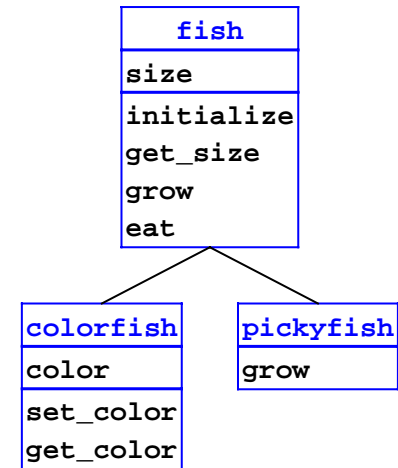


## Outline

- More optimizations for our interpreter
- Types for objects

## Optimization

**Eliminate tree walks:** object creation, method calls



## Object Creation

Current interpreter:

1. Find class
2. Get field list (walk tree)
3. Allocate field array and object

To eliminate tree walks:

2. Extract flat field list from class

## Method Calls

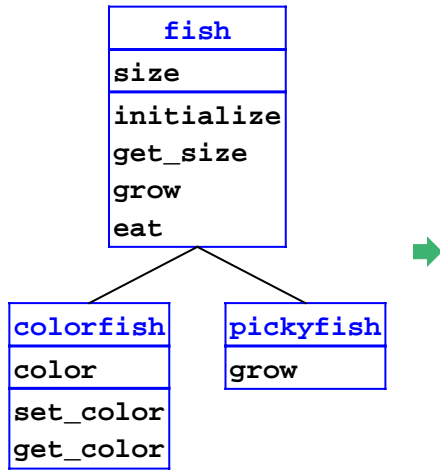
After object and arguments are determined:

1. Lookup object class
2. Find class containing method (walk tree)
3. Get variables for class (walk tree)
4. Create environment: fields + `%super` + `self` + args
5. Evaluate method body

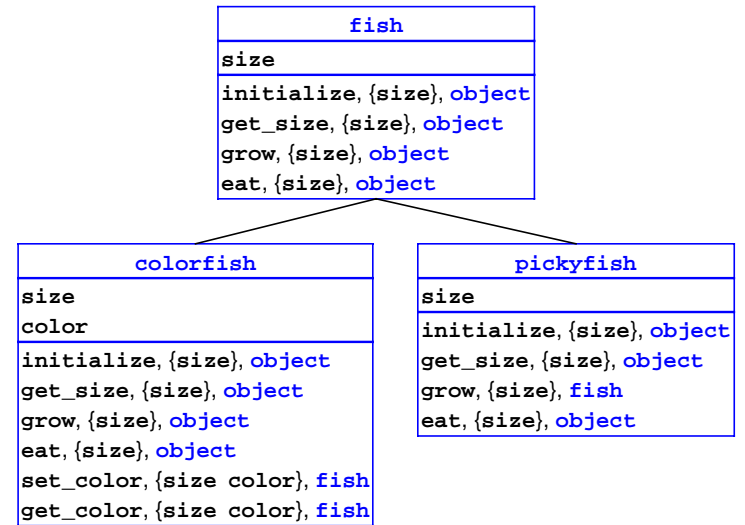
To eliminate tree walks:

- 2 & 3. Find method in current class, extract variable list

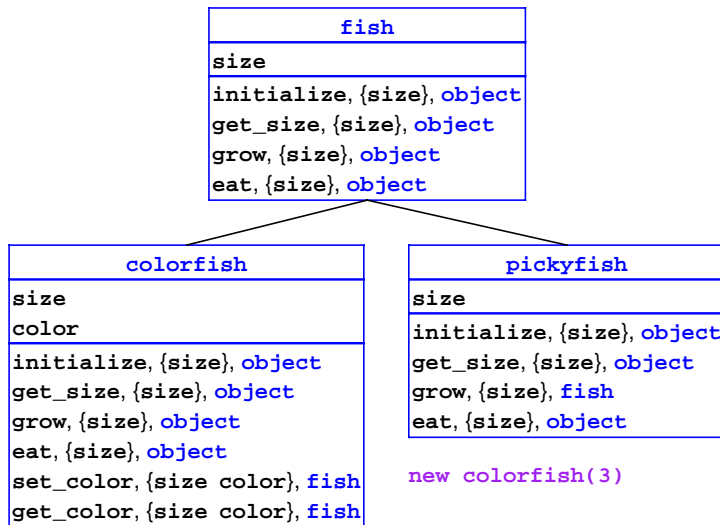
### Class Elaboration



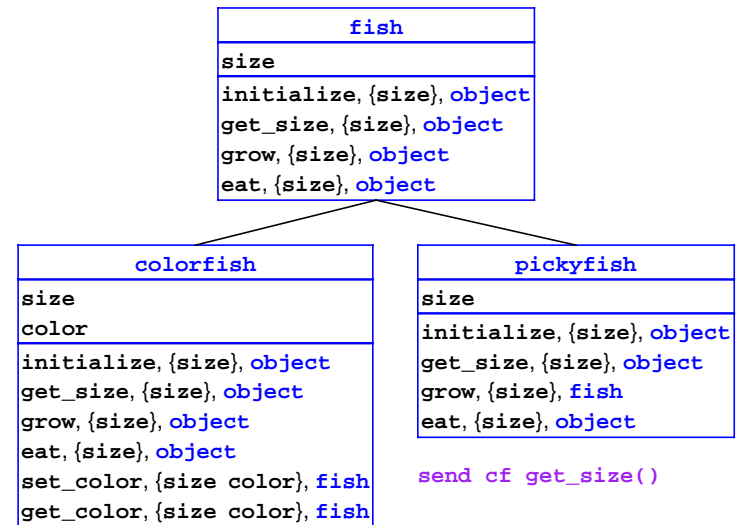
### Class Elaboration



### Class Elaboration



### Class Elaboration



## Implementation

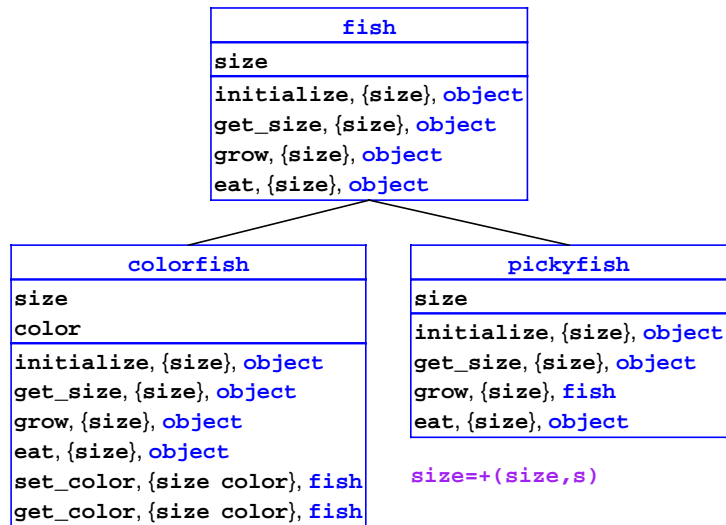
See the book and web page:

- Change `elaborate-class-decls!` to build annotated tree
- Change `new-object` to use class's immediate field list
- Change `apply-method` to work with annotated methods

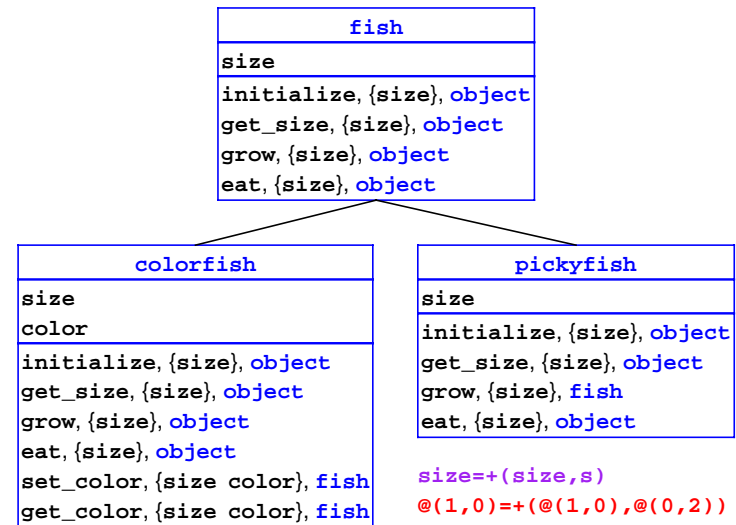
## More Optimization

- Still have list walks: variable lookup, method lookup
  - Can eliminate many with lexical addresses
  - Can eliminate some by pre-computing method positions
  - Need type information to eliminate others

