Honest and Lying Types

Thesis Proposal

Ben Greenman 2019-11-25

Committee:

- 1. Matthias Felleisen
- 2. Amal Ahmed
- 3. Jan Vitek
- 4. Shriram Krishnamurthi
- 5. Fritz Henglein
- 6. Sam Tobin-Hochstadt

Honest and Lying Types

Thesis Proposal

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Thesis Proposal: "Gradual Typing" November 18, 2019

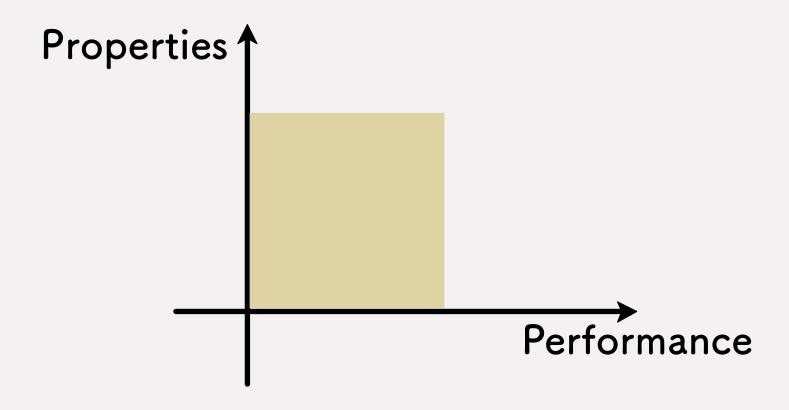


Thesis Proposal: "Gradual Typing" November 25, 2019

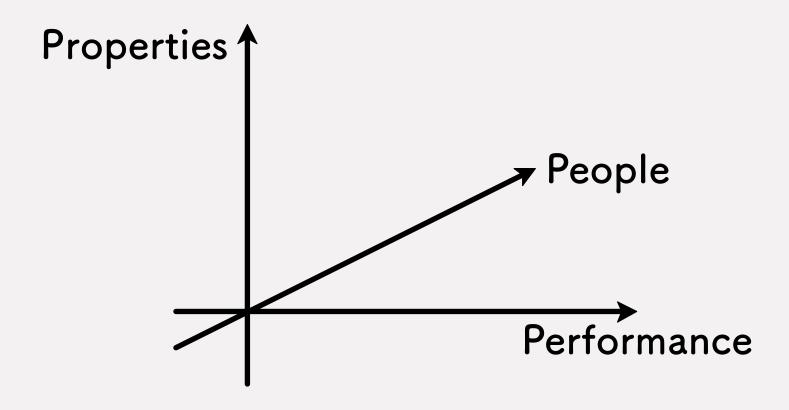
Last Week:



Today:



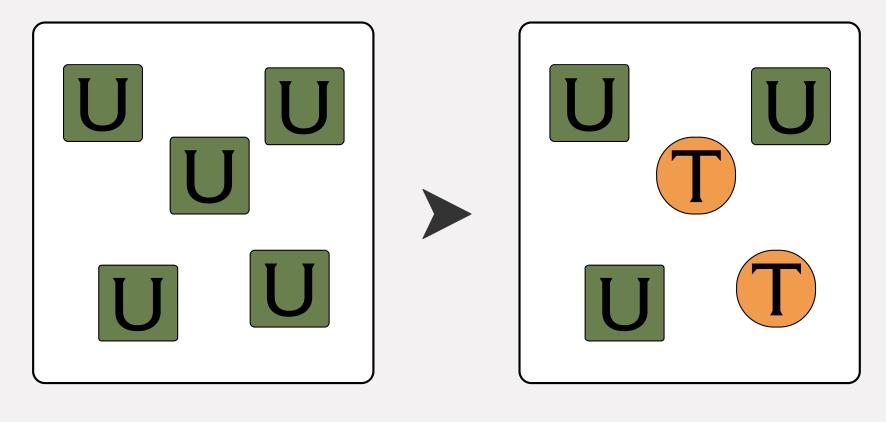
Future:



Migratory Typing

Migratory Typing

Add types to a dynamically-typed language



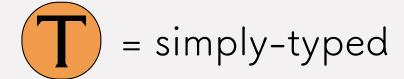
Untyped code

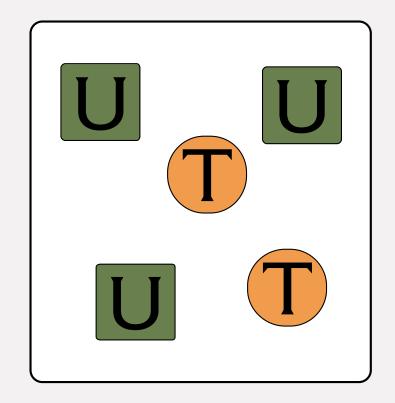
Mixed-Typed code

Migratory Typing

Add types to a dynamically-typed language





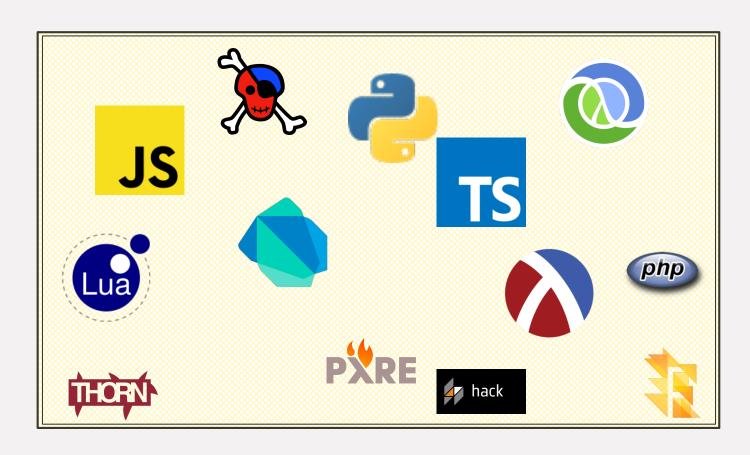


Mixed-Typed code

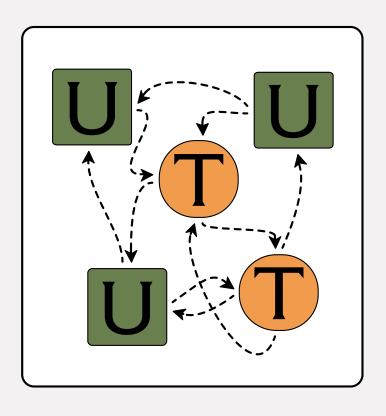
Motivation

Because lots of untyped code exists.

Landscape of Models and Implementations

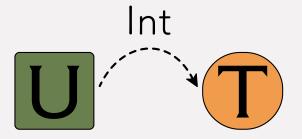


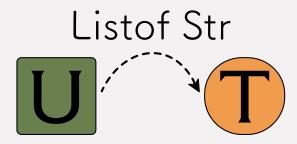
Challenge = Interoperability



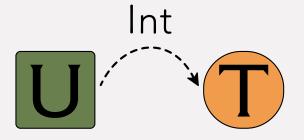
What do types mean when untyped values and typed values interact?

