

Quiz

- Question #1: What is the value of the following expression?

$$+(1,1)$$

- Wrong answer: 0.
- Wrong answer: 42.
- Answer: 2.

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- **Question #2:** What is the value of the following expression?

`+ proc 8`

- **Wrong answer:** error.
- **Answer:** Trick question! `+ proc 8` is not an expression.

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- Question #3: Is the following an expression?

`add1(1, 7)`

- Wrong answer: **No**.
- Answer: **Yes** (according to our grammar).

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- **Question #4:** What is the value of the following expression?

`add1(1, 7)`

- **Answer:** 2 (according to our interpreter).
- But no *real* language (e.g., C++) would accept `add1(1, 7)`.
- Let's agree to call `add1(1, 7)` an ***ill-formed expression*** because `add1` should be used with only one argument.
- Let's agree to never evaluate ill-formed expressions.

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- **Question #5:** What is the value of the following expression?

`add1(1, 7)`

- **Answer:** **None** - the expression is ill-formed.

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- Question #6: Is the following a well-formed expression?

`+ (proc (x) x, 5)`

- Answer: Yes.

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- **Question #7:** What is the value of the following expression?

`+ (proc (x) x, 5)`

- **Answer: None** - it produces an error:

`+`: expects type `<number>` as 1st argument,
given: (closure ((cbv-var x)) (var-exp x)
(empty-env-record)); other arguments were: 5

- Let's agree that a `proc` expression cannot be inside a `+` form.

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- Question #8: Is the following a well-formed expression?

`+ (proc (x) x, 5)`

- Answer: **No.**

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- **Question #9:** Is the following a well-formed expression?

`+((proc(x)x 7), 5)`

- **Answer:** Depends on what we meant by *inside* in our most recent agreement.
 - *Anywhere inside* - **No.**
 - *Immediately inside* - **Yes.**
- Since our interpreter produces **12**, and since that result makes sense, let's agree on *immediately inside*.

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- Question #10: Is the following a well-formed expression?

`+((proc(x)x true), 5)`

- Answer: **Yes**, but we don't want it to be!

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- **Question #11:** Is it possible to define *well-formed* (as a decidable property) so that we reject all expressions that produce errors?
- **Answer: Yes**, obviously: reject *all* expressions!

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- **Question #12:** Is it possible to define *well-formed* (as a decidable property) so that we reject *only* expressions that produce errors?

- **Answer: No.**

`+(1, if ... then 1 else proc(x)x)`

- If we always knew whether ... produces true or false, we could solve the halting problem.

Types

- Solution to our dilemma
 - In the process of rejecting expressions that are certainly bad, also reject some expressions that are good.

```
+ (1, if (prime? 131101) then 1 else proc(x)x)
```

- Overall strategy:
 - Assign a ***type*** to each expression.
 - Compute the type of a complex expression based on the types of its subexpressions.

Types

`1 : int`

`true : bool`

`+(1, 2)`
┌───┴───┐
`int` `int`
│
`int`

`+(1, false)`
┌───┴───┐
`int` `bool`
│
`no type`

Types

if true then 1 else 2
bool int int
 int

if +(1,2) then 1 else 2
 int
 no type

if false then 2 else false
bool int bool
 no type

Types

x : *no type*

$\frac{\text{proc}(\text{bool } x)x}{\text{bool} \rightarrow \text{bool}}$

$\frac{\text{proc}(\text{bool } x)\text{if } x \text{ then } 1 \text{ else } 2}{\text{bool} \rightarrow \text{int}}$

Types

(proc(bool x)if x then 1 else 2 true)

bool -> int bool
 |
 int

(proc(bool x)if x then 1 else 2 5)

bool -> int int
 |
 no type

(7 5)

int int
 |
 no type

Types

$$\frac{\text{proc}(\text{int } x, \text{int } y) + (x, y)}{\text{int} \quad \text{int} \quad \text{int}}$$

$\text{int} * \text{int} \rightarrow \text{int}$

$$\frac{(\text{proc}(\text{int } x, \text{int } y) + (x, y)) \quad 5 \quad 6}{\text{int} * \text{int} \rightarrow \text{int} \quad \text{int} \quad \text{int}}$$

int

$$\frac{(\text{proc}(\text{int } x, \text{int } y) + (x, y)) \quad 5}{\text{int} * \text{int} \rightarrow \text{int} \quad \text{int}}$$

no type

New Interpreter and Checker

- Change our interpreter:
 - Add types for arguments and letrec results to the grammar
- Implement a type-checker:
 - Recursively assign types to subexpressions
 - Check consistency at `if` and application
 - Treat primitives as built-in functions

`+ : int * int -> int`