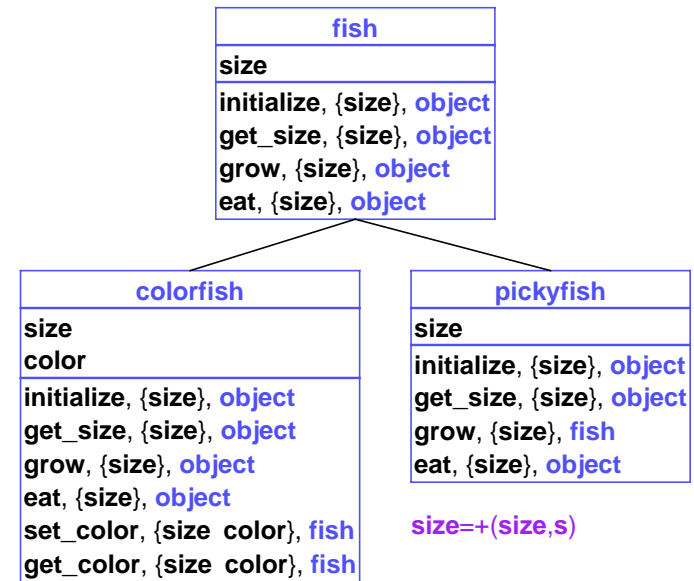


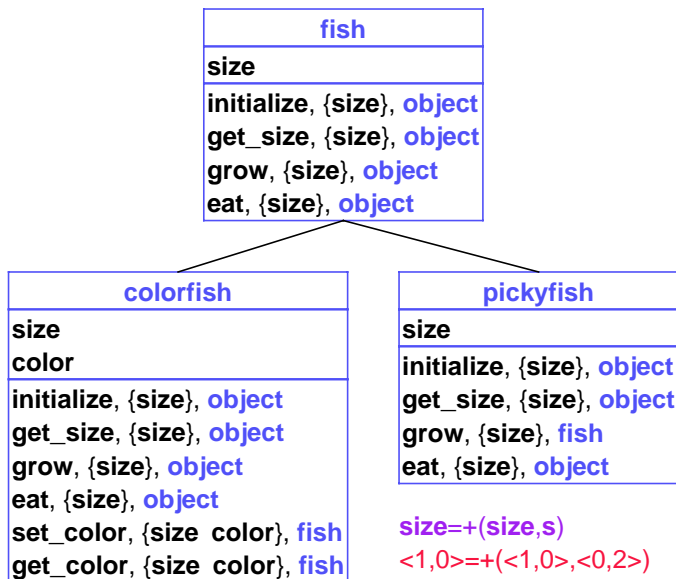
More Optimization

- Still have list walks: variable lookup, method lookup
 - Can eliminate many with lexical addresses
 - Can eliminate some by pre-computing method positions
 - Need type information to eliminate others

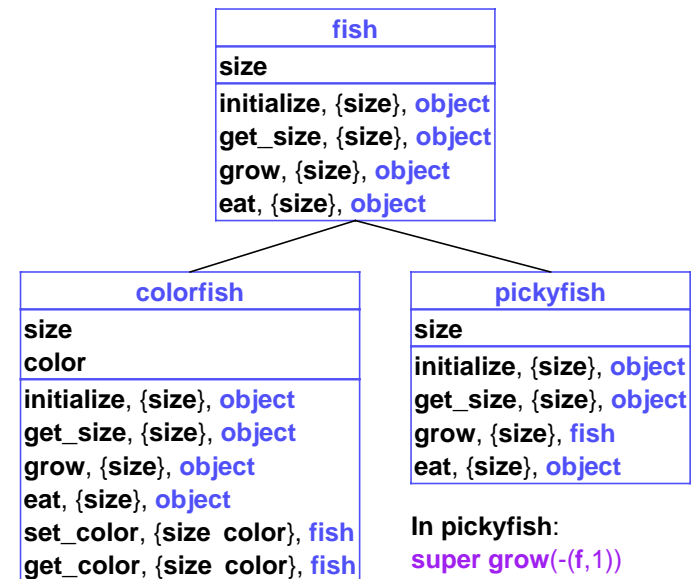
More Optimization: Eliminating List Walks