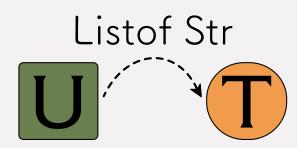
Challenge = Interoperability

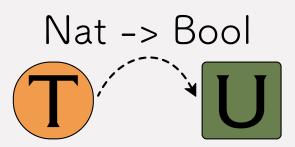


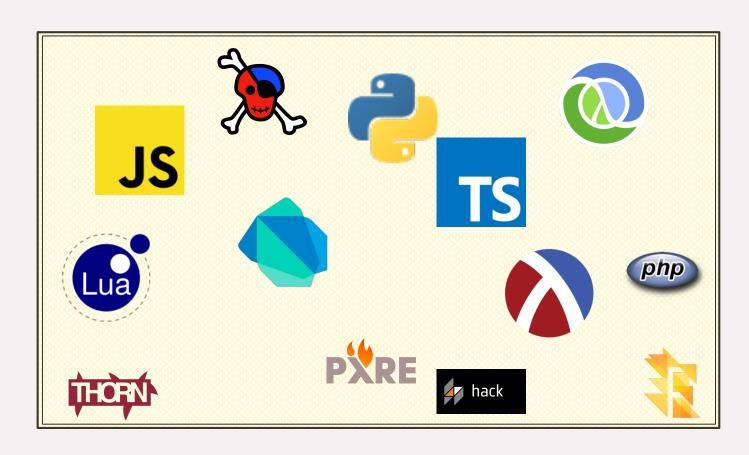


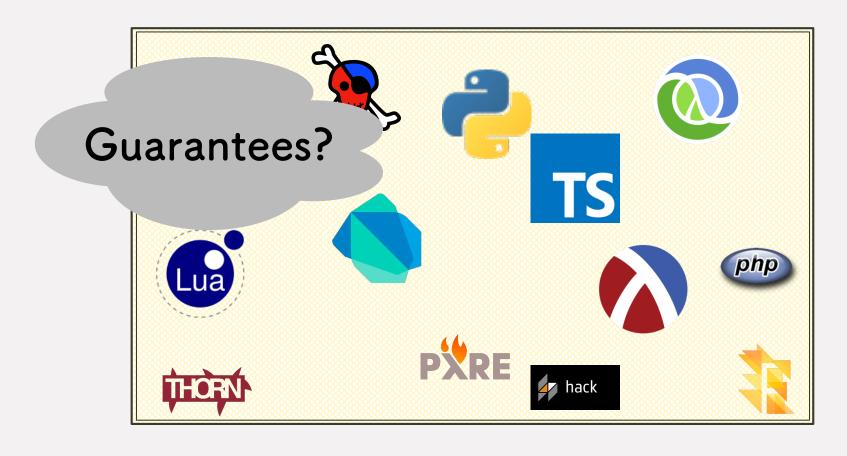
Challenge = Interoperability

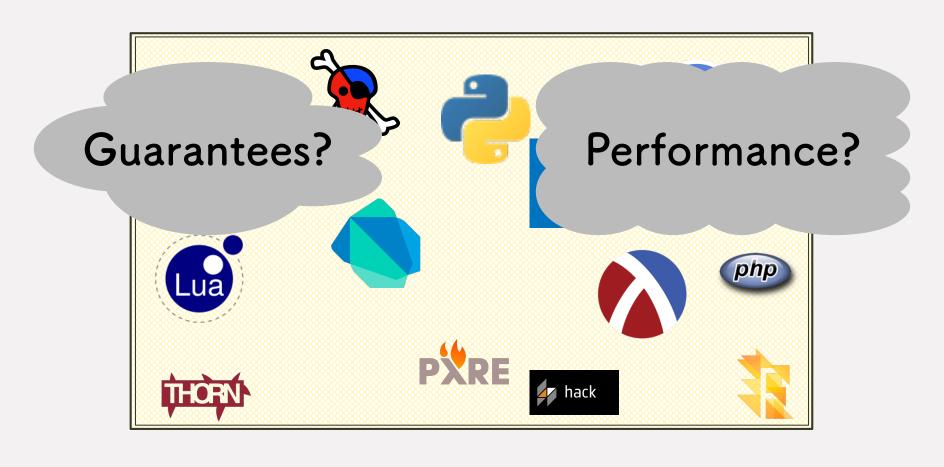


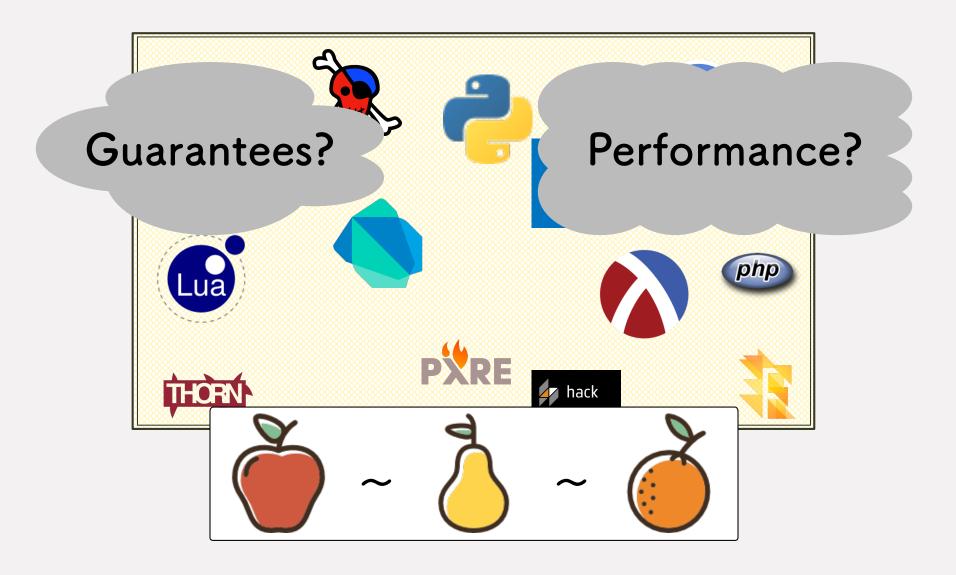












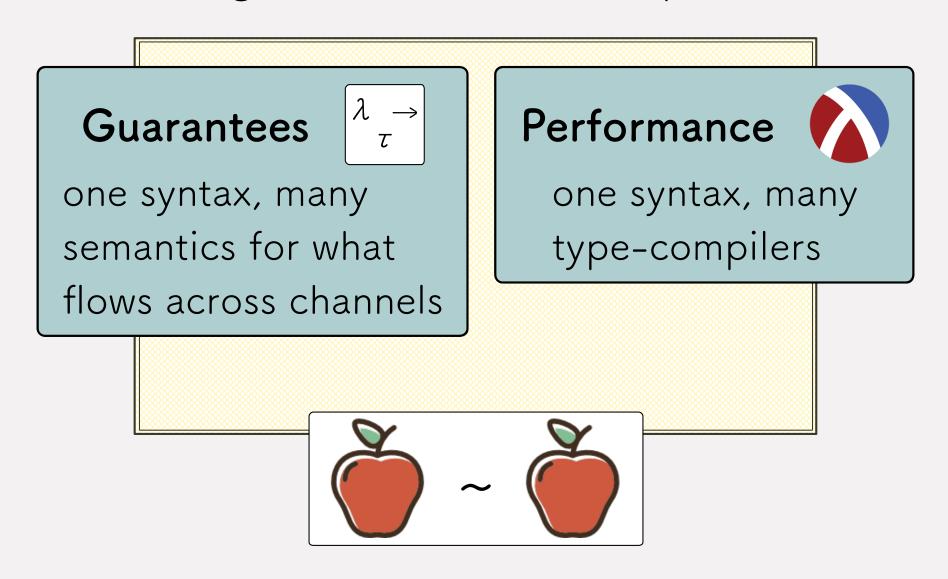
Icons made by Freepik from Flaticon.com

My Research

Research Agenda: Scientific Comparison



Research Agenda: Scientific Comparison



Icons made by Freepik from Flaticon.com

Research Agenda: Results so Far

Design Space Analysis

OOPSLA 19

Ben Greenman, Matthias Felleisen, and Christos Dimoulas

ICFP

18

Ben Greenman and Matthias Felleisen

Research Agenda: Results so Far

Design Space Analysis

OOPSLA 19

Ben Greenman, Matthias Felleisen, and Christos Dimoulas

ICFP

18

Ben Greenman and Matthias Felleisen

Performance Evaluation

JFP 19

Ben Greenman, Asumu Takikawa, Max S. New, Daniel Feltey, Robert Bruce Findler, Jan Vitek, and Matthias Felleisen

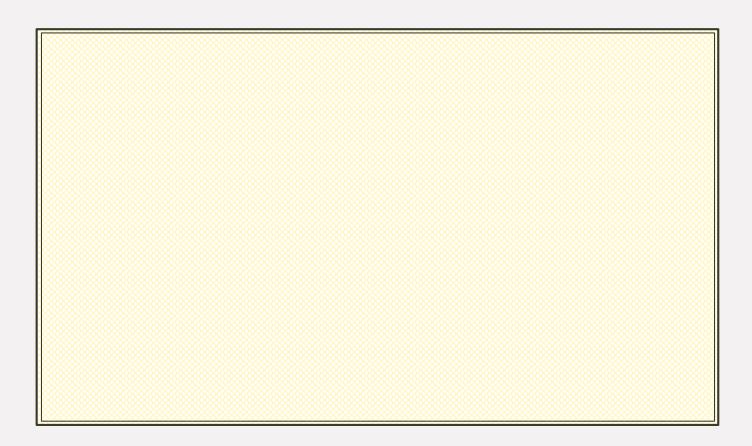
PEPM 18

Ben Greenman and Zeina Migeed

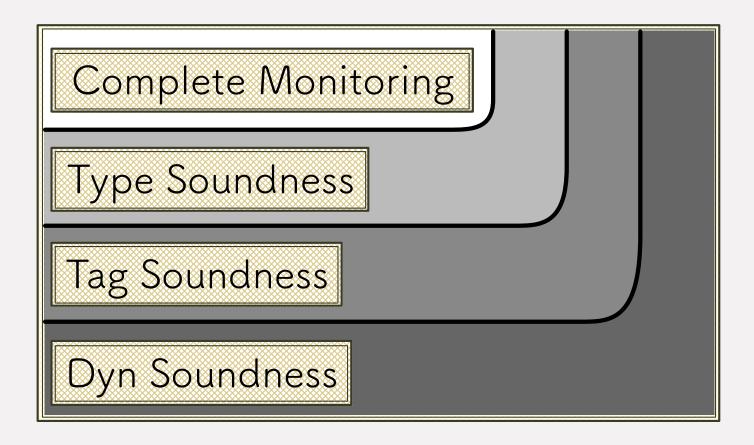
POPL 16

Asumu Takikawa, Daniel Feltey, Ben Greenman, Max S. New, Jan Vitek, and Matthias Felleisen