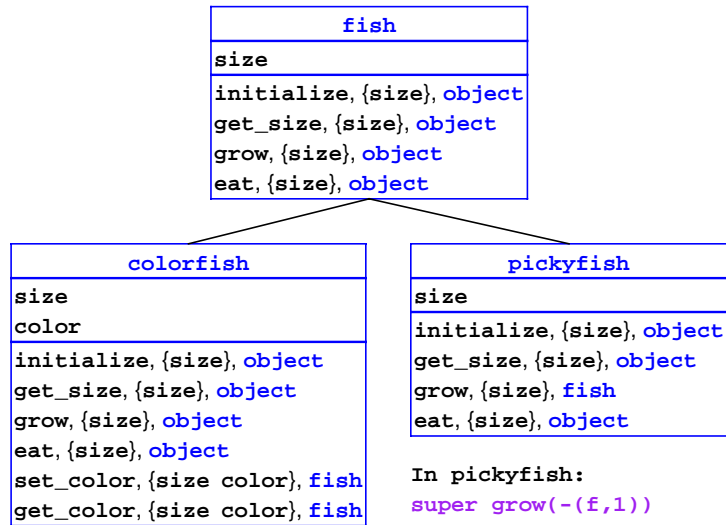
### More Optimization: List Walks



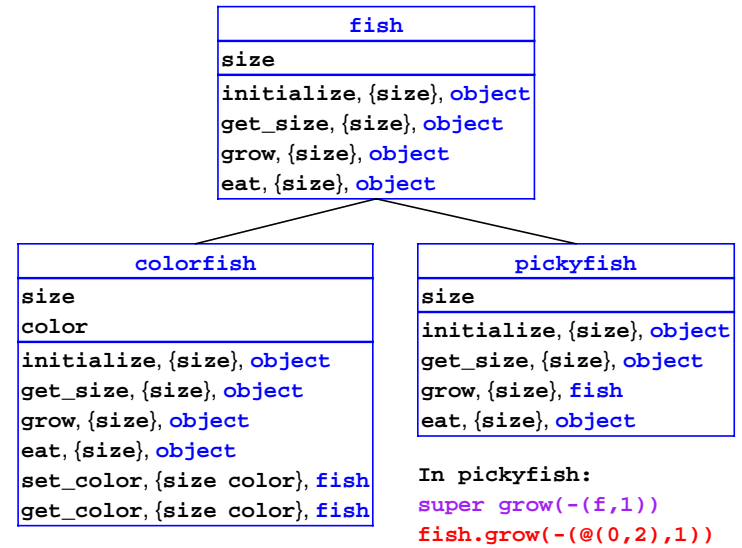
### More Optimization: List Walks



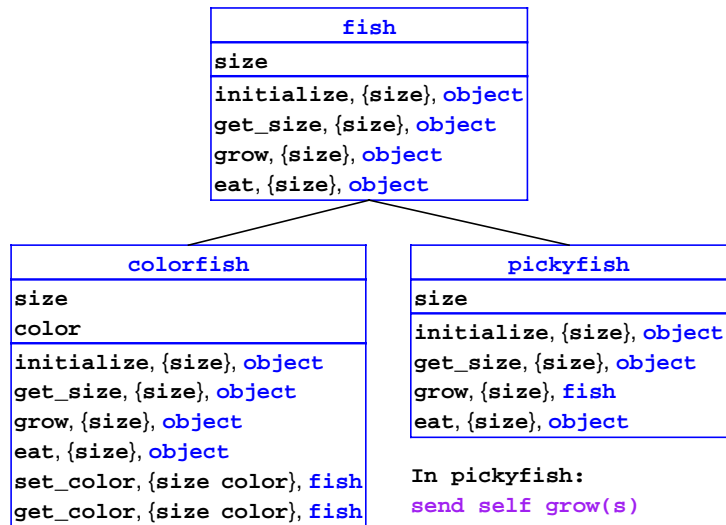
## More Optimization: List Walks



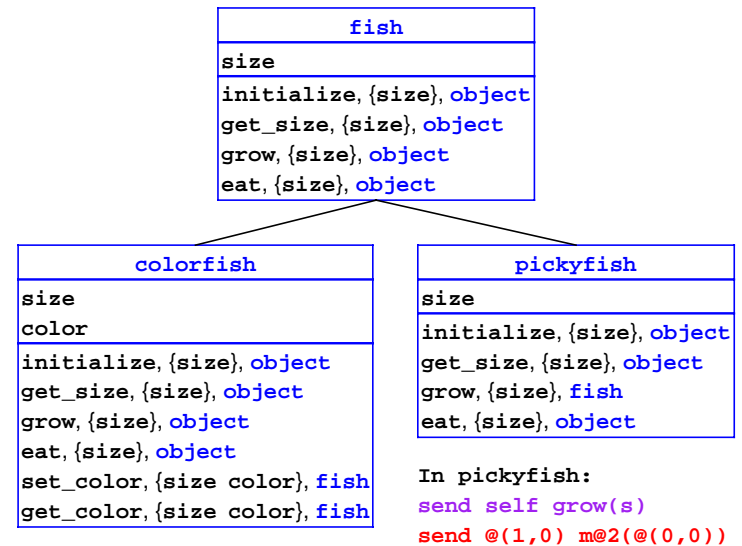
## More Optimization: List Walks



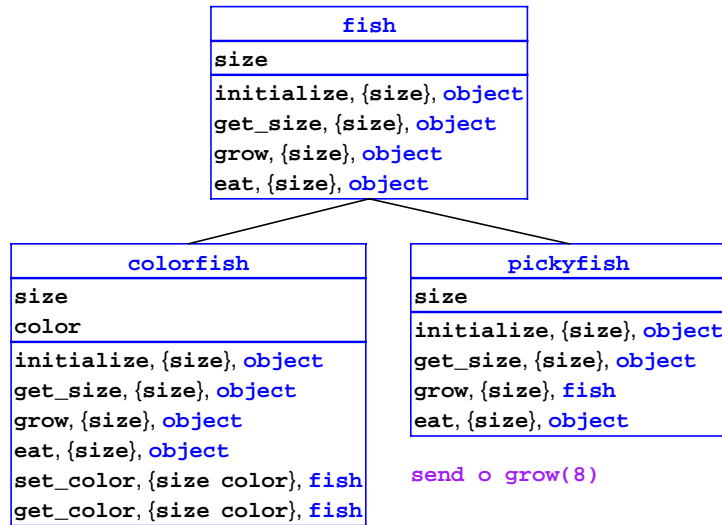