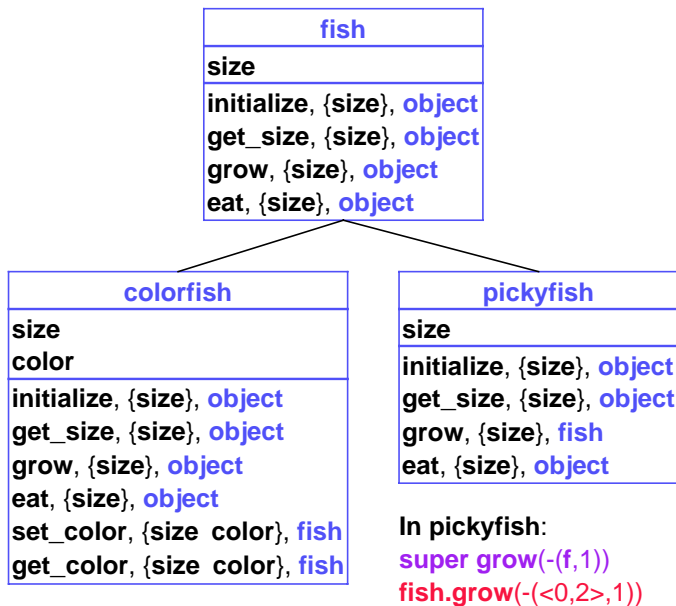
More Optimization: Eliminating List Walks



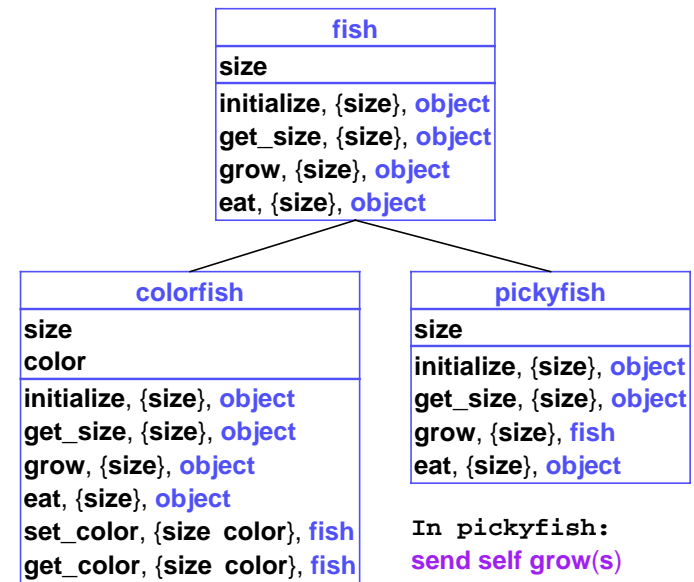
More Optimization: Eliminating List Walks



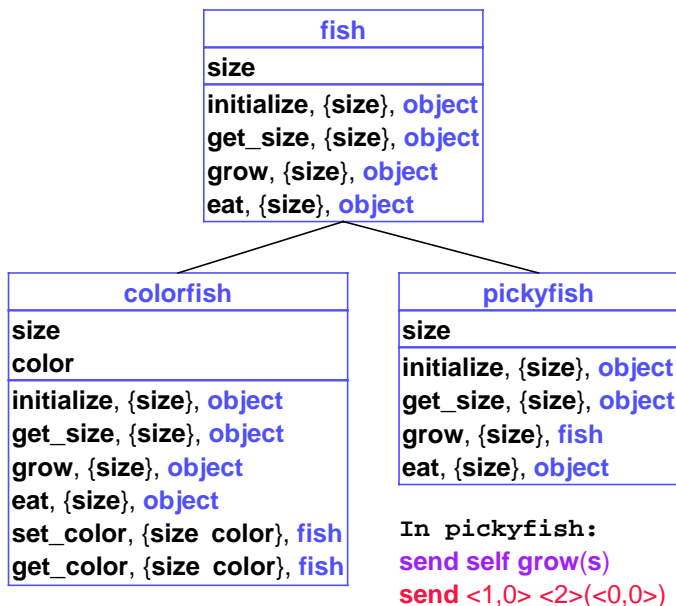
More Optimization: Eliminating List Walks