Landscape: Guarantees



Landscape: Guarantees



(a total spectrum)

types predict behavior

Type Soundness

types predict behavior in typed code, nothing in untyped code

Tag Soundness

types predict shapes in typed code, nothing in untyped code

Dyn Soundness

types predict behavior

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Dyn Soundness

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Honest

Type Soundness

Tag Soundness

Dyn Soundness

Honest

Type Soundness

Lying

Tag Soundness

Dyn Soundness

Honest

Type Soundness

Lying

Tag Soundness

Dyn Soundness

Vacuous

Honest vs. Lying Types

```
Client
 (define path "/tmp/file.txt")
 (define (count acc str)
   (+ 1 acc))
 (t-fold-file path 0 count)
           API
            (provide
            →t-fold-file : (-> Path Num
                                 (-> Num Str Num) Num)))
            (define t-fold-file u-fold-file)
```

```
Client
                                   <u>Library</u>
 (define path "/tmp/file.txt")
                                     (define (u-fold-file path acc f)
                                       ...; read `str` from `path`
                                       ... (f str acc) ...
 (define (count acc str)
   (+ 1 acc))
                                       ...)
 (t-fold-file path 0 count)
           API
            (provide
             →t-fold-file : (-> Path Num
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             (define t-fold-file u-fold-file)
```

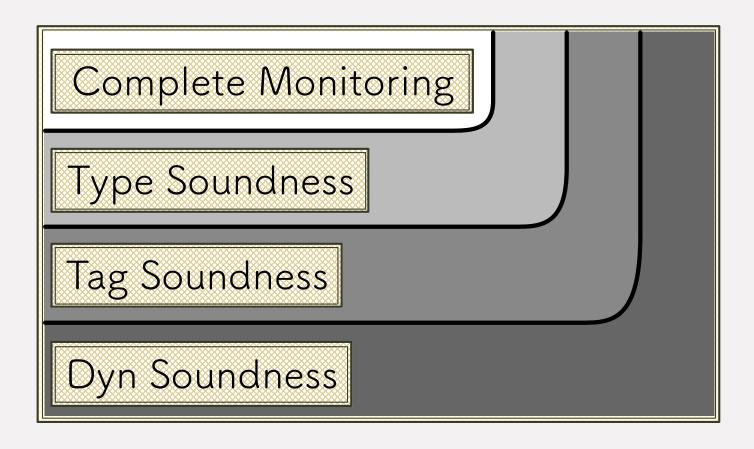
```
Client
                                   <u>Library</u>
 (define path "/tmp/file.txt")
                                    (define (u-fold-file path acc f)
                                      ...; read `str` from `path`
 (define (count acc str)
                                      ... (f str acc) ...
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                                      ...)
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           API
            (provide
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```

```
Client
                                  <u>Library</u>
 (define path "/tmp/file.txt")
                                   (define (u-fold-file path acc f)
                                      ...; read `str` from `path`
 (define (count acc str)

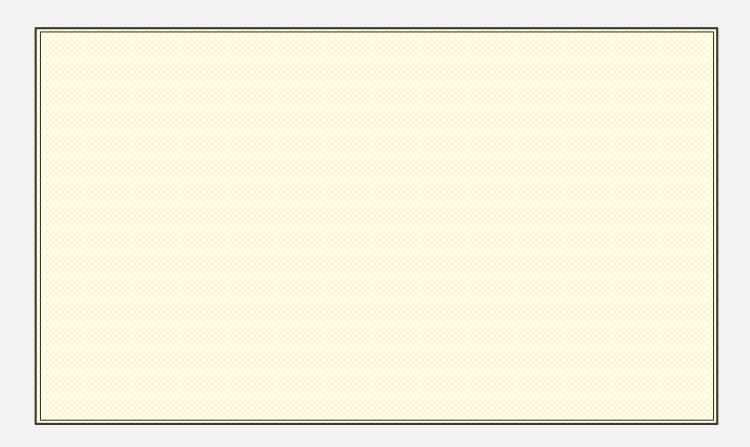
✓
                                      ... (f str acc) ...
   (+ 1 acc))
                                      ...)
          Do the API types protect the Client?
 (t-fold
           API
            (provide
            →t-fold-file : (-> Path Num
                                 (-> Num Str Num) Num)))
            (define t-fold-file u-fold-file)
```