## More Optimization: List Walks



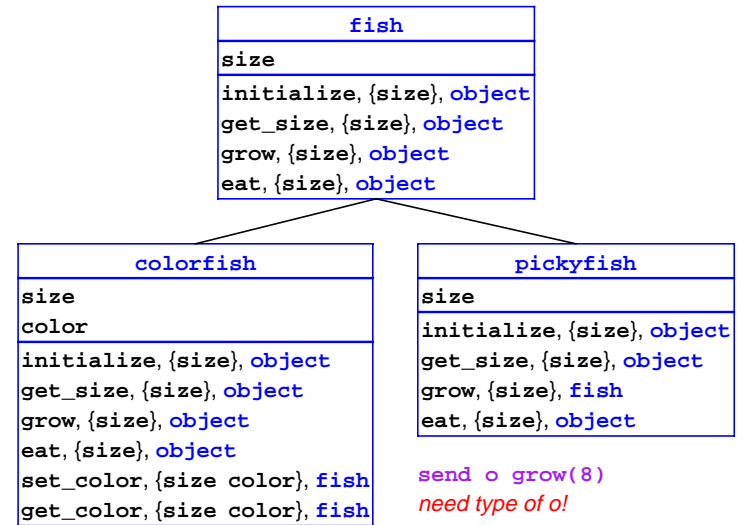
## More Optimization: List Walks



## More Optimization: List Walks



## More Optimization: List Walks



## Object Types

new c1()  
|  
c1

... *if* c1 has an initialize method that takes no arguments

```
class c1 extends ...
  method void initialize() ...
```

