

## Lazy Evaluation of Function Arguments

```
let f = proc(x)0
in (f +(1,(2,(3,(4,(5,6))))))
```

The computed 21 is never used.

What if we were *lazy* about computing function arguments (in case they aren't used)?

## Lazy Evaluation of Function Arguments

One way to laziness:

```
let f = proc(xthunk)0
in (f proc()+1,(2,(3,(4,(5,6))))))
```

```
let f = proc(xthunk)-((xthunk), 7)
in (f proc()+1,(2,(3,(4,(5,6))))))
```

By using **proc** to delay evaluation, we can avoid unnecessary computation.

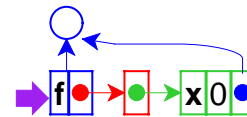
How about making the language compute function arguments lazily in *all* applications?

## Evaluation with Lazy Arguments



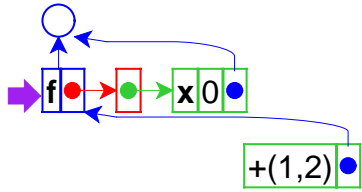
```
let f = proc(x)0
in (f +(1,2))
```

## Evaluation with Lazy Arguments



```
let f = proc(x)0
in (f +(1,2))
```

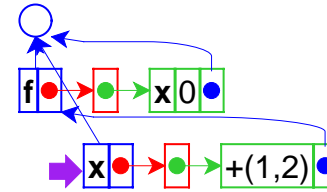
### Evaluation with Lazy Arguments



Application creates a new kind of green box, with two slots: a *thunk*

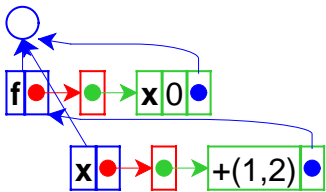
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let f = proc(x)0
in (f +(1,2))
```

### Evaluation with Lazy Arguments



```
let f = proc(x)0
in (f +(1,2))
```

### Evaluation with Lazy Arguments



The result is 0

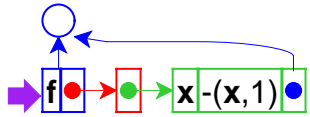
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let f = proc(x)0
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### Evaluation with Lazy Arguments



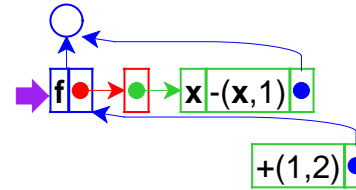
```
let f = proc(x)-(x,1)
in (f +(1,2))
```

### Evaluation with Lazy Arguments



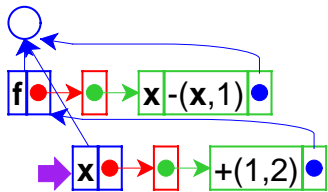
let f = proc(x)-(x,1)  
in (f +(1,2))

### Evaluation with Lazy Arguments



let f = proc(x)-(x,1)  
in (f +(1,2))

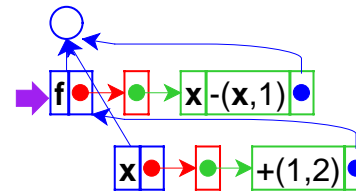
### Evaluation with Lazy Arguments



lookup of x...

let f = proc(x)-(x,1)  
in (f +(1,2))

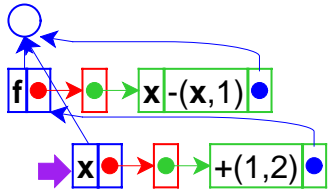
### Evaluation with Lazy Arguments



... forces evaluation of the thunk

let f = proc(x)-(x,1)  
in (f +(1,2))

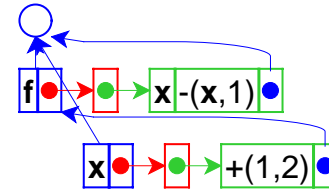
### Evaluation with Lazy Arguments



so 3 is the value of  $x$

```
let f = proc(x)-(x,1)
in (f +(1,2))
```

### Evaluation with Lazy Arguments



The result is 2

```
let f = proc(x)-(x,1)
in (f +(1,2))
```

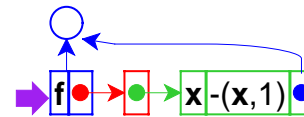
### Evaluation with Lazy Arguments



Lazy expression that needs its environment...

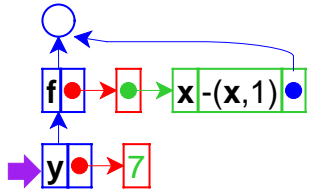
```
let f = proc(x)-(x,1)
in let y = 7
in (f +(1,y))
```

### Evaluation with Lazy Arguments



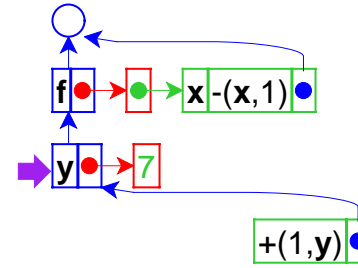
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in let y = 7
in (f +(1,y))
```