```
Client
                                  <u>Library</u>
 (define path "/tmp/file.txt")
                                   (define (u-fold-file path acc f)
                                     ...; read `str` from `path`
 (define (count acc str)≪
                                     ... (f str acc) ...
   (+ 1 acc))
                                     ...)
          Do the API types protect the Client?
 (t-fold
                                     Lying \Rightarrow yes
           Honest \implies yes
            →t-fold-file : (-> Path Num
                                (-> Num Str Num) Num)))
            (define t-fold-file u-fold-file)
```

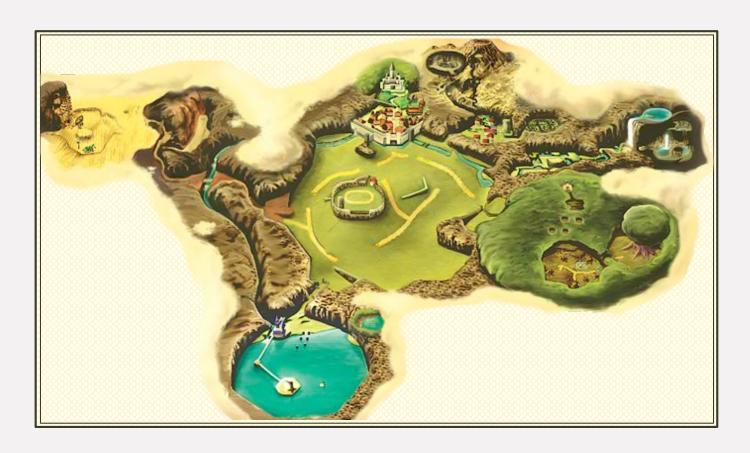
Landscape: Guarantees



Landscape: Performance



Landscape: Performance



Varied space, difficult to rank alternatives

Natural vs. Transient

Natural vs. Transient

Complete Monitoring

guard all boundaries with deep checks

(listof int?)

Tag Soundness

rewrite typed code to tag-check inputs

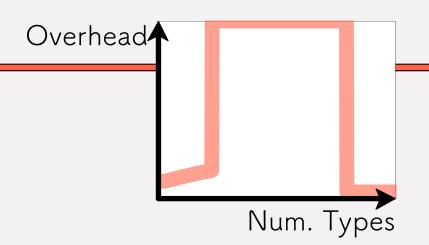
list?

Performance Comparison

ICFP 2018

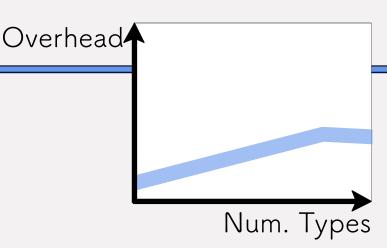
Natural

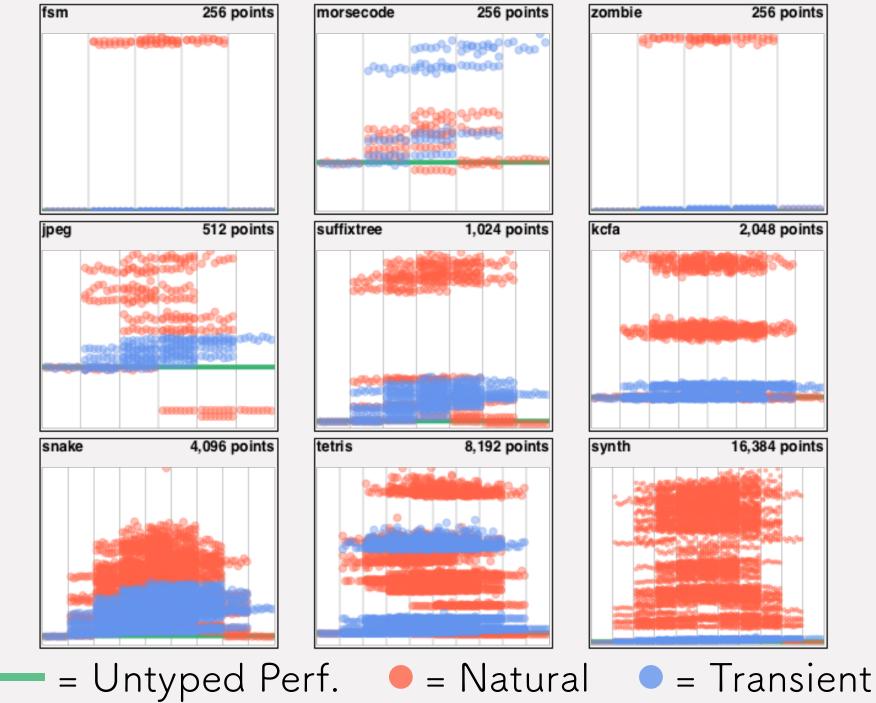
boundaries add "large" overhead



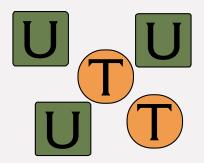
Transient

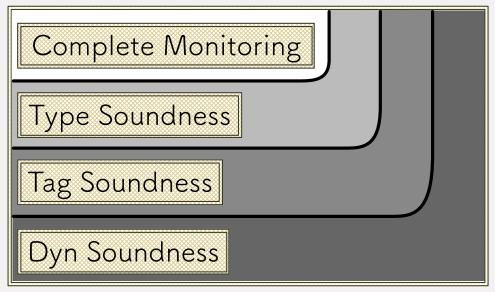
types add "small" overhead

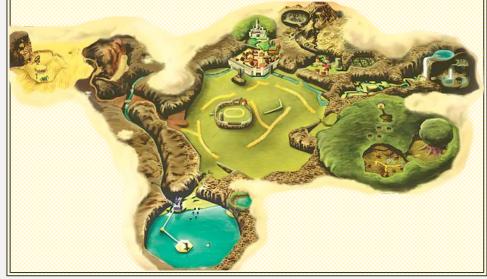




Thesis Question







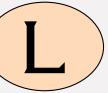
Goal = Migratory Typing Problem = Performance

What to do?

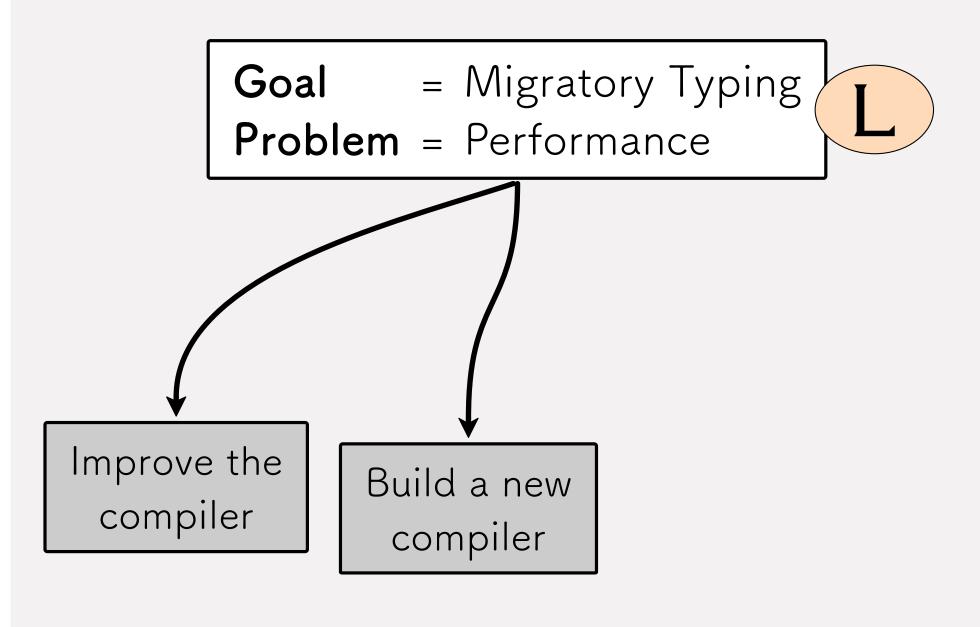
Goal = Migratory Typing
Problem = Performance

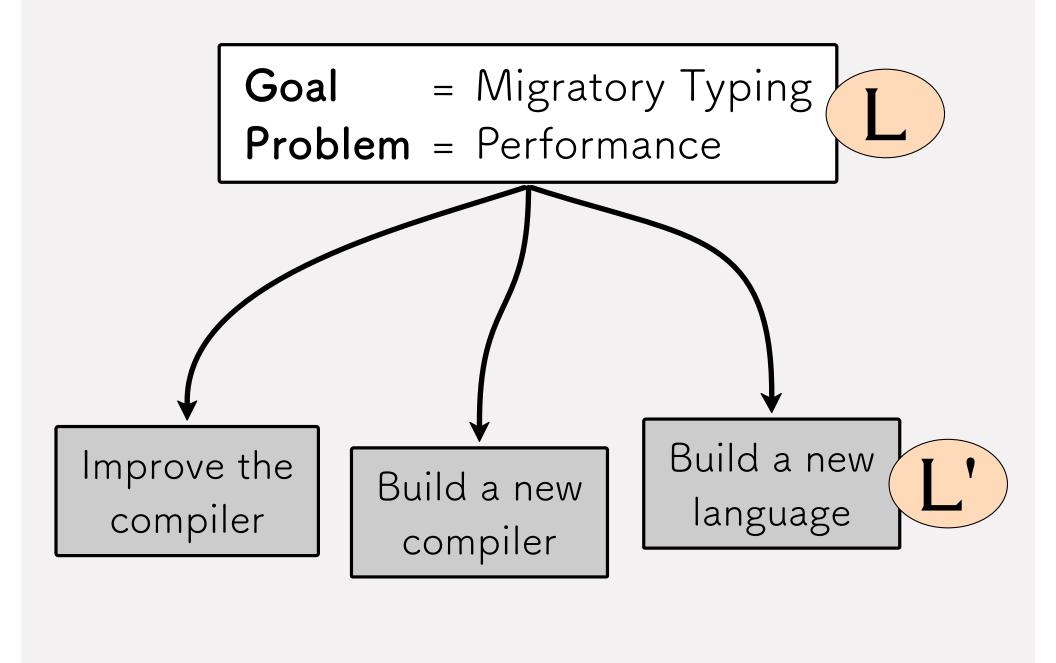


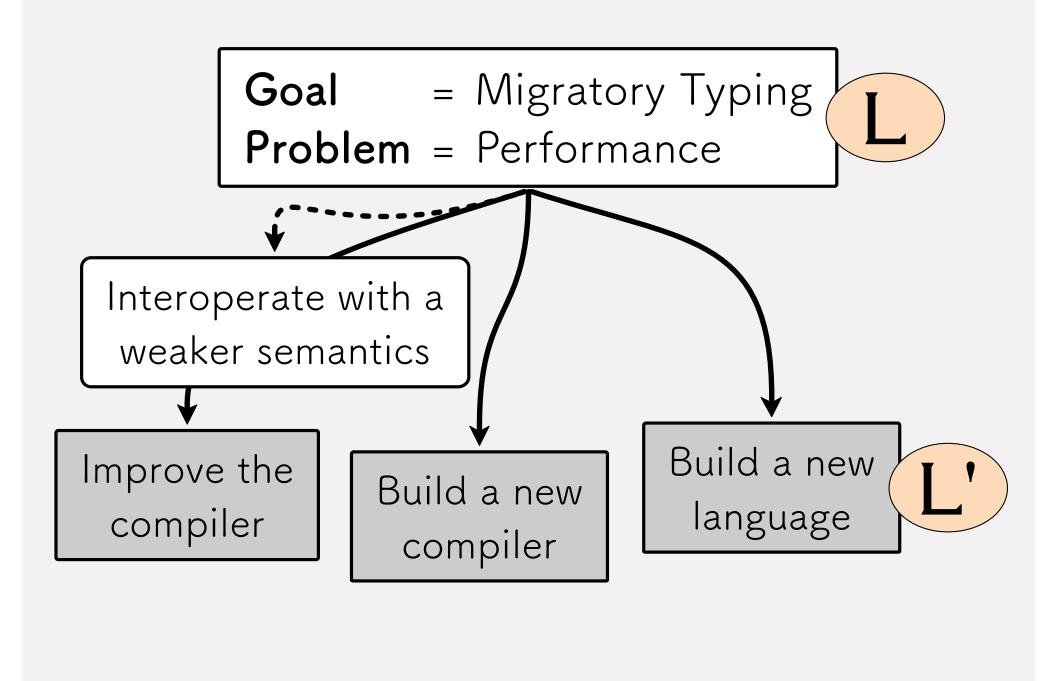
Goal = Migratory Typing Problem = Performance



Improve the compiler







Q Does migratory typing benefit from a combination of honest and lying types?

O Does migratory typing benefit from a combination of honest and lying types?

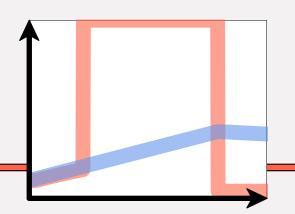
In particular,

Natural + Transient

Complementary Strengths

Natural

types predict full behavior, but need to avoid certain boundaries



Transient

types predict shapes, but add overhead to all typed code

Benefits (1/3): Migration



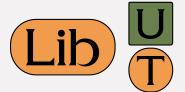






- 1. Begin with **Natural** types
- 2. Switch to **Transient** for performance
- 3. Revisit Natural for debugging
- 4. Return to **Natural** after typing all critical boundaries

Benefits (2/3): Library Interaction

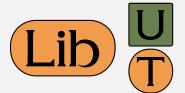








Benefits (2/3): Library Interaction

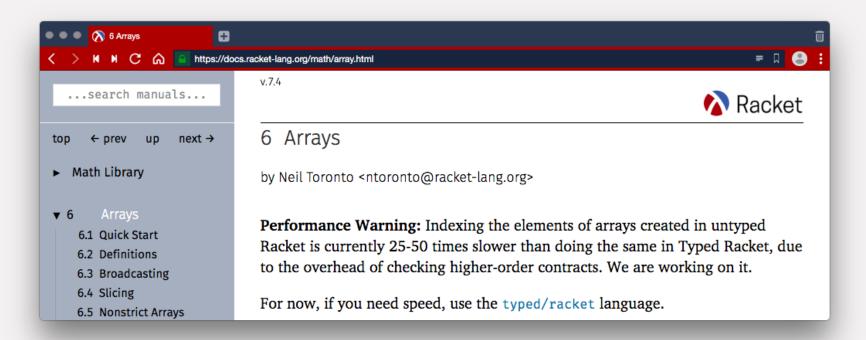








math/array: "25 to 50 times slower"



Benefits (2/3): Library Interaction









Changing a library to **Transient** may improve overall performance





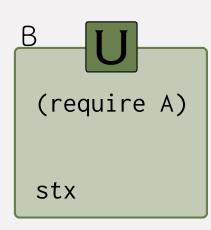




