

Data Mining Seminar : Sampling
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What Properties are Maintained by Random Sampling
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What properties can be recovered from random sampling? What cannot?
- if data is from much bigger distribution, only the first type interesting !

==== density based estimates ====

- + what fraction of points satisfy this property?
- + do more than X fraction of points satisfy this property?
- + what objects have high frequency?

==== shape estimates ====

- + what is most extreme point?
- + score of k-means clustering?

Let P be the phenomenon we are sampling from.
Q subset P is sample.

R = class of subsets of P "ranges"
~ geometric ranges (balls, rectangles, half-spaces) | intervals
~ kernels (weighted subsets)
~ "dual" P is rectangles and r in R is point "stabbing"
~ simple combination of these ranges

---- density -----

for r in R let
 $r(P) = | r \cap P | / | P |$
be fraction of objects from P in range r
"density"

want property:

** for any range r in R **
 $| r(P) - r(Q) | < \text{eps}$
for some parameter eps in $[0,1]$ (think of $\text{eps} = 0.01 = 1/100$)

Q1: Can we do this?

Q2: How big does Q need to be?

A1: Yes, if R is "reasonable"

basically if $|P| = n$, finite, then $|R| < n^v$
bounded VC-dimension (basis of learning theory)

balls $v = d+1$
axis-aligned rectangles $v = 2d$
half-spaces $v = d+1$
 $P =$ rectangles $v = 2d$
 $v \sim$ description complexity
 \sim degrees of freedom

most important:
intervals $v = 2$

A2: $|Q| = (1/\epsilon^2)(v + \log(1/\delta))$
 $\epsilon =$ maximum error
 $\delta =$ probability of failure (success w.p. $> 1-\delta$)
 $v =$ VC-dimension
[Vapnik-Chervonenkis 1971 --> Li, Long, Srinivasan 2001 (via Talagrand 1984)]

eps-sample aka eps-approximation aka agnostic sampling

Chernoff Bound:

r independent events X_1, X_2, \dots, X_r

$A = (1/t) \sum_i X_i$

X_i in $[0,1]$

$\Pr[A - E[A] < \epsilon] < 2 \exp(-2 \epsilon^2 t) < \delta$

X_i is 1 if sample i in r , 0 if X_i not in r (contribution to $r(Q)$)

solve for $|Q| = t > (1/2) (1/\epsilon^2) \ln(2/\delta)$

Frequent Objects

Let R_ϵ subset R such that $r(P) > \epsilon$

$R_\epsilon = \{r \text{ in } R \mid r(P) > \epsilon\}$

Want Q such that

** for all r in R_ϵ **

$r(Q) > 0$

we "hit" all large enough ranges.

Q1: Can we do this?

Q2: How large does Q need to be?

A1: Clearly yes if R satisfies above (v is bounded)

if eps-sample

$|r(P) - r(Q)| < \text{eps}$

\rightarrow if $r(P) > \text{eps}$ \rightarrow $r(Q) > 0$

Small discrete sets (only m possible values) also work.

Here also $v = 1$

since at most $n+1$ distinct ranges with different subsets.

A2: $|Q| = (v/\text{eps}) \log (v/\text{eps} * \text{delta})$

eps = maximum error

delta = probability of failure (success w.p. $> 1-\text{delta}$)

v = VC-dimension

[Hausler + Welzl 1986]

eps-net aka heavy-hitters aka noiseless-learning

discrete sets: heavy-hitters are all sets which occur more than $\text{eps} * n$ times.

Note: We might accidentally hit the small sets, or over-sample large sets.

.... proof from Chernoff bound - small sets are easier class

.... similar to Coupon Collectors problem

extreme points:

max value of set. Will sample recover?

Average value. <sometimes, if variance is low / bounded>

k-means cluster. Can you sample Q subset P

run $C = \text{k-means}(Q)$

compare average cost $|P|/|Q| \text{cost}(C,Q) - \text{cost}(C,P) | ?$

No?

Trick: don't try to recover density, since won't work.

Sample C directly:

basic: for max, just choose maximum point.

No need to sample.

approx-convex hull:

Let u be a unit vector in some direction

$\text{wid}(u,P) = \max_{\{p \text{ in } P\}} \langle u,p \rangle - \min_{\{p \text{ in } P\}} \langle u,p \rangle$
 eps-kernel Q:
 for **any** u : $(\text{wid}(u,P) - \text{wid}(u,Q))/\text{wid}(u,P) < \text{eps}$
 (other forms, but this settled upon)

1. normalize so P fits in $[-1,1]^d$
2. place $(1/\text{eps})^{\{(d-1)/2\}}$ points G evenly on each side of $[-2,2]^d$
3. Select to Q the closest point in P to each g in G

or

- 2b. Take one point from each $[\text{eps}]^d$ grid cell

k-means clustering:

let $\text{phi}_C : P \rightarrow C$
 $\text{phi}_C(p) = \text{argmin}_{\{c \text{ in } C\}} \|p - c\|$

construct $C_1 = c_1$ at random ($q \text{ in } P$)
 $C_{\{i+1\}} = C_i \cup c_{\{i+1\}}$
 choose $c_{\{i+1\}}$ proportional to $(\text{phi}_{\{C_i\}}(p))^2$

$\text{cost}(C,P)$ is 8-approximation to optimal centers
 $\text{cost}(C,P) = \sum_{\{p \text{ in } P\}} \text{phi}_C(p)$

more general:

- $\text{phi}(p)$ = sensitivity of p
 - ~ $\text{phi}(p)$ proportional to how $\text{cost}(Q)$ changes to $\text{cost}(Q / p)$
 where Q is random subset
 - ~ "extreme" have higher $\text{phi}(p)$
 - ~ points near many other points have lower $\text{phi}(p)$
- complicated to describe in some specific settings

 Question: MAP Estimate?