## Object Types

new c1(5)  
|  
c1  
|  
int

... *if* c1 has an initialize method that takes one integer

```
class c1 extends ...
  method void initialize(int v) ...
```

## Object Types

```
send new c1() m(false)
      |   |
      c1  bool
      |
      int
```

... if c1 has an m method that takes bool and returns int

```
class c1 extends ...
  method void initialize() ...
  method int m(bool v) ...
```

## Object Types

```
class fish extends object
  field int size
  method void initialize (int s) ...
  method void eat(fish other) ...
class colorfish extends fish
  ...
```

```
send new fish(8) eat(new colorfish(1))
      |           |
      fish       colorfish
      colorfish doesn't match fish
```

## Subtyping

- **Subtype:** An instance of class C can be used as an instance of class C' if C is derived from C'

$C <: C'$

- Subtype rule:

If  $e : T$  and  $T <: T'$ , then  $e : T'$

## Object Types

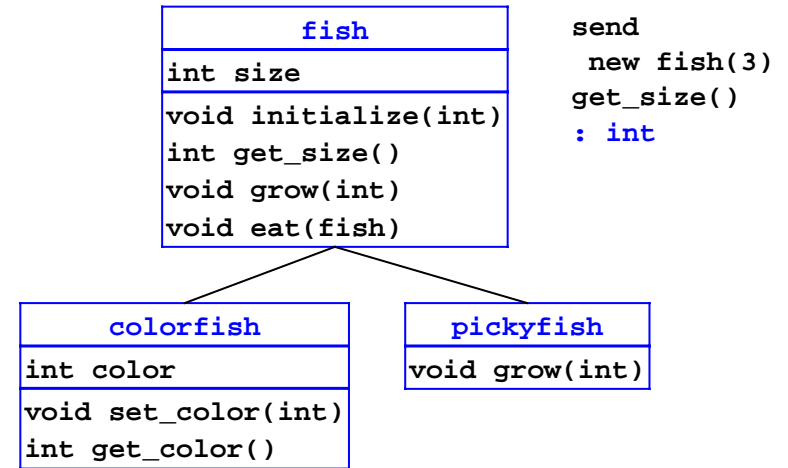
```
class fish extends object
  field int size
  method void initialize (int s) ...
  method void eat(fish other) ...
class colorfish extends fish
  ...
```

```
send new fish(8) eat(new colorfish(1))
      |           |
      fish       colorfish <: fish
                  |
                  void
```

## Language Changes

- Add types to field declarations
- Add types to method arguments and result
- Add `abstract class` and `abstractmethod`
- Add `instanceof`
- Add `cast`

## Program Checking



## Things to Check

`cast` and `instanceof`:

- Operand has an object type (for any class)
- Target class exists

`cast o c1`      `instanceof o c7`

## Things to Check

`cast` and `instanceof`:

- Operand has an object type (for any class)
- Target class exists

`cast` only:

- Class for operand and target are comparable
  - Otherwise, cast cannot possibly succeed

```
class c1 extends object ...
class c2 extends object ...
cast new c1() c2
```

## Things to Check

Object creation:

- Class exists, and is not abstract
- Class has an `initialize` method
- `initialize`'s argument types match the operand types

```
class c1 extends object
  method void initialize(int x, bool y)
  ...

new c1(1, false)
```

## Things to Check

Method calls:

- Receiver expression is an object
- Method is in the object-type's class
  - Except `initialize`...
- Method's argument types match the operand types

```
class c1 extends object
  method void initialize() ...
  method void m(int x, bool y)
  ...

let o1 = new c1()
in send o1 m(1, false)
```

## Things to Check

`super` calls:

- Expression is within a method
- Method is in the superclass, and not abstract
- Method's argument types match the operand types

```
class c1 extends object
  method void m(int x, bool y)
  ...

class c2 extends c1
  method void n()
    super m(1, false)
  ...
```