```
(define stx

#`#,(vector 0 1))

(provide stx)
```







```
(define stx

#`#,(vector 0 1))

stx

(provide stx)
```

Type Check: Ok





```
(define stx

#`#,(vector 0 1))

stx

(provide stx)
```

Type Check: Ok

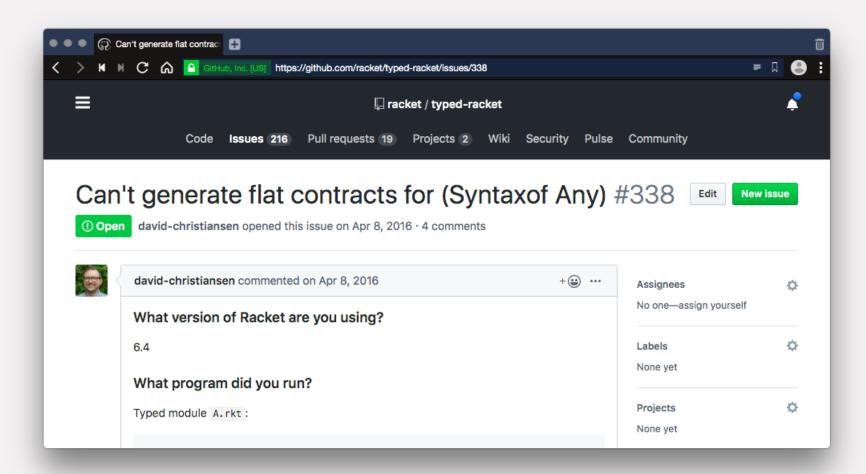
Runtime: Error could not convert type to a contract



















Typed Racket provides 203 base types; 12 lack runtime support (wrappers)









Typed Racket provides 203 base types; 12 lack runtime support (wrappers)

(Async-Channel T) (Custodian-Box T) (C-Mark-Key T)
 (Evt T) (Ephemeron T) (Future T)
 (MPair T T') (MList T) (Prompt-Tag T T')
 (Syntax T) (Thread-Cell T) (Weak-Box T)









Typed Racket provides 203 base types; 12 lack runtime support (wrappers)

Transient does not need wrappers, so more code can run



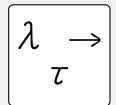
Q Does migratory typing benefit from a combination of honest and lying types?

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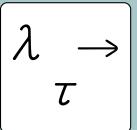
- Q1. Can honest and lying types coexist?
- Q2. Are the benefits measurably significant?

Q1. Can honest and lying coexist?

Model:

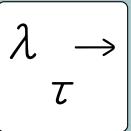


- develop a combined model
- formally prove basic properties
- reduce overlap in runtime checks





Q1. Can honest and lying coexist?





Q1. Can honest and lying coexist?