### Evaluation with Lazy Arguments



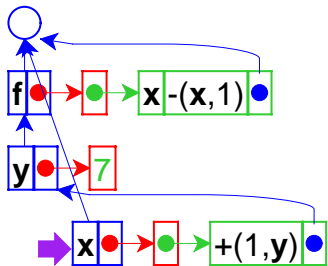
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let f = proc(x)-(x,1)
  in let y = 7
    in (f +(1,y))
```

### Evaluation with Lazy Arguments



```
let f = proc(x)-(x,1)
  in let y = 7
    in (f +(1,y))
```

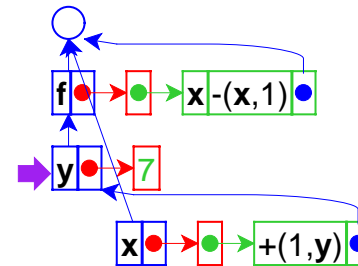
### Evaluation with Lazy Arguments



Evaluation of *x* forces the thunk...

```
let f = proc(x)-(x,1)
  in let y = 7
    in (f +(1,y))
```

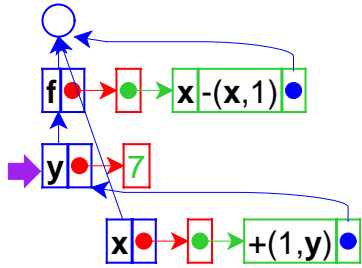
### Evaluation with Lazy Arguments



Triggering evaluation with the *think*'s environment, not the current one

```
let f = proc(x)-(x,1)
  in let y = 7
    in (f +(1,y))
```

## Evaluation with Lazy Arguments



(The result will be 7)

```
let f = proc(x)-(x,1)
in let y = 7
in (f +(1,y))
```

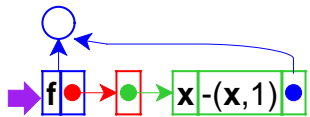
## Evaluation with Lazy Arguments



What if the right-hand side for **y** is an expression, instead of a value?

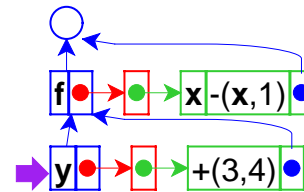
```
let f = proc(x)-(x,1)
in let y = +(3,4)
in (f +(1,y))
```

## Evaluation with Lazy Arguments



```
let f = proc(x)-(x,1)
in let y = +(3,4)
in (f +(1,y))
```

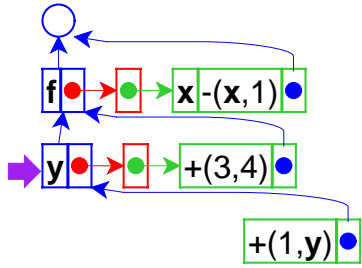
## Evaluation with Lazy Arguments



Added thunk for the value of **y**

```
let f = proc(x)-(x,1)
in let y = +(3,4)
in (f +(1,y))
```

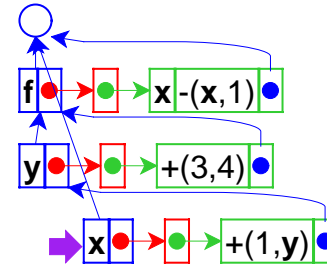
## Evaluation with Lazy Arguments



Another thunk for the argument of f

```
let f = proc(x)-(x,1)
  in let y = +(3,4)
    in (f +(1,y))
```

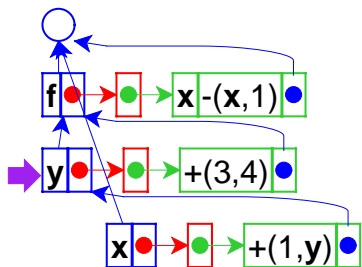
## Evaluation with Lazy Arguments



Evaluation of x forces a thunk...

```
let f = proc(x)-(x,1)
  in let y = +(3,4)
    in (f +(1,y))
```

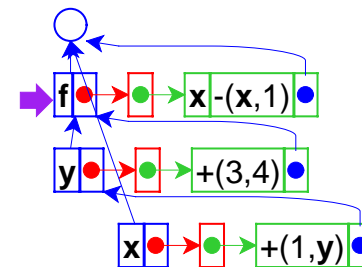
## Evaluation with Lazy Arguments



which, in turn, forces another thunk...

