
Co-Architecting Controllers and DRAM to Enhance DRAM Process Scaling

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Outline

❖ Introduction

❖ DRAM process scaling challenges

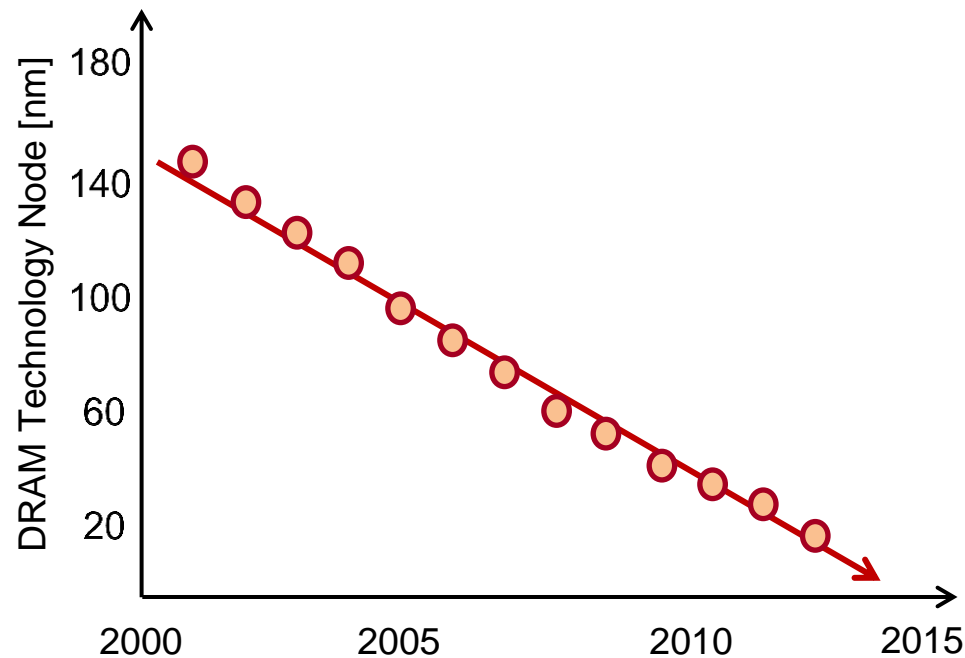
❖ DRAM process scaling enhancing features

- Sub-array level parallelism with tWR relaxation
- Temperature compensated tWR
- In-DRAM ECC with dummy data pre-fetching

❖ Conclusion

Introduction

- ❖ DRAM process scaling enables continued density and speed increase, power decrease, and bit-cost reduction
- ❖ Today's DRAM process is expected to continue scaling, enabling minimum feature sizes below 10nm
- ❖ Main challenges to address are expected to be refresh, write recovery time (tWR), and variable retention time (VRT) parameters



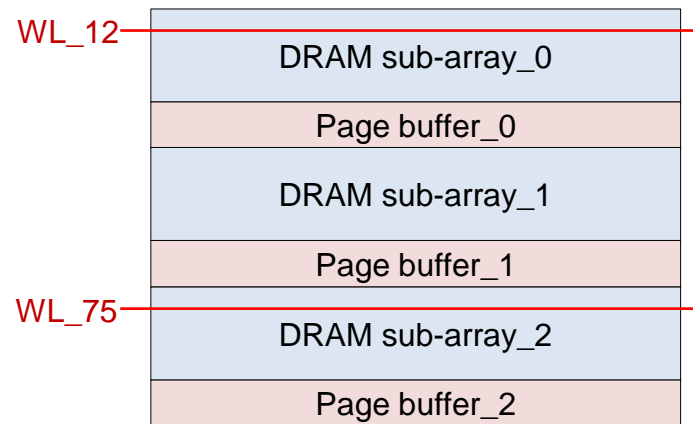
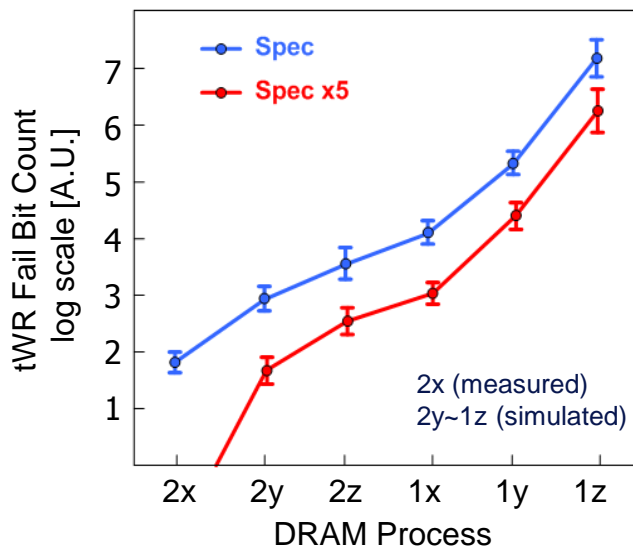
Sub-array Level Parallelism with tWR Relaxation