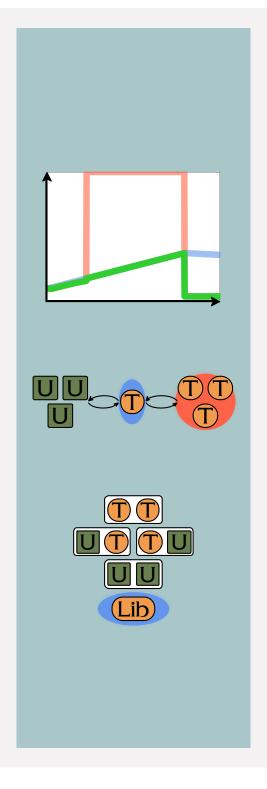
Implementation:



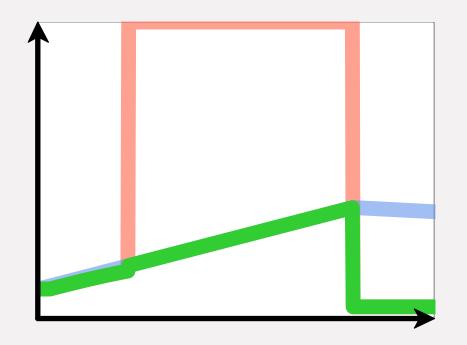


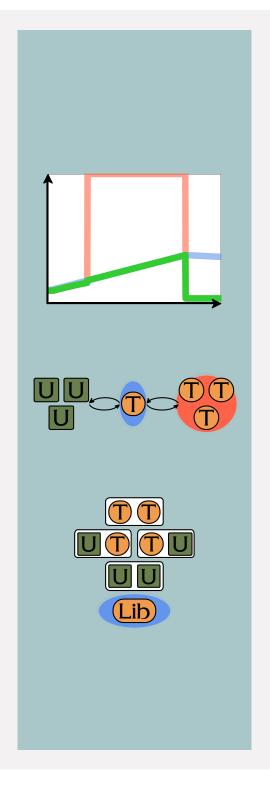
- re-use the type checker
- support all Racket values
- avoid the contract library
- adapt the TR optimizer to lying types

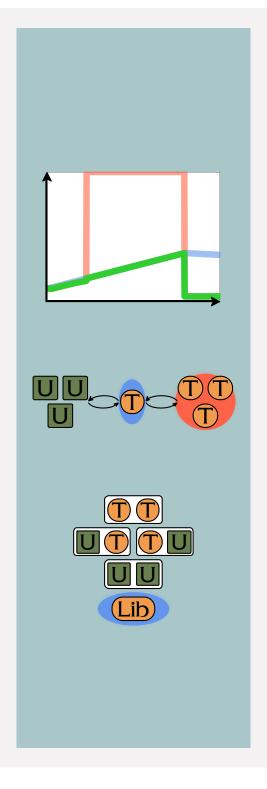




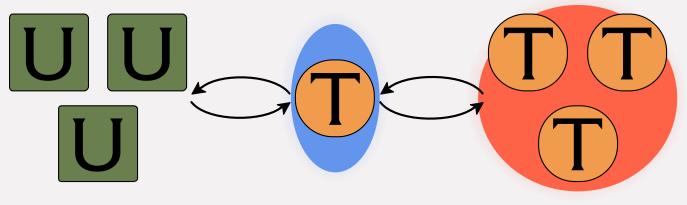
Goal: min(Natural, Transient)

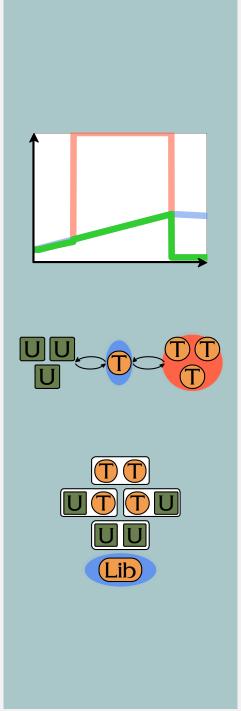




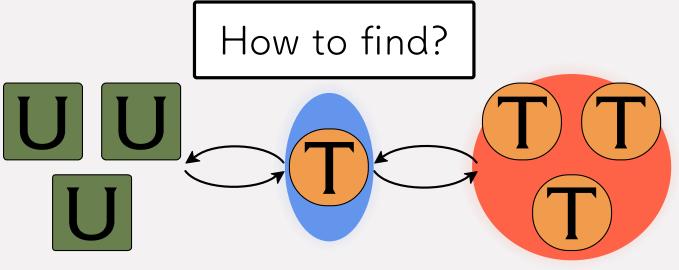


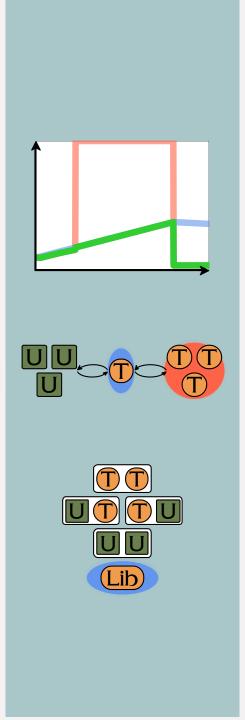
Maybe: reduce cost of U/T edge





Maybe: reduce cost of U/T edge

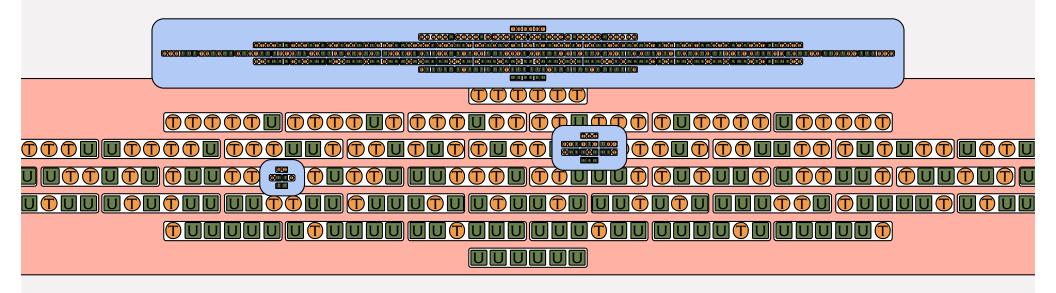




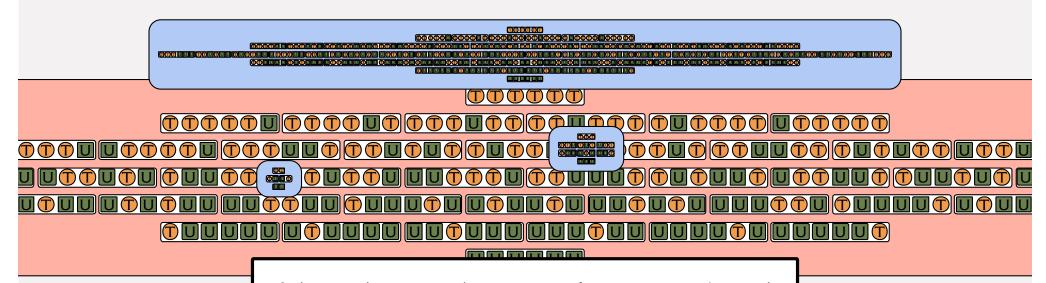
$$(POPL\ 2016) = 2^{N}$$
 measurements

$$|CFP|^{2018} = 2^{(N+1)}$$
 measurements

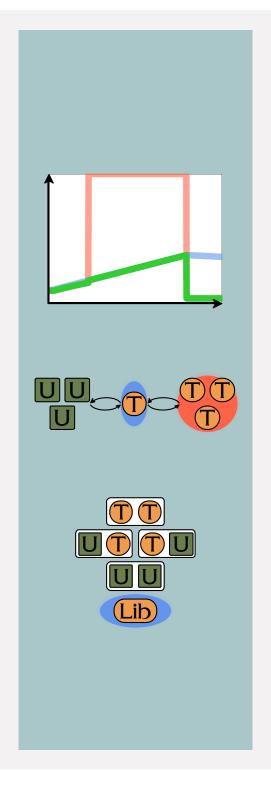
Next = 3^{N} measurements?



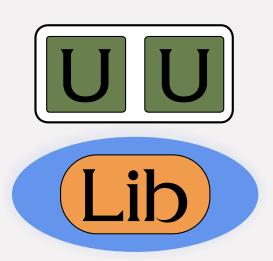
Next = 3^{N} measurements?

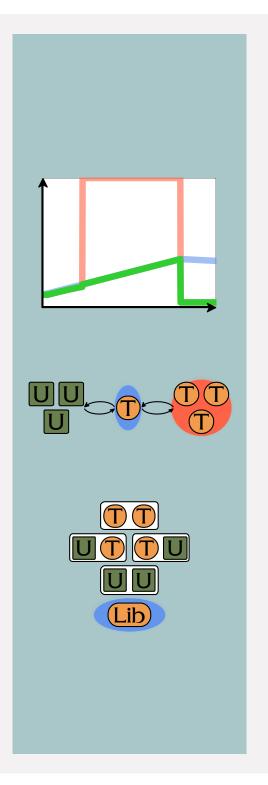


Need an alternative method to measure performance

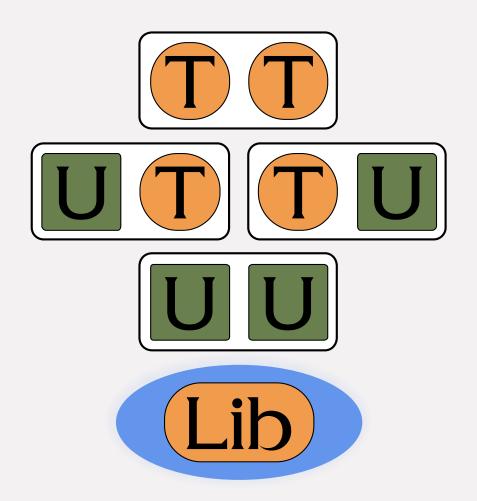


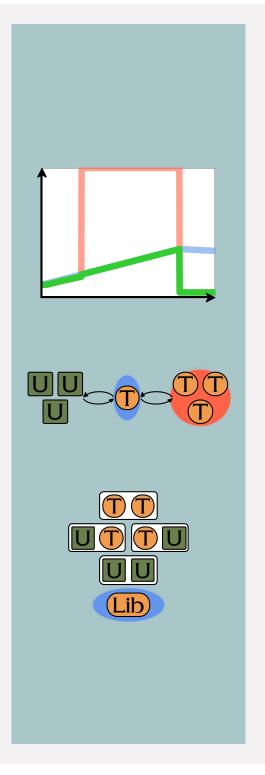
Goal: change lib, improve overall

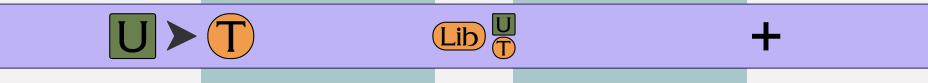


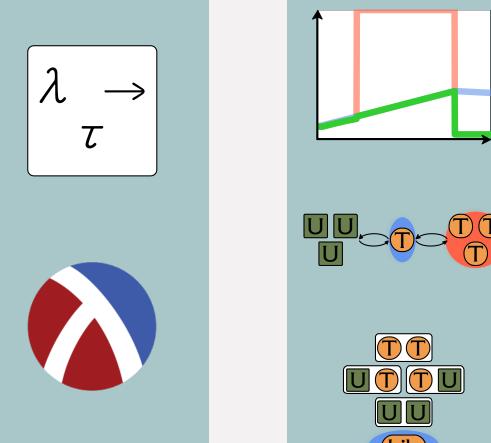


Goal: change lib, improve overall



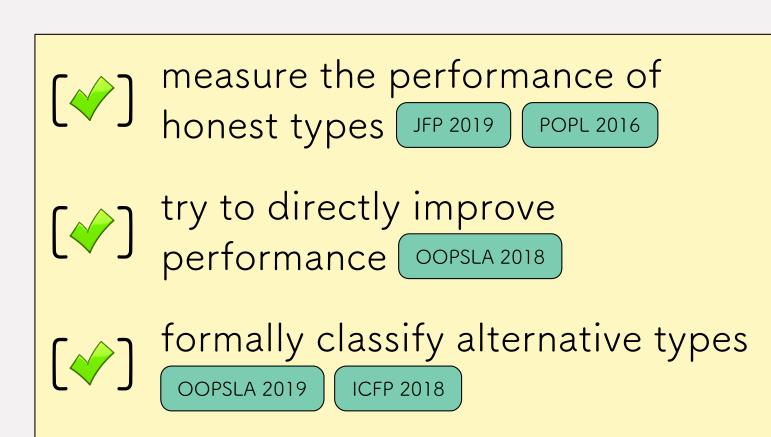




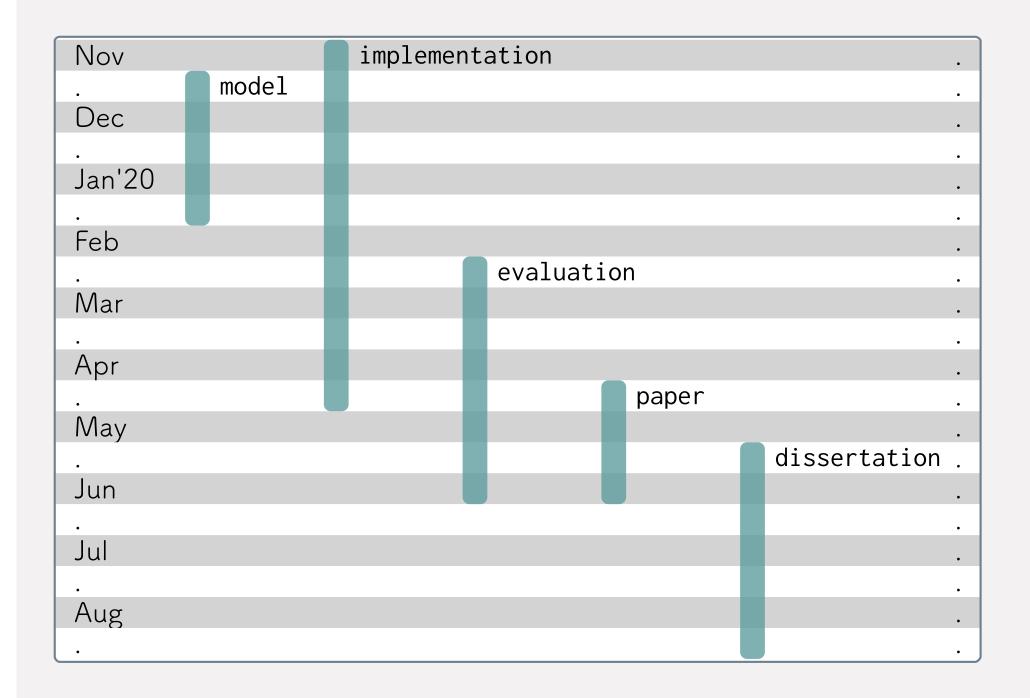


Timeline