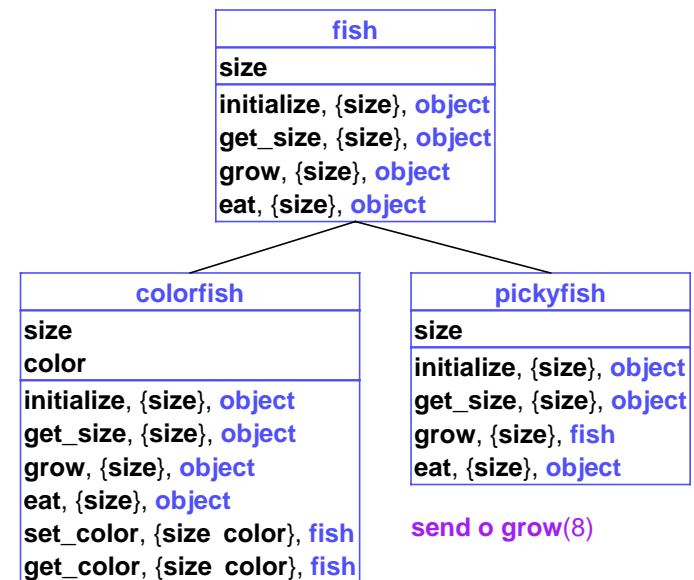
More Optimization: Eliminating List Walks



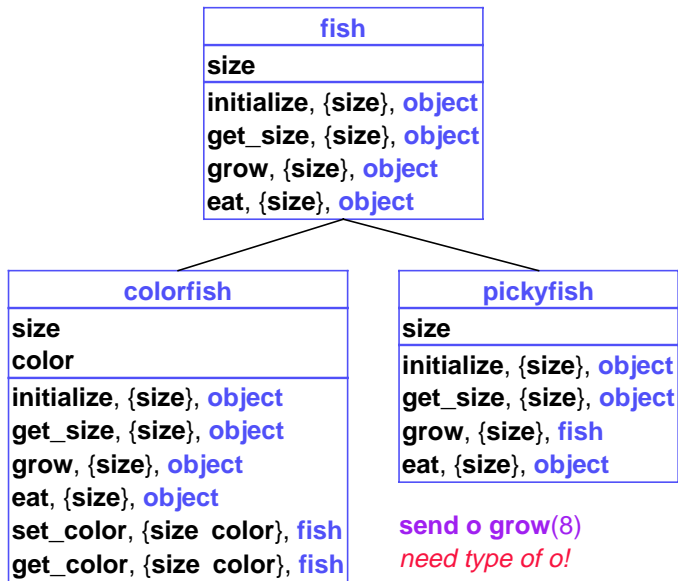
More Optimization: Eliminating List Walks



More Optimization: Eliminating List Walks



More Optimization: Eliminating List Walks



Object Types

```
new c1()
  |
  c1
```

... **if** **c1** has an **initialize** method that takes no arguments

class c1 extends ...
method void initialize() ...

Object Types

```
new c1(5)
  |
  int
  |
  c1
```

... **if** **c1** has an **initialize** method that takes one integer

class c1 extends ...
method void initialize(int v) ...