❖ tWR relaxation

- Relaxing tWR results in DRAM yield improvement but can degrade performance requiring new compensating features
- By increasing tWR 5X (from 15ns to 75ns), fail bit counts are expected to reduce by 1 to 2 orders of magnitudes

❖ Sub-array level parallelism (SALP)

- Allows a page in another sub-array in the same bank to be opened in parallel with the currently activated sub-array
- Results in performance gain by increasing the row access parallelism within a bank
⇒ Used to compensate for the performance loss caused by tWR relaxation



Single bank with multiple sub-arrays

Sub-array Level Parallelism Timing Gain

❖ Read

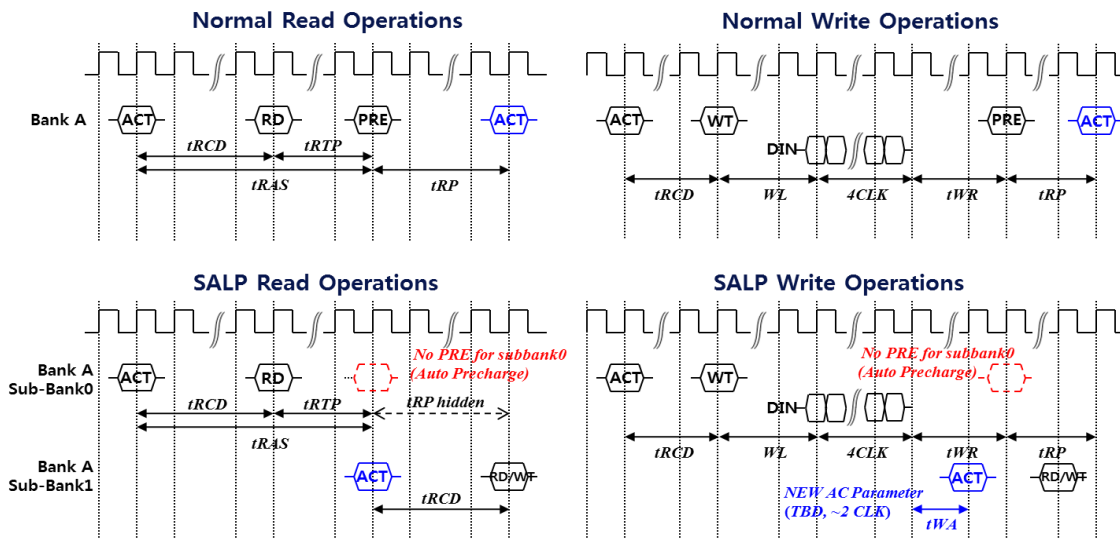
- Next activate command to a different sub-array in the same bank can come after tRAS
⇒ Pulled in by tRP compared to the normal case

❖ Write

- Next activate command to a different sub-array can come after tRCD+WL+4tck+tWA
⇒ Pulled in by tRP+tWR-tWA compared to the normal case