```
[ ] measure the performance of honest types
[ ] try to directly improve performance
        formally classify alternative types
develop a combined model, measure combined performance
```



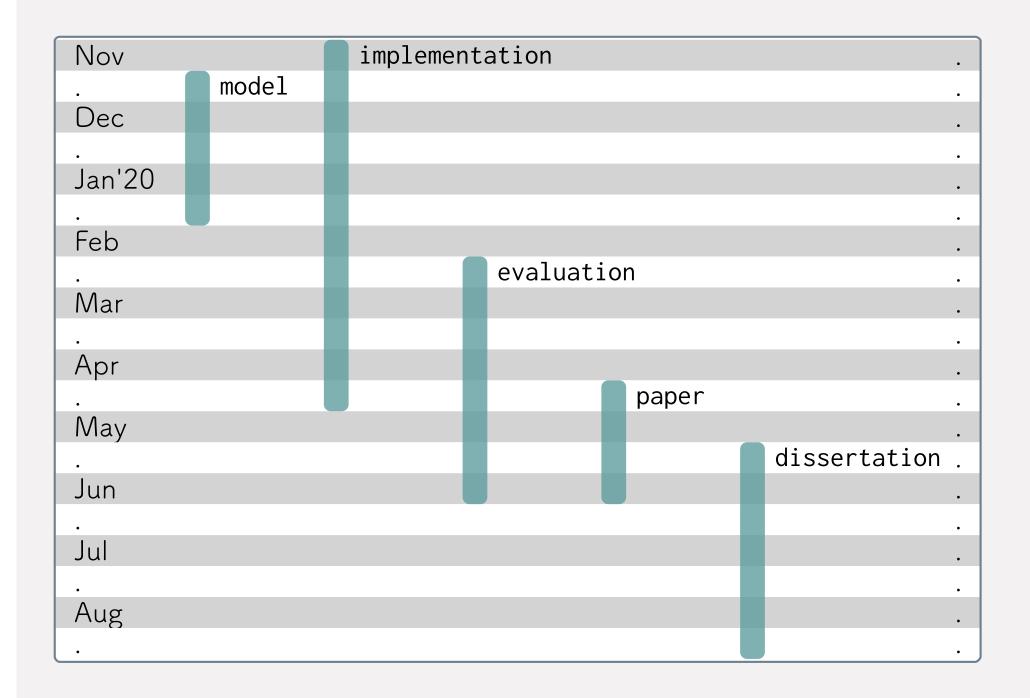
develop a combined model, measure combined performance



Nov		implementation .					
	model						
Dec							
Jan'20							
Feb			_				
			evaluati	.on			
Mar							
Apr							
				paper			
May							
					di	ssertatio	on .
Jun							
Jul							
Aug							
<u>. </u>							

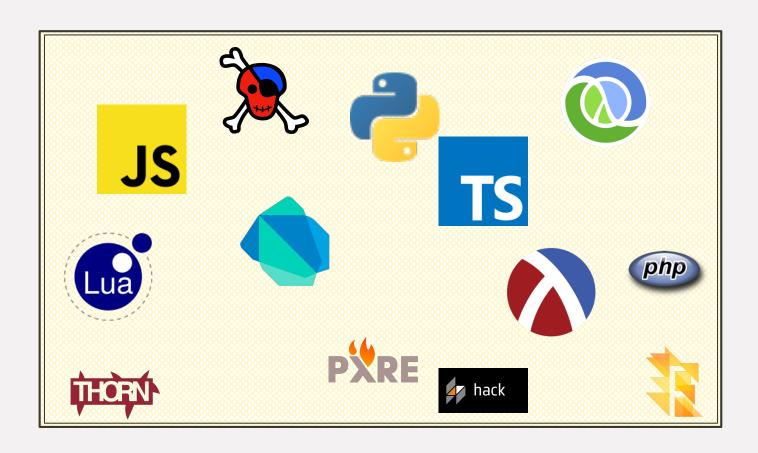
Nov		implemer	implementation .				
	model					.	
Dec							
Jan'20							
Feb							
			evaluatio	on			
Mar							
Apr							
				paper			
May							
					dissertation		
Jun							
Jul							
Aug							
<u> </u>							