Object Types

```
send new c1() m(false)
  |         |
  c1       bool
  |
  int
```

... **if** **c1** has an **m** method that takes **bool** and returns **int**

class c1 extends ...
method void initialize() ...
method int m(bool v) ...

Object Types

```
class fish extends object
  field int size
  method void initialize (int s) ...
  method void eat(fish other) ...
class colorfish extends fish
...
```

```
send new fish(8) eat(new colorfish(1))
      fish      | colorfish
colorfish doesn't match fish
```

Subtyping

- **Subtype:** An instance of class **C** can be used as an instance of class **C'** if **C** is derived from **C'**

$$C <: C'$$

- Subtype rule:

If $E \vdash e : T_1$ and $T_1 <: T_2$, then $E \vdash e : T_2$

$$\frac{E \vdash e : T_1 \quad T_1 <: T_2}{E \vdash e : T_2}$$

Object Types

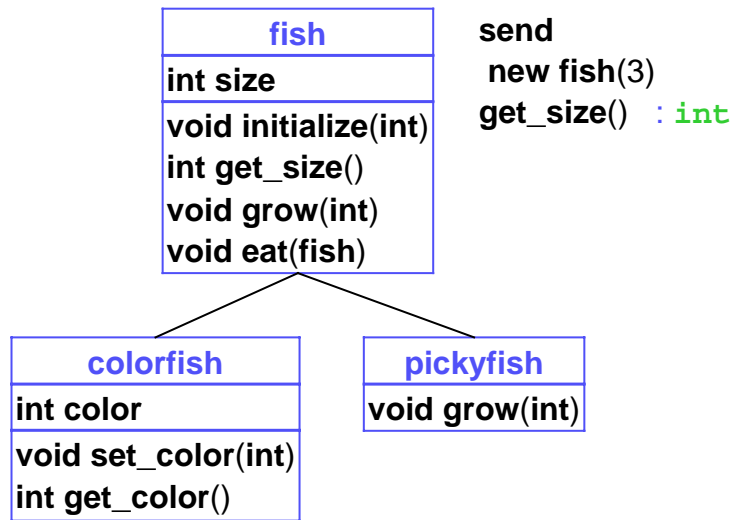
```
class fish extends object
  field int size
  method void initialize (int s) ...
  method void eat(fish other) ...
class colorfish extends fish
...
```

```
send new fish(8) eat(new colorfish(1))
      fish      | colorfish <: fish
                  void
```

Language Changes

- Add types to field declarations
- Add types to method arguments and result
- Add **abstract class** and **abstractmethod**
- Add **cast**

Program Checking



Things to Check

cast:

- Operand has an object type (for any class)
- Target class exists

cast o c1

Things to Check

cast:

- Operand has an object type (for any class)
- Target class exists
- Class for operand and target must be comparable
 - Otherwise, cast cannot possibly succeed

```
class c1 extends object ...
class c2 extends object ...
cast new c1() c2
```

Things to Check

Object creation:

- Class exists, and is not abstract
- Class has an **initialize** method
- **initialize**'s argument types match the operand types

```
class c1 extends object
method void initialize(int x, bool y)
...
new c1(1, false)
```

Things to Check

Method calls:

- Receiver expression is an object
- Method is in the object-type's class
 - Except **initialize...**
- Method's argument types match the operand types

```
class c1 extends object
  method void initialize() ...
  method void m(int x, bool y)
  ...
let o1 = new c1()
in send o1 m(1, false)
```

Things to Check

super calls:

- Expression is within a method
- Method is in the superclass, and not abstract
- Method's argument types match the operand types

```
class c1 extends object
  method void m(int x, bool y)
  ...
class c2 extends c1
  method void n()
  super m(1, false)
  ...
```

Things to Check

class declarations:

- Superclass exists, and no cyclic inheritance
- Methods bodies ok
 - Use host class for type of **self**
- Overriding method signatures are the same as in superclass
 - Except for **initialize**

```
class c2 extends c1
  method void m(int x, bool y)
  if y then +(2, x) else send self w()
```

The Initialize Method

```
class c1 extends obj
  field int x
  method void initialize()
  set x = 3
  method int m()
  send self initialize()

class c2 extends c1
  field int y
  method void initialize(int v)
  set y = v
  super initialize()
  ...
```

- Derived class needs different signature for **initialize**

The Initialize Method

```
class c1 extends obj
  field int x
  method void initialize()
    set x = 3
  method int m()
    send self initialize()
```

```
class c2 extends c1
  field int y
  method void initialize(int v)
    set y = v
    super initialize()
  ...
```

- Disallow **send** to **initialize**

The Initialize Method

```
class c1 extends obj
  field int x
  method void initialize()
    set x = 3
  method int m()
    send self initialize()
```

```
class c2 extends c1
  field int y
  method void initialize(int v)
    set y = v
    super initialize()
  ...
```

- **super** call to **initialize** is ok

Field Initializations

Not checked: field initializations

```
class interior_node extends tree
  field tree left
  field tree right
  method void initialize(tree l, tree r)
  begin
    send left sum();
  ...
end
```

- Can get "bad object 0 for method call"
- This is analogous to the `null` error in Java

Type Checking and Errors

Disallowed errors:

- Object has no such method, or Super method not found
- Can't call method of non-object, non-0
- No such field, no such variable
- Illegal primitive argument (except car of empty)

Allowed errors:

- Can't call method of 0
- Cast failed
- Car of empty

Mixing Subtyping and Procedures

Our language still has procedures:

```
let feed = proc(colorfish f)
  send f grow(10)
  o1 = new colorfish(0)
in
  (feed o1)
```

Mixing Subtyping and Procedures

And higher-order procedures:

```
let feed = proc(colorfish f)
  send f grow(10)
  o1 = new colorfish(0)
  o2 = new colorfish(1)
in let toboth = proc((colorfish -> void) p)
  begin
    (p o1);
    (p o2)
  end
in (toboth feed)
```

Mixing Subtyping and Procedures

Subtyping on procedure arguments:

```
let feed = proc(fish f)
  send f grow(10)
  o1 = new colorfish(0)
in
  (feed o1)
```

- This works, and is allowed by our subtyping rule

Mixing Subtyping and Procedures

Subtyping on procedure arguments:

```
let feed = proc(fish f)
  send f grow(10)
  o1 = new colorfish(0)
  o2 = new colorfish(1)
in let toboth = proc((colorfish -> void) p)
  begin
    (p o1);
    (p o2)
  end
in (toboth feed)
```

- This works, but is **not** allowed by our subtyping rule

(fish → void) versus (colorfish → void)

Procedure Subtyping Rule

If $T_1 <: T_{10}$ and $T_2 <: T_{20}$

then $(T_{10} \rightarrow T_2) <: (T_1 \rightarrow T_{20})$

Another example:

- **dog** <: **animal**
 - a dog can go anywhere an animal can go
- **(animal → hairstyle)** <: **(dog → hairstyle)**
 - a groomer for all animals can groom a dog
 - a groomer who only works with dogs doesn't work for all animals

Procedure Subtyping Rule

If $T_1 <: T_{10}$ and $T_2 <: T_{20}$

then $(T_{10} \rightarrow T_2) <: (T_1 \rightarrow T_{20})$

General intuition:

- $T_1 <: T_{10}$ means T_{10} is more general than T_1



- A function that is willing to accept a more general argument is itself more specific



Procedure Subtyping Rule

If $T_1 <: T_{10}$ and $T_2 <: T_{20}$

then $(T_{10} \rightarrow T_2) <: (T_1 \rightarrow T_{20})$

- Procedure types are **contravariant** with respect to their argument types
- Procedure types are **covariant** with respect to their result types