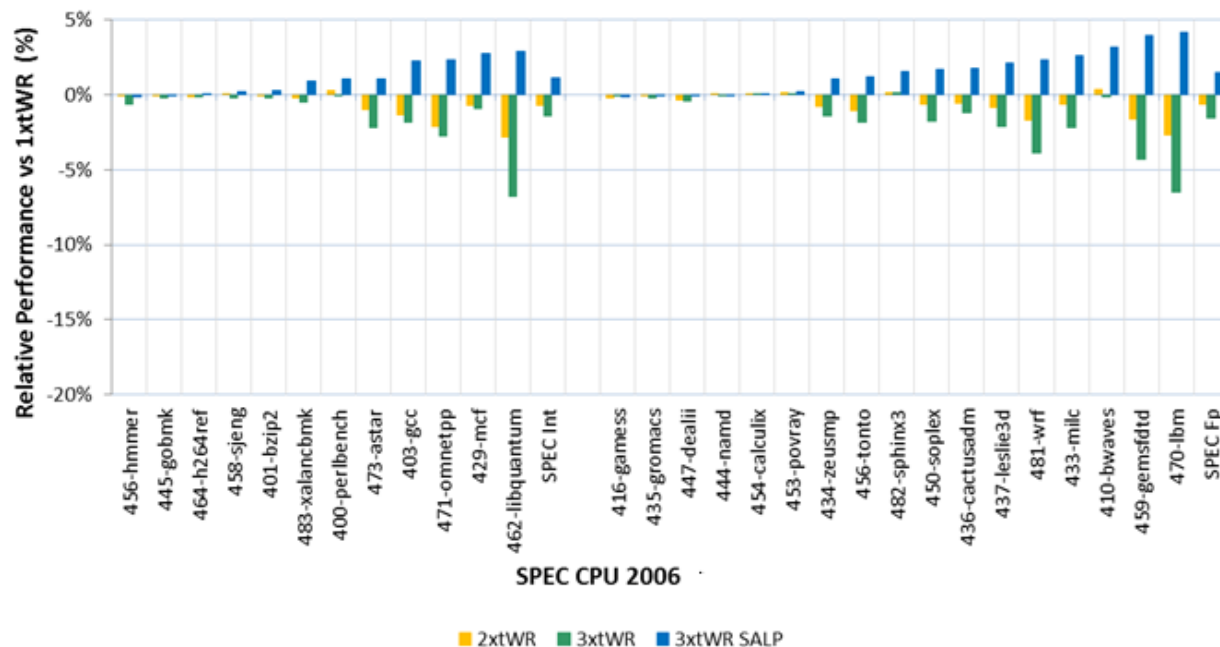
❖ Mutually exclusive sub-array activation

- All open pages in other sub-arrays are closed automatically once a new activate request to a closed sub-array page is made ⇒ simplifies operation



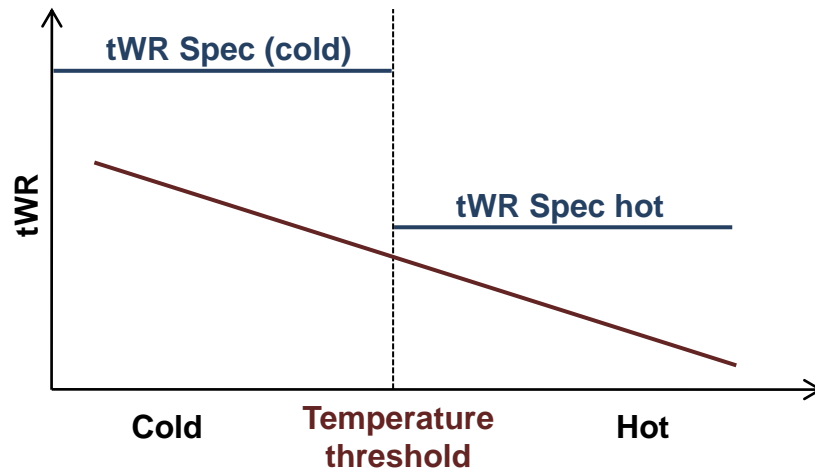
Performance Impact of SALP and tWR relaxation

- ❖ Performance simulations run for various workloads when tWR is relaxed by 2X and 3X, and when SALP is applied with 2 sub-banks
- ❖ Results show that performance is reduced by ~5% and ~2% in average if tWR is relaxed by 3X and 2X, respectively
- ❖ Results also show that performance is compensated, and even improved to up to ~3% in average when SALP is applied, even with tWR relaxed by 3X



Temperature Compensated tWR

- ❖ Specifies different tWR values below and above a certain given temperature threshold
 - Longer tWR spec at cold temperature since tWR is worse at cold
 - Relaxing tWR spec at cold temperature helps to increase DRAM yield
 - DRAM temperature is checked by the memory controller periodically
- ❖ Combining temperature compensated tWR with SALP can improve performance when ambient system operation temperatures exceed the temperature threshold most of the time