```
let f = proc(x)-(x,1)
  in let y = +(3,4)
    in (f +(1,y))
```

## Evaluation with Lazy Arguments



and so on (to get 7)

```
let f = proc(x)-(x,1)
  in let y = +(3,4)
    in (f +(1,y))
```

## Implementing Lazy Evaluation

Interpreter changes:

- Change `eval-fun-rands` to create thunks
- Change variable lookup to force thunk evaluation

(Implement in DrScheme)

## Call-by-Name and Call-by-Need

The lazy strategy we just implemented is *call-by-name*

- Advantage: unneeded arguments are not computed
- Disadvantage: needed arguments may be computed many times

```
let f = proc(x)+(x,+(x,x))
in (f +(1,+(2,+(3,+(4,+(5,6))))))
```

Best of both worlds: *call-by-need*

- Evaluates each lazy expression once, then remembers the result

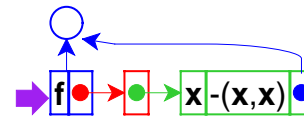
## Evaluation with Lazy Arguments



Start as before...

```
let f = proc(x)-(x,x)
in (f +(1,2))
```

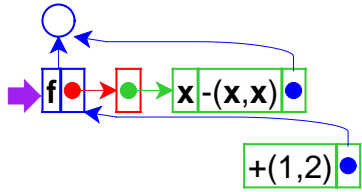
## Evaluation with Lazy Arguments



```
let f = proc(x)-(x,x)
in (f +(1,2))
```

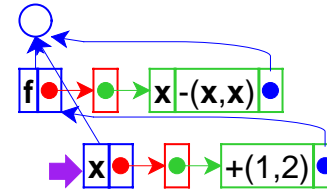


### Evaluation with Lazy Arguments



let f = proc(x)-(x,x)  
in (f +(1,2))

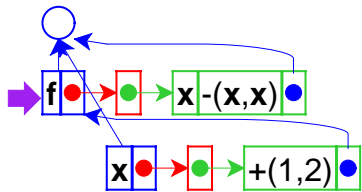
### Evaluation with Lazy Arguments



let f = proc(x)-(x,x)  
in (f +(1,2))

lookup of x...

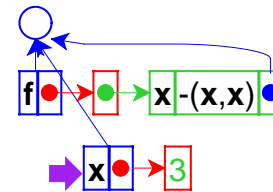
### Evaluation with Lazy Arguments



let f = proc(x)-(x,x)  
in (f +(1,2))

... forces evaluation of the thunk  
to get 3

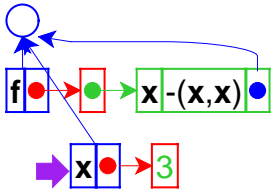
### Evaluation with Lazy Arguments



let f = proc(x)-(x,x)  
in (f +(1,2))

so change x to 3 --- which is the  
essence of call-by-need

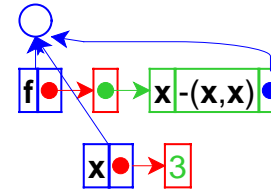
## Evaluation with Lazy Arguments



lookup of **x** again gets 3

```
let f = proc(x)-(x,x)
in (f +(1,2))
```

## Evaluation with Lazy Arguments



(The result is 0)

```
let f = proc(x)-(x,1)
in (f +(1,2))
```

## Implementing Call-by-Need

Interpreter changes:

- Change variable lookup to replace thunks in locations with their values

(Implement in DrScheme)

## Calling Convention Terminology

- Call-by-name and call-by-need = *lazy* evaluation
- Call-by-value = *eager* evaluation

Call-by-reference can augment either

## Popular Calling-Convention Choices

- Most languages are call-by-value
  - C, C++, Pascal, Scheme, Java, ML, Smalltalk...
- Some provide call-by-reference
  - C++, Pascal
- A few are call-by-need
  - Haskell
- Practically no languages are call-by-name

## Popularity of Laziness

Why don't more languages provide lazy evaluation?

- Disadvantage: evaluation order is not obvious

```
let x = 0 f = ...
in let y = set x=1
    z = set x=2
    in { (f y z) ; x }
```

## Popularity of Laziness

Why do some languages provide lazy evaluation?

- Evaluation order does not matter if the language has no **set** form
- Such languages are called *purely functional*
  - Note: call-by-reference is meaningless in a purely functional language
- A language with **set** can be called *imperative*

## Laziness and Eagerness

Even in a purely functional language, lazy and eager evaluation can produce different results

```
let f = proc(x)0
in (f [loop forever])
```

- Eager answer: none
- Lazy answer: 0