## Things to Check

`class` declarations:

- Superclass exists, and no cyclic inheritance
- Methods bodies ok
  - Use host class for type of `self`
- Overriding method signatures the same as in superclass
  - Except for `initialize`

```
class c2 extends c1
  method void m(int x, bool y)
    if y then +(2, x) else send self w()
```



## The Initialize Method

```
class c1 extends obj
  field int x
  method void initialize()
    set x = 3
  method int m()
    send self initialize()

class c2 extends c1
  field int y
  method void initialize(int v)
    set y = v
    super initialize()
  ...
```

- Derived class needs different signature for `initialize`

## The Initialize Method

```
class c1 extends obj
  field int x
  method void initialize()
    set x = 3
  method int m()
    send self initialize()

class c2 extends c1
  field int y
  method void initialize(int v)
    set y = v
    super initialize()
  ...
```

- Disallow `send` to `initialize`

## The Initialize Method

```
class c1 extends obj
  field int x
  method void initialize()
    set x = 3
  method int m()
    send self initialize()

class c2 extends c1
  field int y
  method void initialize(int v)
    set y = v
    super initialize()
  ...
```

- `super` call to `initialize` is ok

## Field Initializations

**Not** checked: field initializations

```
class interior_node extends tree
  field tree left
  field tree right
  method void initialize(tree l, tree r)
    begin
      send left sum();
      ...
    end
```

- Can get "bad object 0 for method call"
- This is analogous to the `null` error in Java

## Type Checking and Errors

### Disallowed errors:

- Object has no such method, or Super method not found
- Can't call method of non-object, non-0
- No such field, no such variable
- Illegal primitive argument (except car of empty)

### Allowed errors:

- Can't call method of 0
- Cast failed
- Car of empty

## Implementation

See the book and web page

## Mixing Subtyping and Procedures

Our language still has procedures:

```
let feed = proc(colorfish f)
            send f grow(10)
    o1 = new colorfish(0)
in
    (feed o1)
```

## Mixing Subtyping and Procedures

And higher-order procedures:

```
let feed = proc(colorfish f)
            send f grow(10)
    o1 = new colorfish(0)
    o2 = new colorfish(1)
in let toboth = proc((colorfish -> void) p)
    begin
        (p o1);
        (p o2)
    end
in (toboth feed)
```

## Mixing Subtyping and Procedures

Subtyping on procedure arguments:

```
let feed = proc(fish f)
  send f grow(10)
  o1 = new colorfish(0)
in
  (feed o1)
```

- This works, and is allowed by our subtyping rule

## Mixing Subtyping and Procedures

Subtyping on procedure arguments:

```
let feed = proc(fish f)
  send f grow(10)
  o1 = new colorfish(0)
  o2 = new colorfish(1)
in let toboth = proc((colorfish -> void) p)
  begin
    (p o1);
    (p o2)
  end
in (toboth feed)
```

- This works, but is *not* allowed by our subtyping rule

(fish -> void) versus (colorfish -> void)

## Procedure Subtyping Rule

If  $T1 <: T1'$  and  $T2 <: T2'$   
then  $(T1' \rightarrow T2) <: (T1 \rightarrow T2')$

Another example:

- `dog <: animal`
  - a dog can go anywhere an animal can go
- `(animal -> hairstyle) <: (dog -> hairstyle)`
  - a groomer for all animals can groom a dog
  - a groomer who only works with dogs doesn't work for all animals

## Procedure Subtyping Rule

If  $T1 <: T1'$  and  $T2 <: T2'$   
then  $(T1' \rightarrow T2) <: (T1 \rightarrow T2')$

General intuition:

- $T1 <: T1'$  means  $T1'$  is more general than  $T1$



- A function that is willing to accept a more general argument is itself more specific



## Procedure Subtyping Rule

If  $T1 <: T1'$  and  $T2 <: T2'$   
then  $(T1' \rightarrow T2) <: (T1 \rightarrow T2')$

- Procedure types are **contravariant** with respect to their argument types
- Procedure types are **covariant** with respect to their result types