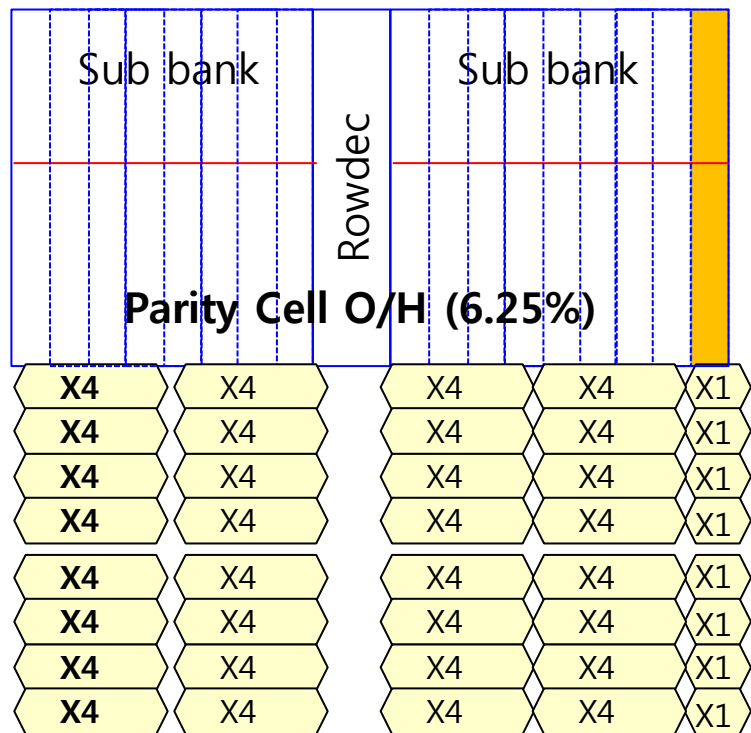


	<Temp Threshold	>Temp Threshold
tWR Spec	m+n [ns]	m [ns]

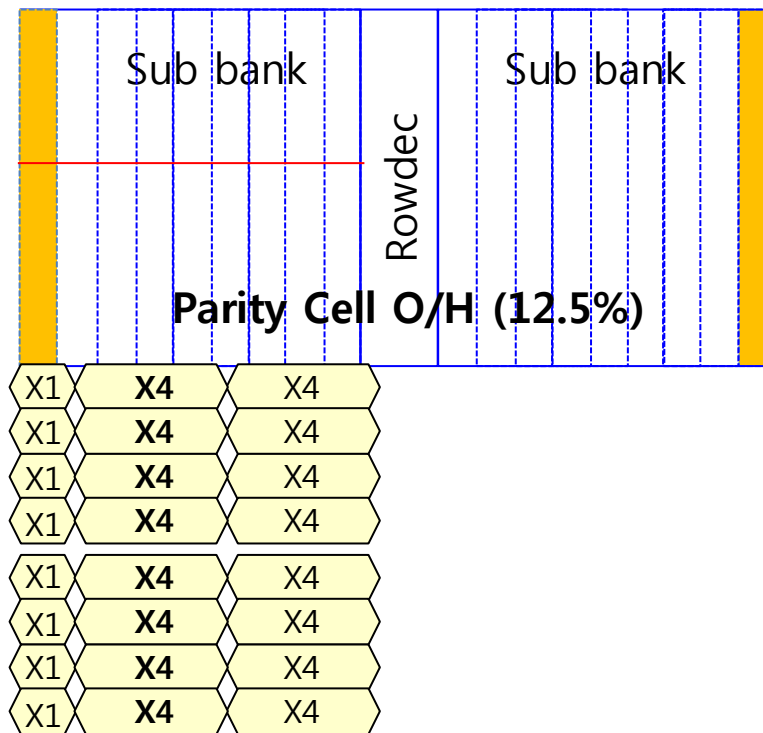
In-DRAM ECC with Dummy Data Pre-fetching

- ❖ In-DRAM ECC is known to have high repair efficiency in repairing single fail bits, but has in general large chip size overhead (~20% for X4 DRAM)
- ❖ Proposing a method where extra dummy data bits are pre-fetched internally during reads and writes, resulting in parity bit chip size overhead reduction
 - Ex) External: $X4 * BL8 = 32b$ | Internal: $X4 * BL8 * 4 = 128b+8b$ (Parity O/H: 6.25%)
 - Internal: $X4 * BL8 * 2 = 64b+8b$ (Parity O/H: 12.5%)