```
Timeline
 #lang typed/racket/base #:locally-defensive
(provide make-timeline)
(require typed/racket/class typed/racket/draw typed/pict)
 (require/typed ppict/2
  F#:opaque Coord refpoint-placer?7
  [coord (-> Real Real Symbol Coord)])
 (require/typed "ppict-simple.rkt"
  [ppict (-> Pict (Listof (Pairof Coord Pict)) Pict)])
 (require/typed pict-abbrevs
  [add-rounded-border
         [#:radius Real #:y-margin Real #:frame-width Real #:frame-color String]
         Pict)])
(define-type Pict pict)
 (: make-timeline-bar (-> Real Real (U #f String) (-> String Pict) Pict))
 (define (make-timeline-bar w h label tcodesize)
  (define color (if label "light gray" "white"))
  (define bar (filled-rounded-rectangle w h 1 #:color color #:draw-border? #f))
  (ppict
    (list (cons (coord 2/100 48/100 'lc) (tcodesize (or label ".")))
          (cons (coord 98/100 48/100 'rc) (tcodesize ".")))))
 (: make-timeline-span : (-> Real String (-> String Pict) (Instance Color%) Pict))
 (define (make-timeline-span h label ct timeline-span-color)
  (define span-radius 7)
  (define bar-pict (filled-rounded-rectangle 25 h span-radius #:color timeline-span-color #:draw-border? #f))
  (define label-pict (ct label))
  (ht-append 10 bar-pict label-pict))
 (: make-timeline (-> Real Real (Instance Color%) (-> String Pict) (-> String Pict) Pict))
 (define (make-timeline w h timeline-span-color ct tcodesize)
  (let* ((month*
           '("Nov" "Dec" "Jan'20" "Feb" "Mar" "Apr" "May" "Jun" "Jul" "Aug"))
         (bar-h
           (/ h (* 2 (length month*))))
         (make-span-h
           (lambda ((i : Real)) (* i bar-h)))
         (make-span-%
           (lambda ((i : Real)) (/ (make-span-h i) h)))
         (base
           (for/fold : Pict
                    ((acc : Pict (blank)))
                     ((m : String (in-list month*)))
            (vl-append 0 acc
                       (make-timeline-bar w bar-h m tcodesize)
                       (make-timeline-bar w bar-h #f tcodesize))))
         (timeline
           (ppict
            base
              (cons (coord 14/100 (make-span-% 1) 'lt) (make-timeline-span (make-span-h 5) "model" ct timeline-span-color))
               (cons (coord 29/100 0 'lt) (make-timeline-span (make-span-h 12) "implementation" ct timeline-span-color))
               (cons (coord 44/100 (make-span-% 7) 'lt) (make-timeline-span (make-span-h 8) "evaluation" ct timeline-span-color))
               (cons (coord 59/100 (make-span-% 11) 'lt) (make-timeline-span (make-span-h 4) "paper" ct timeline-span-color))
               (cons (coord 74/100 (make-span-% 13) 'lt) (make-timeline-span (make-span-h 7) "dissertation" ct timeline-span-color))))))
    (add-rounded-border
      #:radius 5 #:y-margin 6 #:frame-width 3 #:frame-color "slategray"
      timeline)))
```



The End

Q Does migratory typing benefit from a combination of honest and lying types?



Complete Monitoring

types predict behavior

Type Soundness

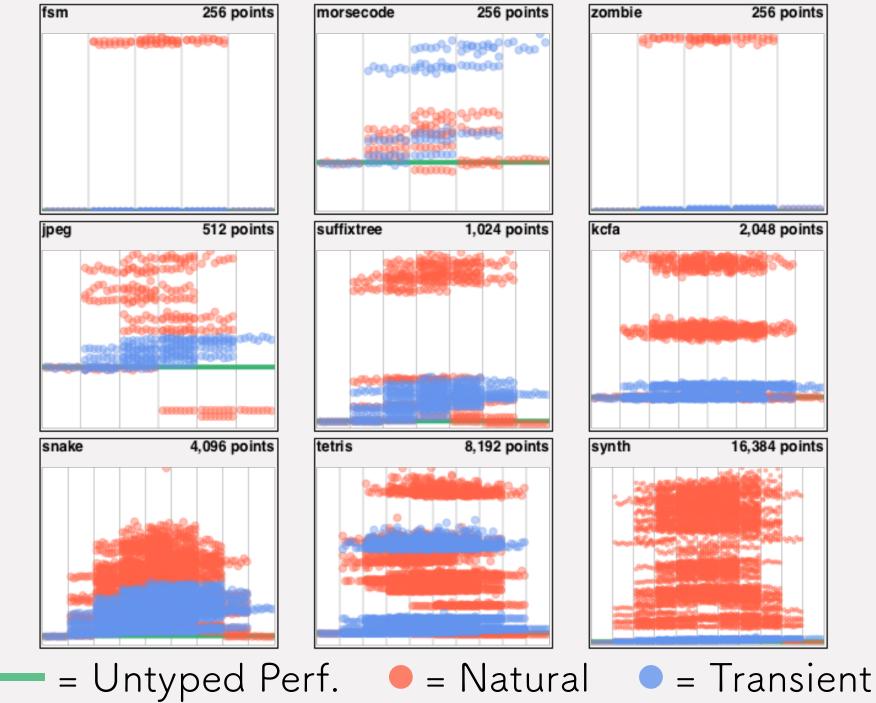
types predict behavior in typed code, nothing in untyped code

Tag Soundness

types predict shapes in typed code, nothing in untyped code

Dyn Soundness

types predict nothing

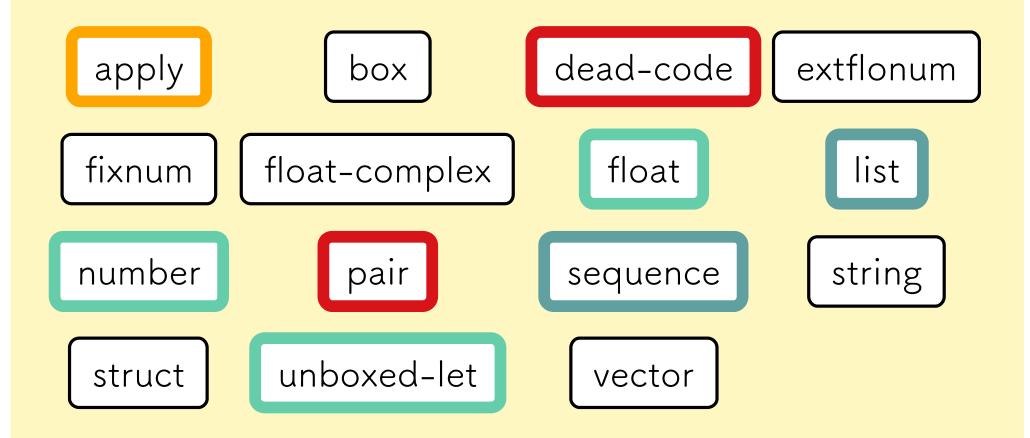


Expressiveness

DLS 18

Preston Tunnell Wilson, Ben Greenman, Justin Pombrio, and Shriram Krishnamurthi

TR Optimizations



dead-code = unsafe for Transient

```
(: g (-> Str Str))
(define g
    (case-lambda
       [(x) x]
       [(x y) y]))
```



```
(define g
  (case-lambda
  [(x) x]
  [(x y) (void)]))
```

Problem: untyped code can call (g 0 1)

pair = unsound for Transient

```
(: x (Pairof (Pairof Nat Int) Str))
(cdar x)
```



```
(unsafe-cdr (unsafe-car x))
```

Problem: no guarantee (car x) is a pair

apply = safe but risky for Transient

```
(: h (-> Str Str))
(: xs (Listof Str))
(apply + (map h xs))
```



```
(+ (h (unsafe-car xs))
  (h (unsafe-car (unsafe-cdr xs))) ...)
```

Caution: h must check inputs

list sequence = force choice for LTJ

```
(: xs (List Str Str))
(list-ref xs 1)
```



(unsafe-list-ref xs 1)

Note: [List Str Str] needs more than a tag check

number = LTJ is more than a tag check

Natural

Exact-Nonnegative-Integer

Nonpositive-Inexact-Real

ExtFlonum-Negative-Zero

unboxed-let = safe with escape analysis

```
(: f (-> Float-Complex Any))
(define (f n)
 . . . . )
```



```
(define (f n-real n-imag)
 . . . . )
```

float = false alarm

(flrandom)



(unsafe-flrandom
 (current-pseudo-random-generator))

Ok because the PRNG parameter checks inputs