

Data Mining Seminar : Sampling
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OVERVIEW

Aug 22, 2012

0. Aug 22 (today)
Overview

0.5 Aug 29
Jeff is at conference, not seminar.

1. Sept 05

What properties can be recovered from random sampling? What cannot?
- if data is from much bigger distribution, only the first type interesting !

==== density based estimates ====

- + what fraction of points satisfy this property?
- + do more than X fraction of points satisfy this property?
- + what objects have high frequency?

==== shape estimates ====

- + what is most extreme point?
- + score of k-means clustering?

2. Sept 12

Importance and Rejection Sampling:
- how to improve sampling with clever weighting.

Importance Sampling:

- Some areas more important than others (specific questions)
- a. select proportionally more points from important areas
 - b. proportionally down-weight points from important areas
- > lower variance in important areas

Rejection Sampling:

- Computational Tool for complex distribution.
- a. Sample from simple distribution close to true distribution.
 - b. Keep those points with some probability (based on relative probabilities).

3. Sept 19 (Yan)

Markov Chains: Definitions

- powerful rejection sampling for really crazy distributions.

X = state space

$P : X \rightarrow X$ = transition probability matrix

q = state

Starting at q_0 in X , move $q_{i+1} = P q_i$. Where is q_{∞} likely to end up?

What are high-level properties or when this is not "strange"

4. Sept 26 (Haibo)

Markov Chains: Rapidly Mixing + Convergence

- how long do we run this thing?

$q_n = P^n q_0$. As $n \rightarrow \infty$, for any start q_0 , the same q_{∞} is reached. How long does this take.

If for **all** states x in X , and any q_0 , if

$|\Pr[q_{\infty} = x] - \Pr[q_n = x]| / |\Pr[q_{\infty} = x] - \Pr[q_{2n} = x]| < 2$
for all $n > n_0$, then (P, X, q_0) is said to be "rapidly mixing".

What does this mean?

Some scenarios...

5. Oct 03 (Parasaran)

Metropolis Algorithm + Gibbs

- how do we use this to sample from weird distributions?

Given "probe-only distribution" μ

- can only pick point x , and evaluate $\phi(x)$
such that $\phi(x) = c \mu(x)$ (c is fixed but unknown)

Then $y = P * x$ is an algorithm:

- choose random q (nearby x perhaps)
if $\phi(q) > \phi(x)$ $\rightarrow y=q$
else (w.p. $\phi(q)/\phi(x)$) $y=q$
otherwise $y=x$ (stays put)

Gibbs, X = multi-dimensional:

\rightarrow Speed up if, we fix subset of dimensions, we can sample from non-fixed set exactly and efficiently (rejection sampling, or direct).

- \rightarrow Fix each set in turn = one-step.

No rejection...

6. Oct 17 (XXXXX)

Advanced Sampling

- what techniques are state of the art for speeding this up?

Main problem is "auto-correlation".

- a) trying to move too far, and not moving
- b) only trying nearby, and not exploring space.

Think of q as random Gaussian kernel around x . What is variance?

- if change dynamically, may not converge!

Have small probability of changing variance in P . (e.g. double, cut-by-half)
(like running many in parallel, at diff variance)

7. Oct 24 (Prasana)

Coupling from the Past

- selecting a "true" random sample

To sample: run many many steps, then take 1 sample, start over again for next.
in practice: run 1000 steps (burn-in), then retain all following samples
(e.g. 5000)

What can we say in theory - can we ever guarantee a truly random sample: Yes!

Pretend we started chains at all q_0 in X . Use same randomness for all.

- If they all reach same spot, the next is random.
(bit more subtle)

continuous state spaces X too.

8. Oct 31 (Reza)

Reservoir Sampling

- one-pass random sample

Streaming model: get to read each data point once. Memory too small to store all.

- Important: (a) network traffic (actually like this), (b) I/O Efficient.

How do we draw a random sample: Keep i th point w.p. k/i (if kept, kick out random point)

How do we do this in parallel, when data is distributed.

... or L_p sampling.

9. Nov 7 (MiaoMiao)

Variance-Optimal Sampling

- best way to sample weighted points

Stream of points, with weights \rightarrow output: weighted subset of points (diff weight)

Want subset, so estimate of sum of weights in any subset preserved.

- select high weight points with prob 1.
- select remaining points proportional to weight.

For any subset query

- right expected value
- minimize variance of estimate

can you trade off zero-bias for less variance?

"regularization"

10. Nov 14 (Suresh)

Precision Sampling

- variation of VarOpt Sampling with negative weights

Maintains set estimates \hat{a}_i with "precision" u_i (error in count)

- needs $1/u_i$ space for each
- averages them to get error bound
- each an improvement over random sample

11. Nov 21 (Samira)

Discrepancy

- get same (density) guarantees, with fewer points.

(P, R) is range space (subsets R of P , e.g. contained in ball)

$\chi: P \rightarrow \{-1, +1\}$

$\text{disc}_\chi(P, R) = \max_{\{r \in R\}} \left| \sum_{\{p \in P \cap r\}} \chi(p) \right|$
(as close to 0 as possible)

Remove points colored -1, now size $|P|/2$.

Repeat on these points for about $\log(|P|)$ rounds

\rightarrow for good χ , better than random sampling on (P, R) .

12. Nov 28 (XXXXX)

Constructing the Coloring

- how to construct this coloring

Hyperbolic Cosine:

$H(r) = \cosh(\sqrt{\ln |R| / |P|} * \text{disc_chi}(P,r))$

$H = \sum_{r \in R} H(r)$

choose point that decreases H most

$\cosh(x) = (e^x + e^{-x})/2$

Edge Walking:

produces "optimal" discrepancy

Special Markov Chain on relaxed coloring $\chi : P \rightarrow [-1,1]$

when point hits $\{-1,1\}$ becomes fixed

13. Dec 05 (XXXXX)

Streaming Weighted points for Ranges

- beats random sampling in stream with structure

Pair-aggregation. Pairs points in similar ranges.

if $p_i + p_j < 1$: $p_i = p_i + p_j$ // $p_j = 0$
So don't sample both. Reduces variance.

(idea should be similar to discrepancy -- often constructs pairing +1,-1)