ECC codeword : 128bit + 8bit

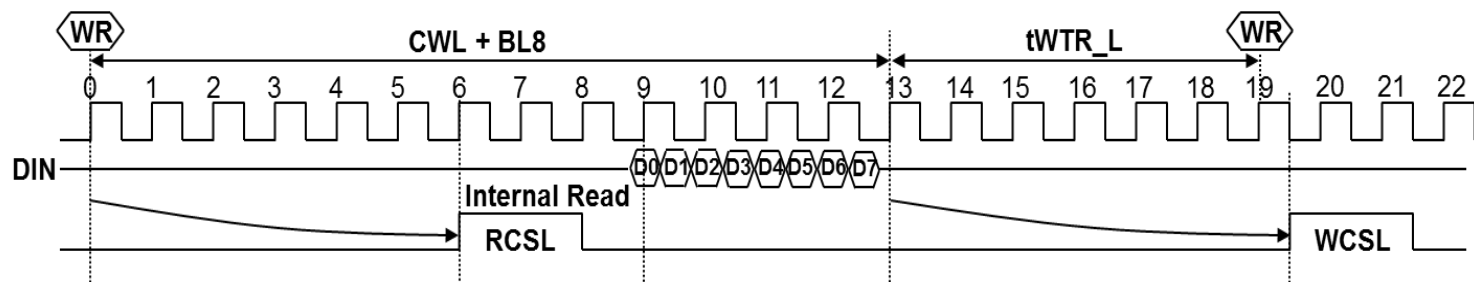


ECC codeword : 64bit + 8bit



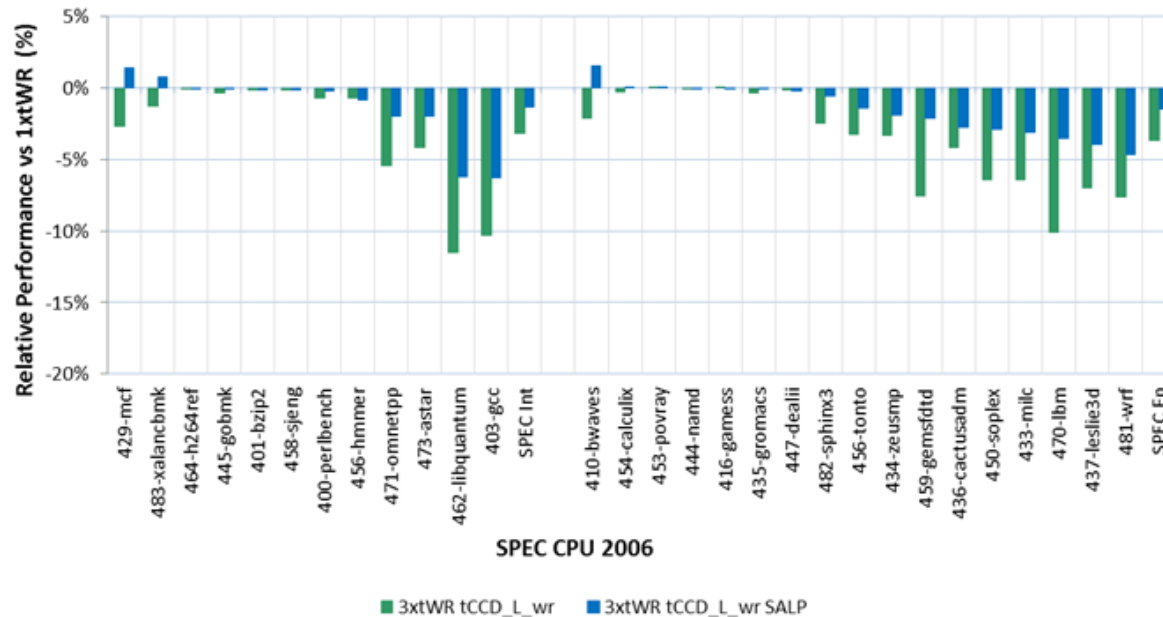
In-DRAM ECC Timing

- ❖ Dummy data bit pre-fetching requires read-modify-write for all writes
 - Only parts of the data code words are overwritten, requiring parity bit update for every write
- ❖ Read-modify-write requires increase in $t_{CCD_L_WR}$ by “ $CWL + BL8 + t_{WTR_L}$ ”
 - t_{WTR_L} needs to be observed from the last write-data burst to the next write command within the same bank group, since the next write actually starts with an internal read
 - $t_{CCD_L_WR}$ needs to be increased from 6~9 clock cycles to 25~29 clock cycles



