(name, exprs) ->
        let decl = findClass name cdecls
        in ObjV(decl, List.map recur exprs)
      ...
    | ...
```

CAE Interpreter

```
let rec interp : (cae * cdecl list * caeValue * caeValue
                  -> caeValue )
  = function (expr, cdecls, this, arg) ->
    let recur = fun e -> interp(e, cdecls, this, arg)
    in match expr with
      ...
    | Get(expr, fname) ->
        (match recur expr with
         ObjV(Class(_, fields, _), vals) ->
           getField fname fields vals
         | _ -> raise (Failed "not an object for get"))
      ...
    
```

CAE Interpreter

```
let rec interp : (cae * cdecl list * caeValue * caeValue
                  -> caeValue )
  = function (expr, cdecls, this, arg) ->
    let recur = fun e -> interp(e, cdecls, this, arg)
    in match expr with
    ...
  | DSend(expr, mname, argExpr) ->
    (match recur expr with
     (ObjV(Class(_, _, methods), _) as this) ->
      let Method(name, body) = findMethod mname methods
      in interp(body, cdecls, this, recur argExpr)
    | _ -> raise (Failed "not an object for send"))
  ...
  ...
```

CAE Interpreter

```
let rec interp : (cae * cdecl list * caeValue * caeValue
                  -> caeValue )
= function (expr, cdecls, this, arg) ->
  let recur = fun e -> interp(e, cdecls, this, arg)
  in match expr with
  ...
| SSend(expr, cname, mname, argExpr) ->
  let this = recur expr
  in let Class(_, _, methods) = findClass cname cdecls
  in let Method(name, body) = findMethod mname methods
  in interp(body, cdecls, this, recur argExpr)
```

CAE Helpers

```
exception NoSuch of string * string

let rec find = fun what nameOf name vals ->
  match vals with
    [] -> raise (NoSuch (what, name))
  | a::rest ->
    if (name = nameOf(a))
    then a
    else find what nameOf name rest

let findClass = (find "class"
                  (fun (Class(name, _, _)) -> name))

let findMethod = (find "method"
                      (fun (Method(name, _)) -> name))

let getField = fun name fields vals ->
  snd (find "field" (fun (Field(name), _) -> name))
  name (List.map2 (fun x y -> (x,y)) fields vals))
```