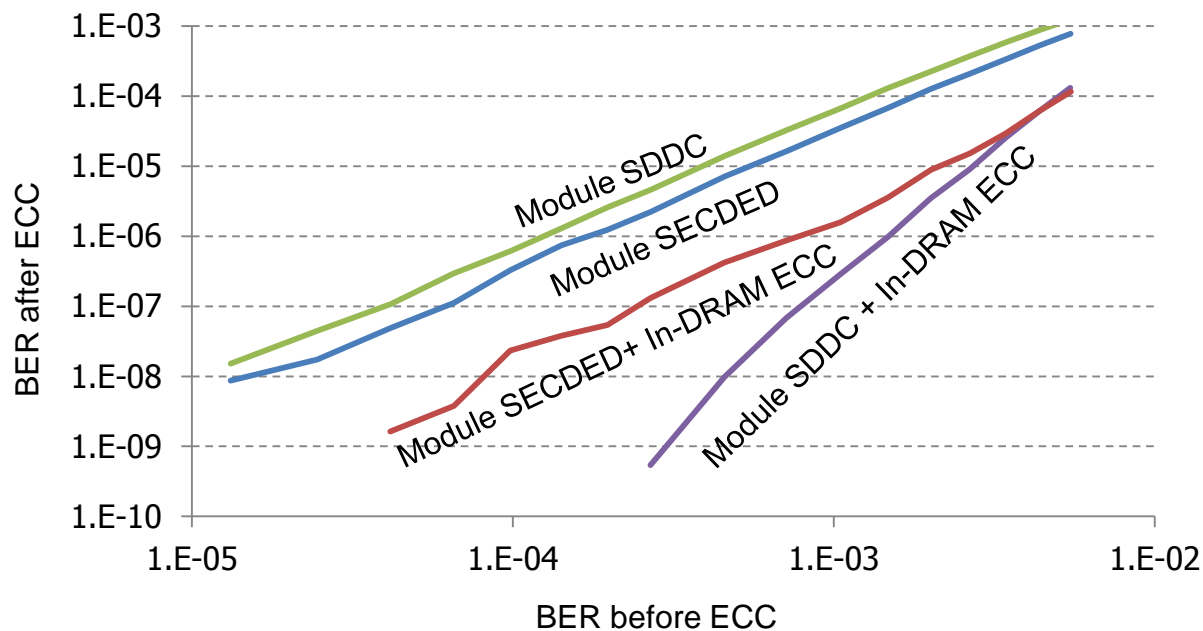
Performance Impact of In-DRAM ECC

- ❖ Performance simulations run when tWR is relaxed 3X, tCCD_L_WR is increased to 24tck, and SALP is used to compensate for the performance loss
 - Assumption: tCCD_L_WR=24tck, tWR 3X, DDR4-2400, 6 channels, 1 DIMM/channel, 2 rank/DIMM
- ❖ Results show an average performance loss of ~2% with SALP and tWR relaxed 3X
- ❖ By applying in-DRAM ECC only without increasing tWR, performance loss is expected to break even when applying SALP



Field Reliability and in-DRAM ECC

- ❖ Field failure rate probability analysis with and without in-DRAM ECC, and with SECDED and SDDC (chip-kill) were performed
 - Assumption: 4Gb X4 DRAM, single bit failures only, errors randomly scattered
- ❖ Results show that the bit error rate decreases by orders of magnitudes when in-DRAM ECC is adopted



Conclusion

- ❖ To enhance DRAM process scaling, challenging process parameters including refresh, tWR, and VRT need to be addressed
- ❖ Combining SALP with TCWR enables tWR relaxation and compensation of performance loss caused by this parameter relaxation
- ❖ In-DRAM ECC allows efficient repair of fail bits coping with VRT and refresh failures resulting in increased field reliability
- ❖ By enabling the proposed scaling features, future DRAM process scaling is expected to accelerate further beyond 10nm



Thank you

